

無線網路多媒體系統

Wireless Multimedia System

(Topic 3)

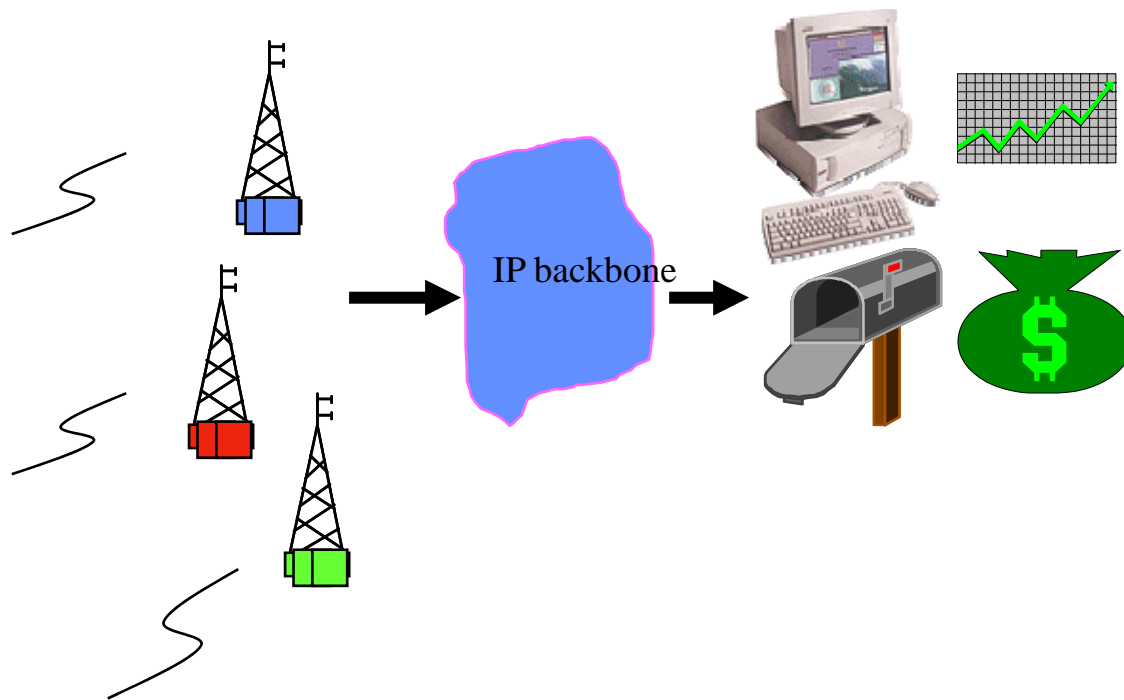
Wireless Link I: Fundamental issues of Modulation
and Multiple Access

吳曉光博士

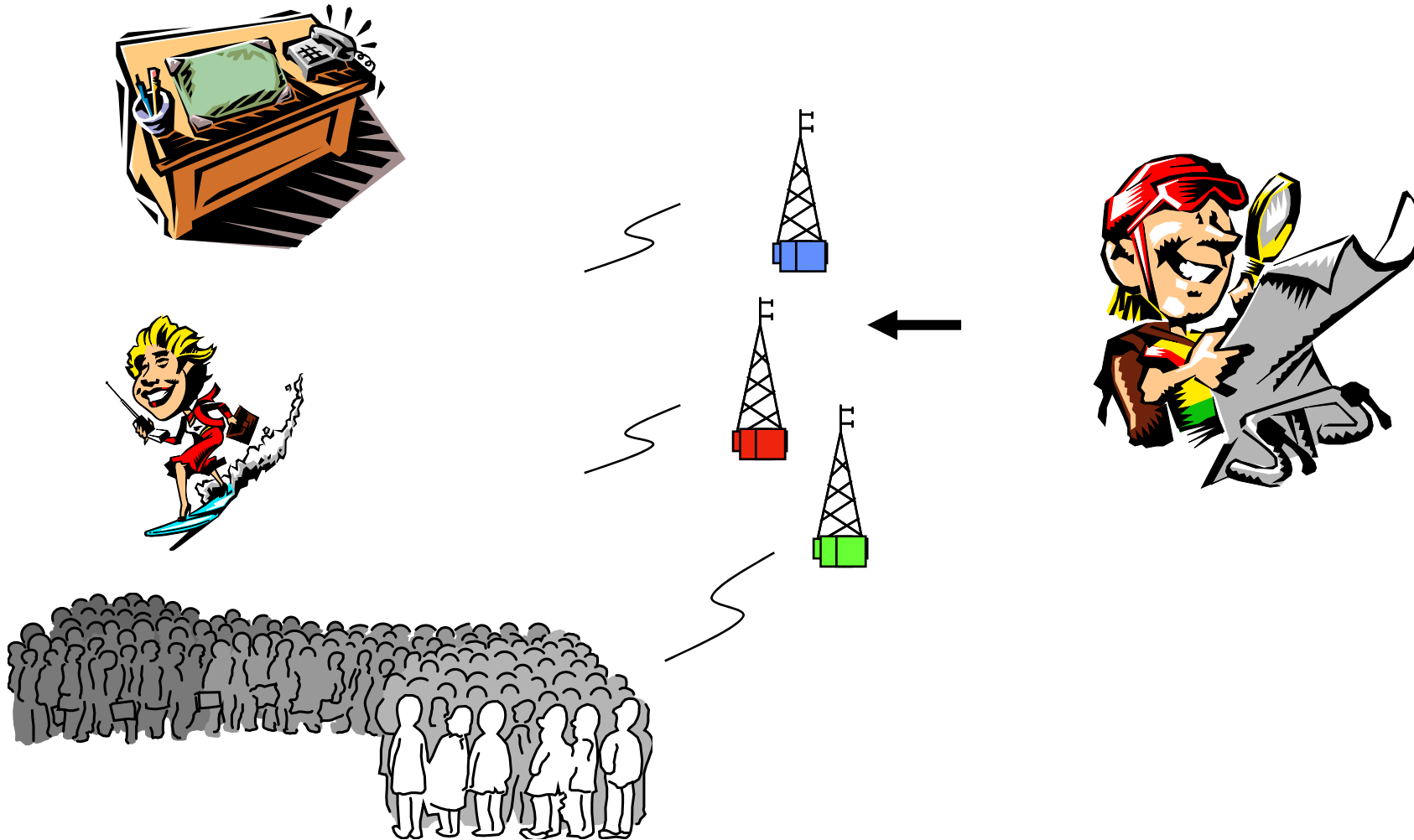
<http://wmlab.csie.ncu.edu.tw>

We
provide
無線網路多媒體實驗室
Wireless
Wireless Network & Multimedia Laboratory
Solution

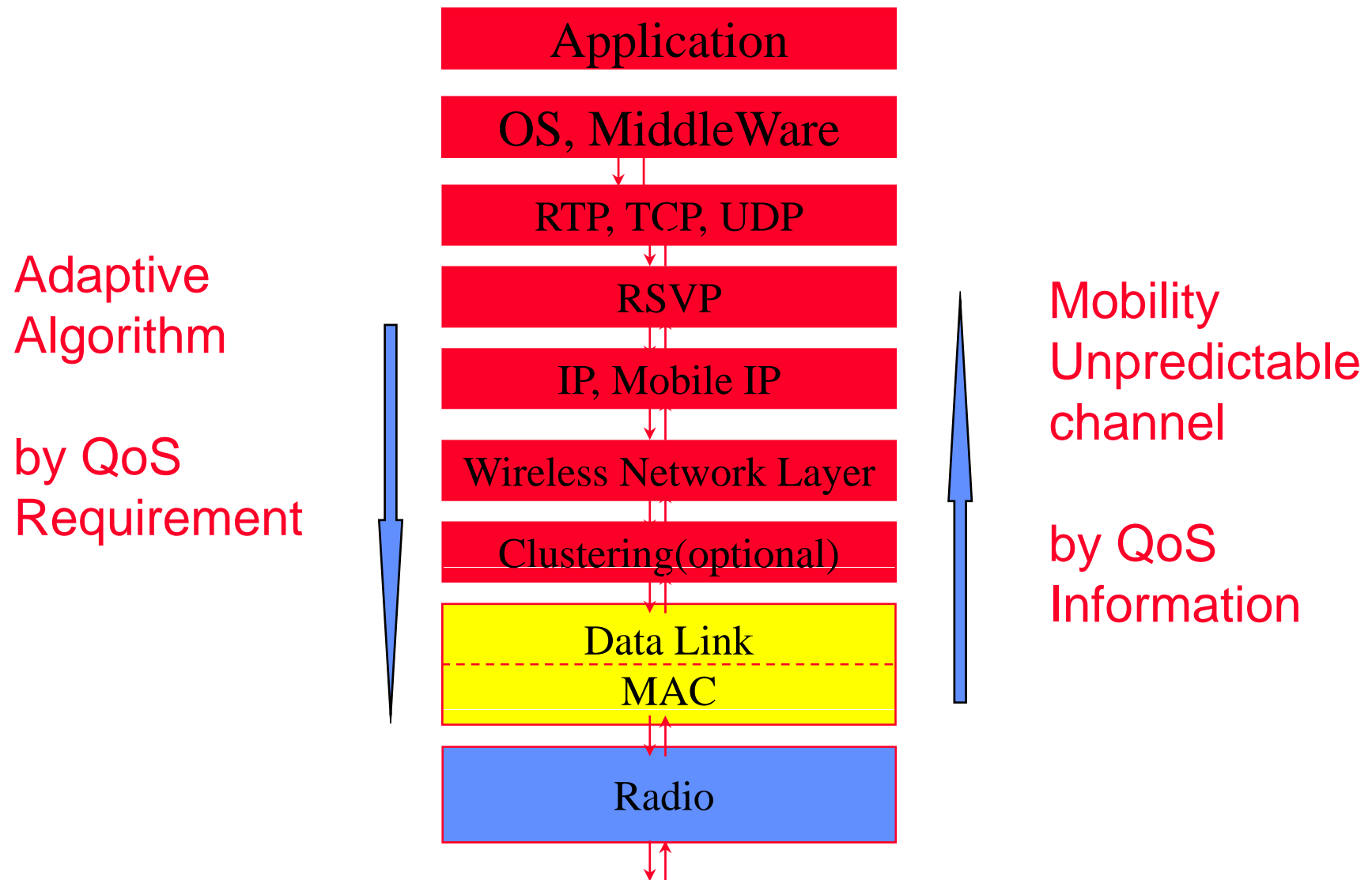
How to deal with Radio Propagation



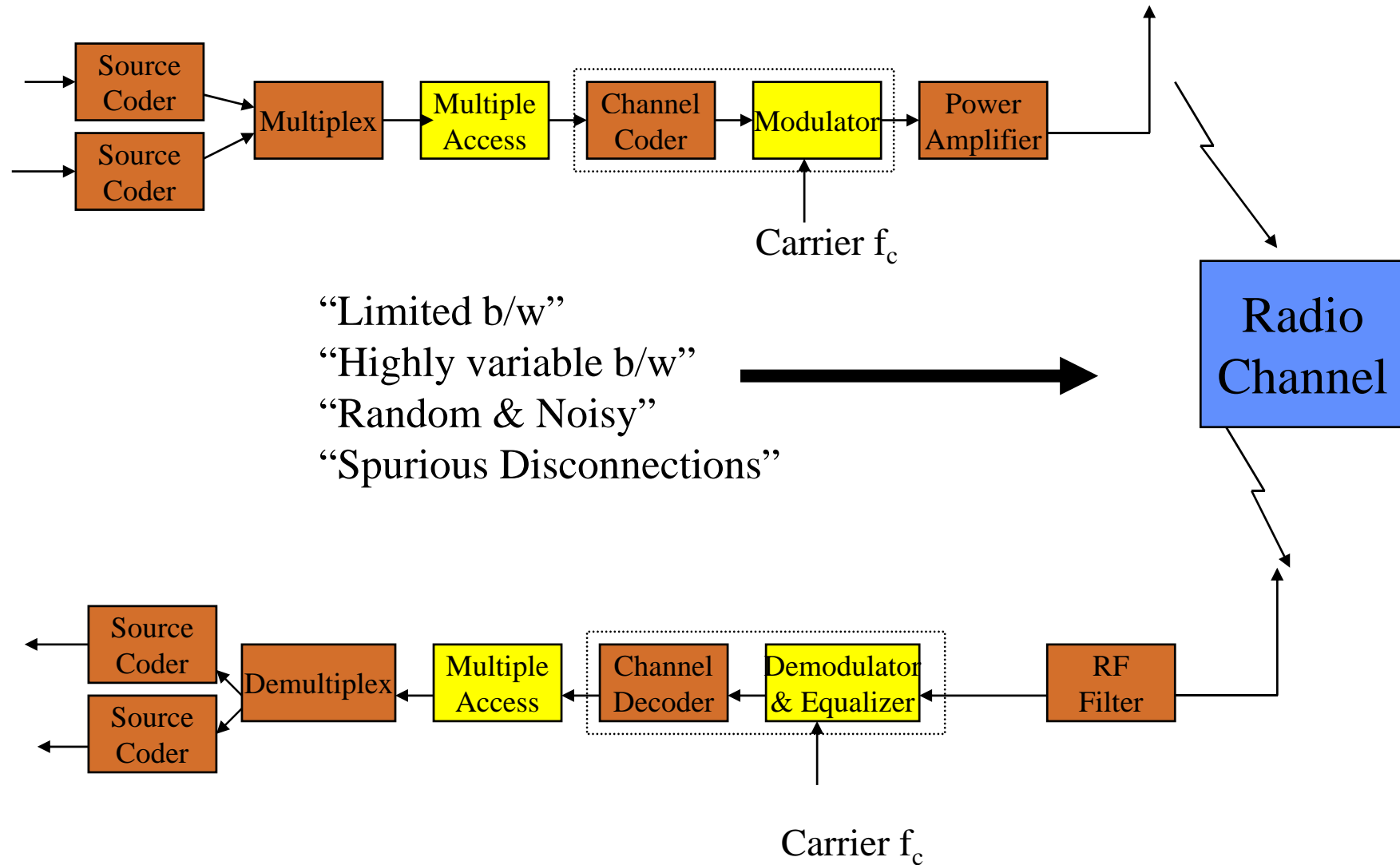
Where are you from?



QoS and Multimedia Traffic Support



Multiple Access & Modulation



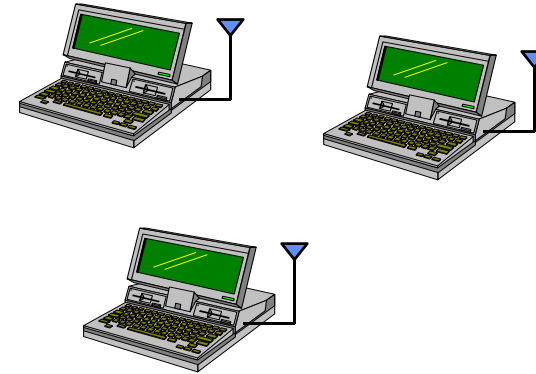
Topic III Agenda

◆ Wireless Link

- Deployment of “Pervasive Computing” and “Seamless Telecom services”
- Channel resource sharing in time, frequency, and code dimensions
- Spread Spectrum-direct sequence, frequency hopping, interference resistance
- Static techniques: TDMA, FDMA, CDMA
- Random access techniques: MACA, MACAW, 802.11 etc



What kind of multiple access environments?



Reading list for This Lecture

◆ Required Reading:

- (Bharghavan94) V. Bharghavan, A. Demers, S. Shenker, L. Zhang, "MACAW: A Medium Access Protocol for Wireless LANs, Proceedings of SIGCOMM'94
- (J.J.97) L. Fullmer and J.J. Garcia-Luna-Aceves, Solutions to Hidden Terminal Problems in Wireless Networks, Proceedings of SIGCOMM'97
- (Jing 2006) J. Zhu, B. Metzler, X. Guo, Y. Liu, "Adaptive CSMA for Scalable Network Capacity in High-Density WLAN: A Hardware Prototyping Approach", Proceedings of Infocom 2006.

Further Reading

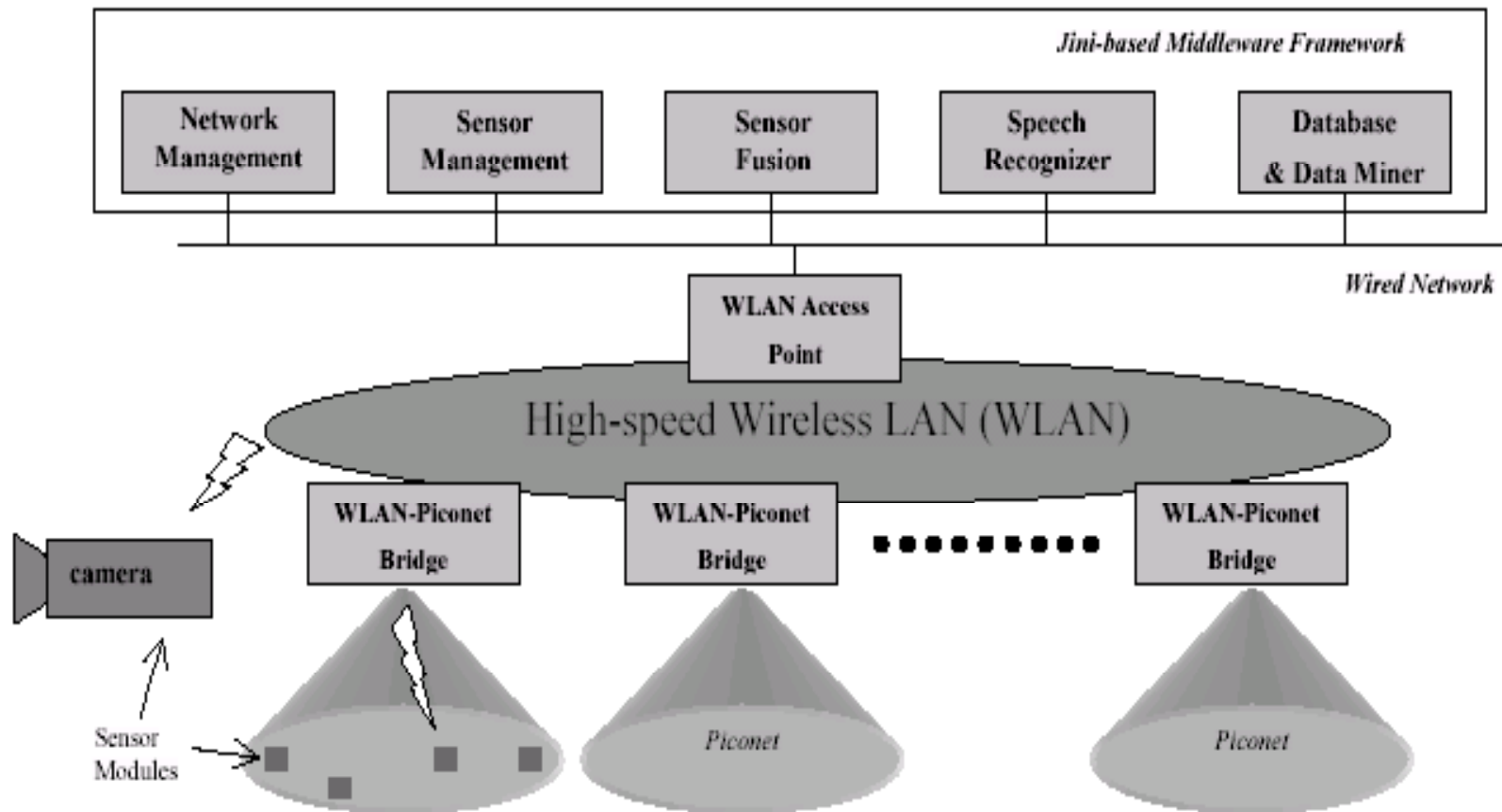
- (David 95) David D. Falconer, F. Adachi, and B. Gudmundson, "Time Division Multiple Access Methods for Wireless Personal Communications", IEEE Communication Magazine January 1995
- (Vadu2000) Vaduvur Bharghavan, "Achieving MAC Layer Fairness in Wireless Packet Networks". IEEE MobileCom2000
- (Songwu Lu2000) Haiyun Luo, Songwu Lu, Vaduvur Bharghavan, "A New Model for Packet Scheduling in Multihop Wireless Networks". IEEE MobileCom2000
- (J.J.2001) L. Bao A New Approach to Channel Access Scheduling for Ad hoc Networks, IEEE MobileCom2001
- (Alex2001) A. Woo, David E. Culler, "A Transmission Control Scheme for Media Access in Sensor Networks", IEEE MobileCom2001
- (Gavin2001) G. Holland, N. Vaidya, P. Bahl, "A Rate-Adaptive MAC Protocol for Multi-Hop Wireless Network, IEEE MobileCom2001

Pervasive Computing Projects

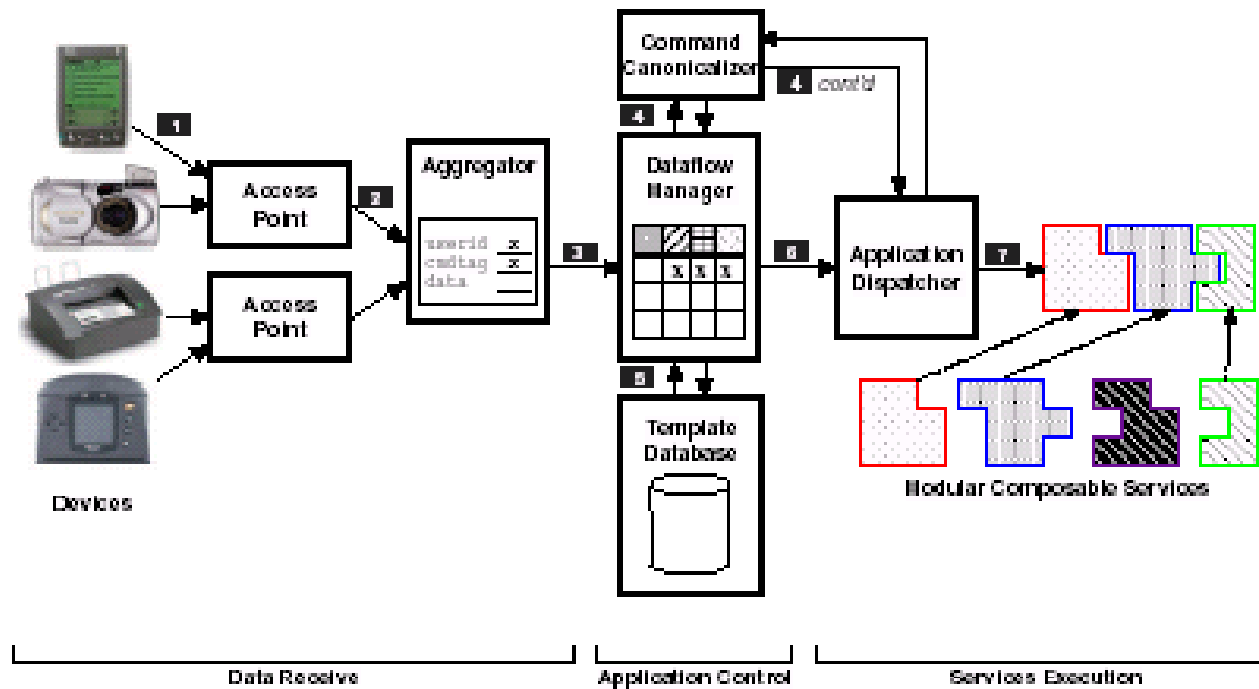


Packet Oriented -> Multimedia Traffic

Smart Kindergarten (UCLA)

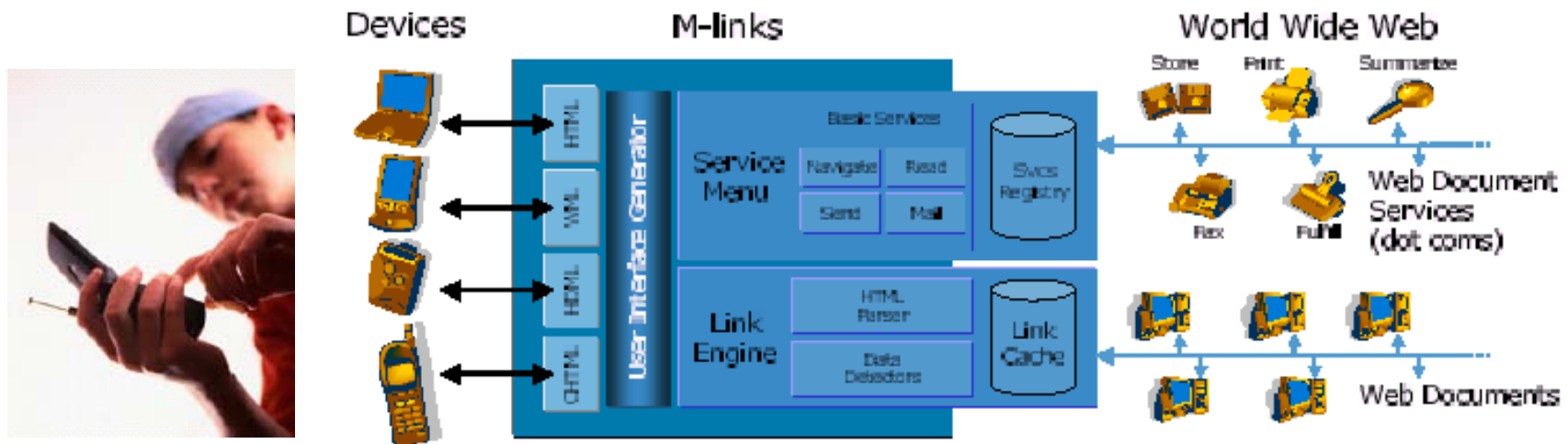


Making Computer Disappear (Stanford) ADS (Appliance Data Services)



- 1 data sent to Access Point
- 2 Aggregator receives data, which completes the triple: (userid, cmdtag, data)
- 3 completed triple sent to Dataflow Manager
- 4 Canonicalizer converts cmdtag to plaintext
- 5 (userid, cmdtag) looked up in Template Database to find matching application template
- 6 application template and data sent to Application Dispatcher when required data is received
- 7 Application Dispatcher invokes services on data as specified in the application template

M-Links (Xerox)



Seamless Telecom Deployments



Circuit Services -> Data Services -> Multimedia

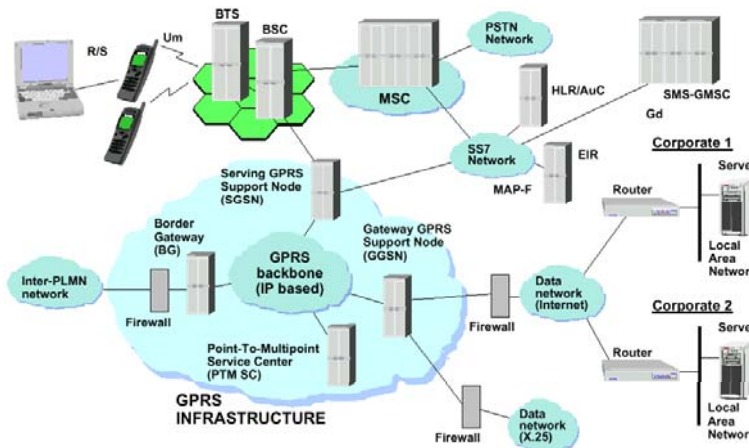
2.5 G & 3 G

Packet Radio



Packet Backbone

System Integration
Multimedia Services
Mobile Computing



Network	Radio	Terminal	Interconnectivity
Servers (IP)	WAP	Applications • Multimedia • Infotainment	Bluetooth
Packet Networks	JAVA	User Interface	
IS-41, GSM MAP Evolution	3G Radio	EPOC 32 SWD Radio	

Wireless Networking Technology

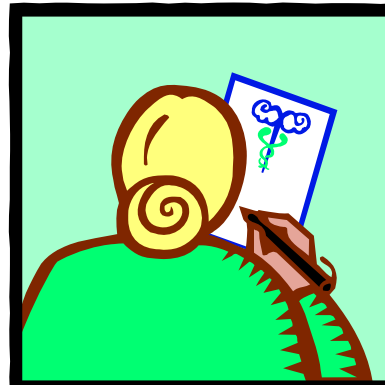


Telecom & Datacom

Circuit & Packet

MAC Design Issues

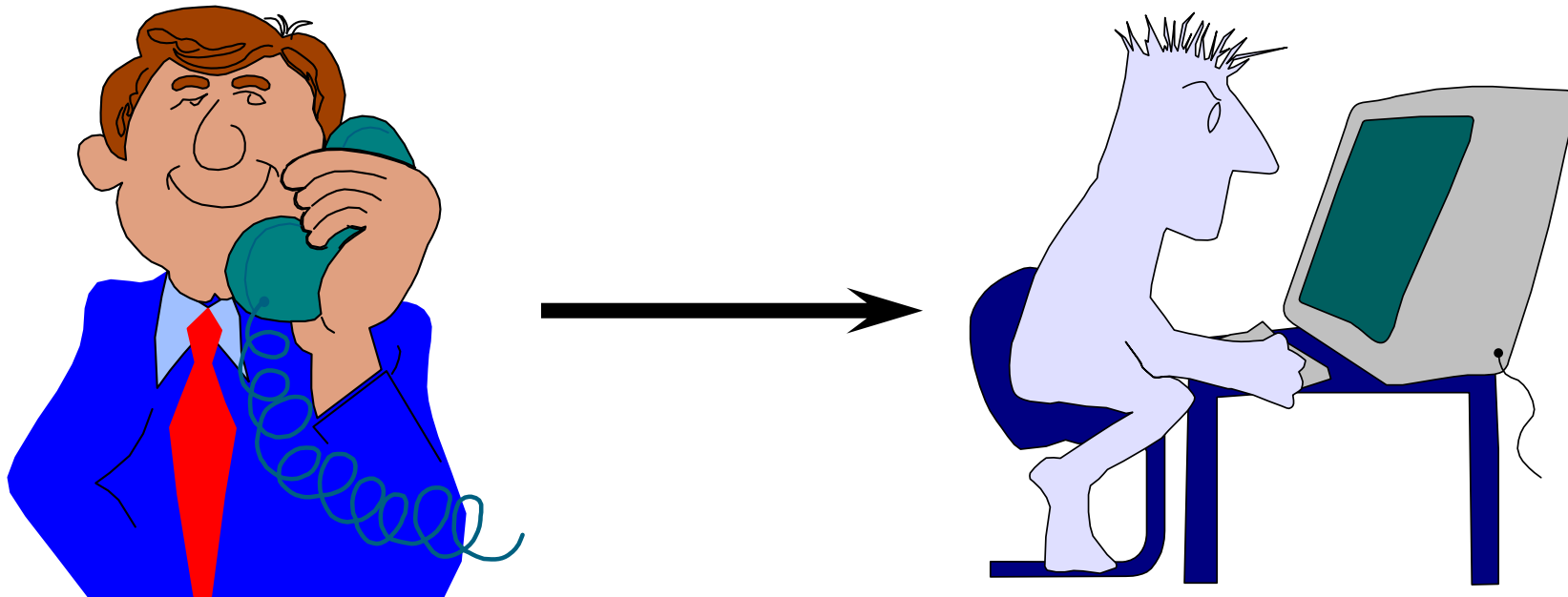
- ◆ What kind of Resource we have?
- ◆ How much you need and how often and how regular you need?
- ◆ How often you will initial request?
- ◆ How much traffic you could afford?
- ◆ How much “Promise” you could provide?
- ◆ How fair you are going to be?
- ◆ Control or “Let it be”?
- ◆ Power Saving Issues?
- ◆ Complexity?



Circuit Switch

- ◆ **Cellular System**
 - AMPS
 - GSM
- ◆ **Voice System**
 - Continue Traffic
- ◆ **Circuit Set up**
 - Reserve A trunk

HOW about Data



Packet Radio

- ◆ **Packet Nature**
 - If we could deliver information by packet
 - Bursty Type of Traffic
 - Packet Size

CSMA with Collision Detection/Avoidance

- ◆ **CSMA/CD: enhancement to slotted or unslotted CSMA schemes**
- ◆ **Node monitors its own transmission**
 - If collision detected, transmission is aborted without waiting for a NACK backoff and re-transmission procedure started
 - A jamming signal may be sent to get everybody else to abort too
- ◆ **Problem: does not work with RF wireless**
 - Cannot easily sense the channel while transmitting
 - ◆ MH's signal will dominate, need different receiving and transmitting antenna patterns
- ◆ **But, does work well with infrared wireless.. Directional receivers**
- ◆ **Wireless networks stick with ACK/NACK approach**
 - Popular called CSMA/CA
 - 802.11

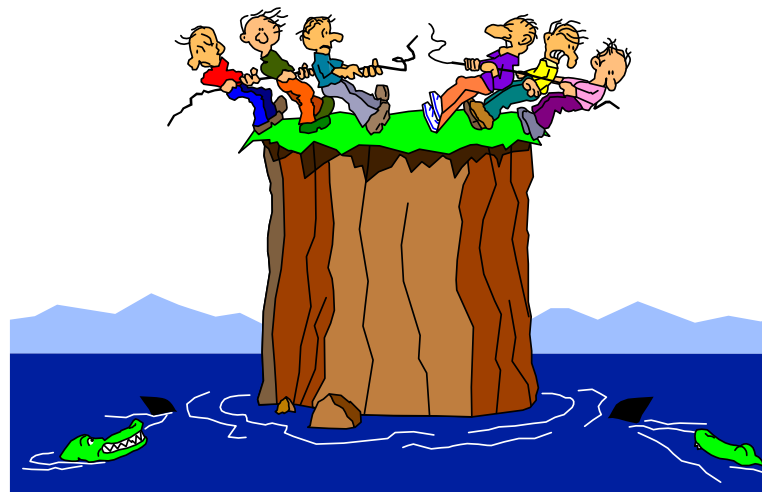
RANDOM Access

- ◆ Give everybody freedom



Hawaii Story

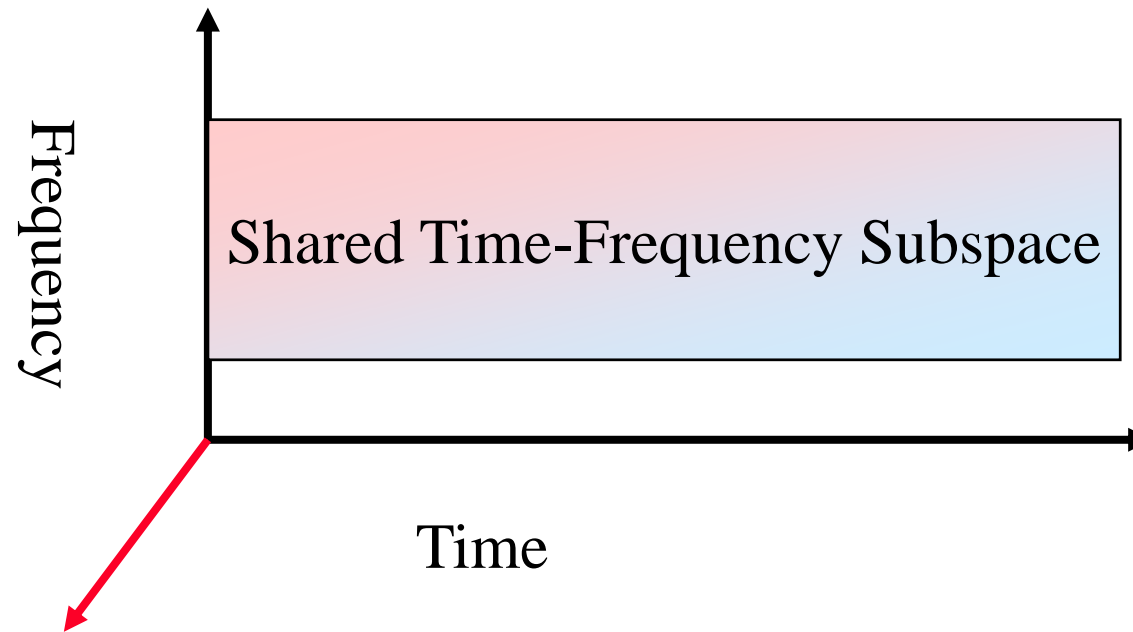
- ◆ University of Hawaii
- ◆ ALOHA
 - Hello and Goodbye



Multiple Access

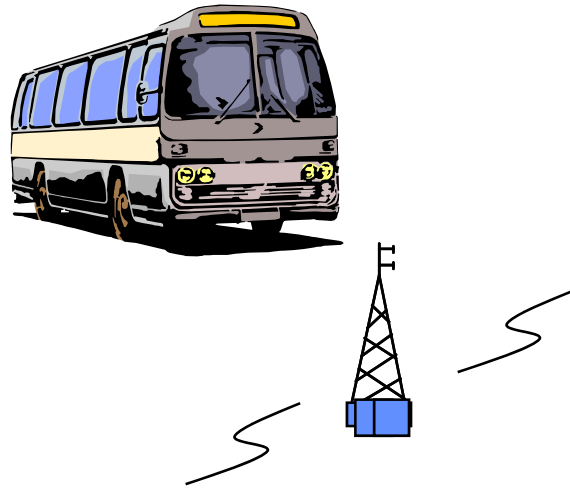
◆ Fundamental Problem

- How to share the Time-Frequency Space among multiple co-located transmitters?

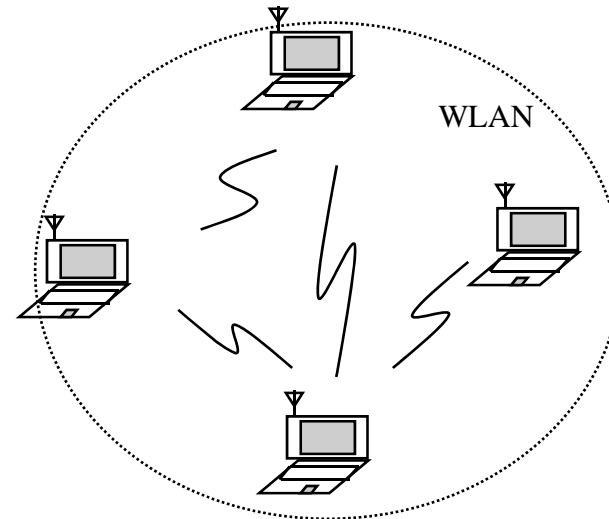


CDMA approach

Base-station versus Peer-to-Peer Models

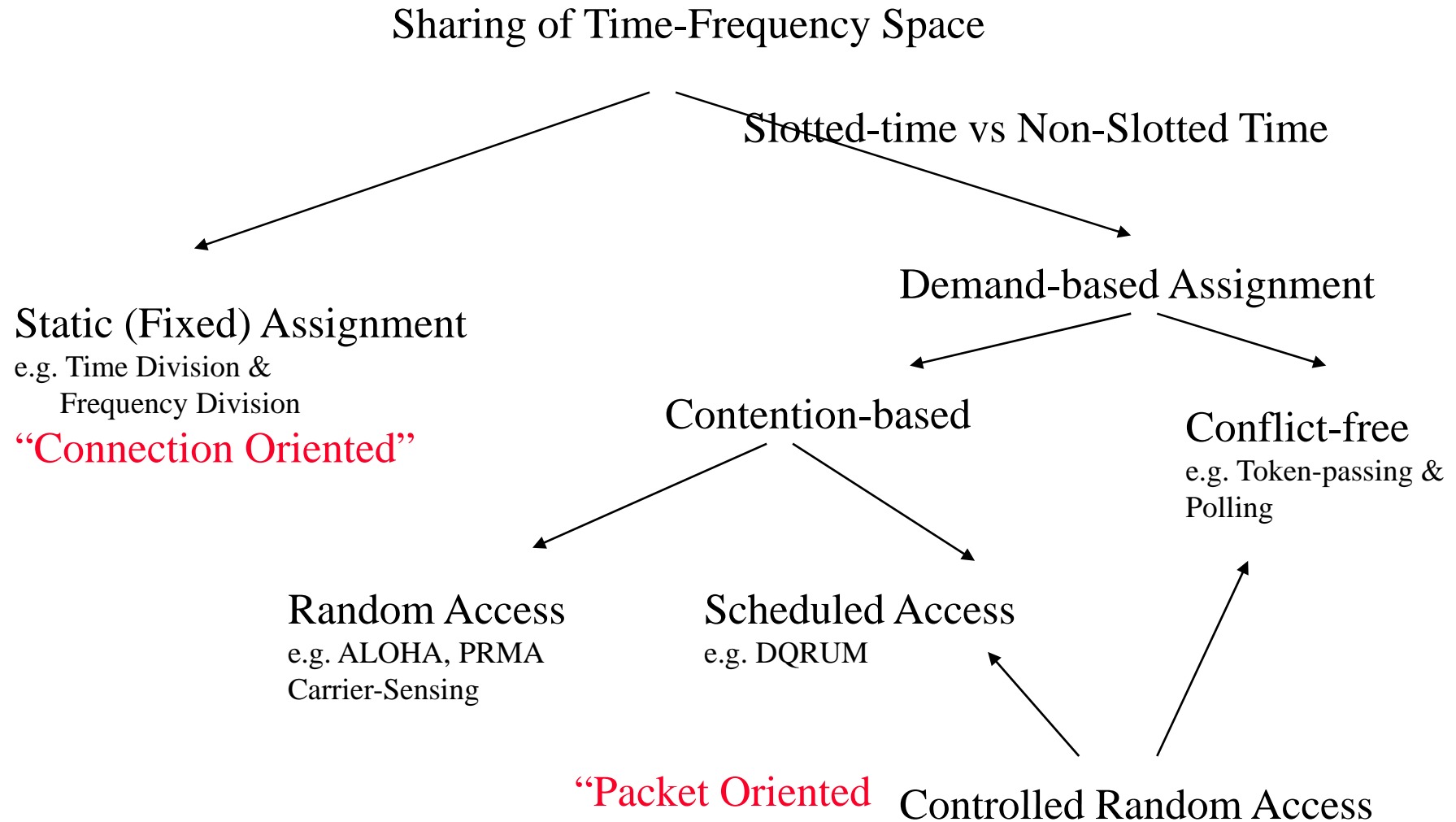


Base-station
(infrastructure-centralized)



Peer-to-Peer
(ad hoc network-
Fully-connected vs multihop)

Approaches to Wireless Multiple Access

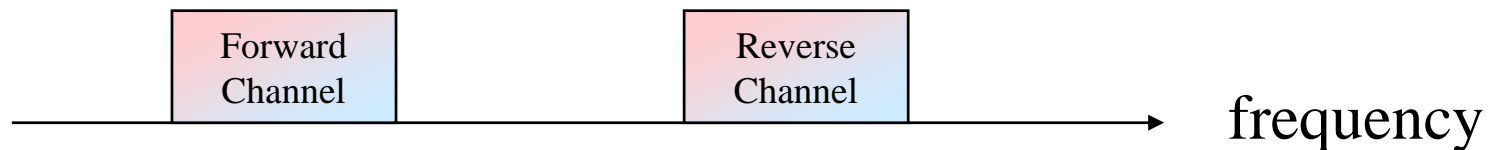


Frequency Division & Time Division Duplexing



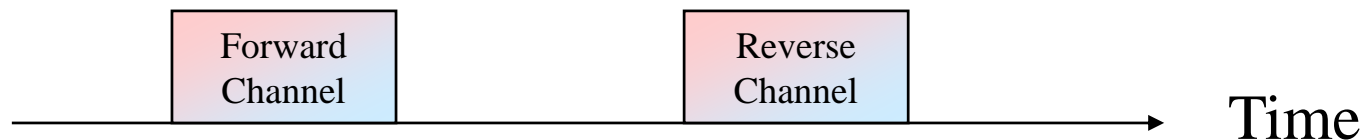
◆ Frequency Division Duplexing (FDD)

- Two distinct frequency at the same time for the two directions
- Frequency separation must be coordinated to allow cheap RF technology
- Coordination with out-of-band users between the two bands
- Geared towards providing individual frequencies for each user



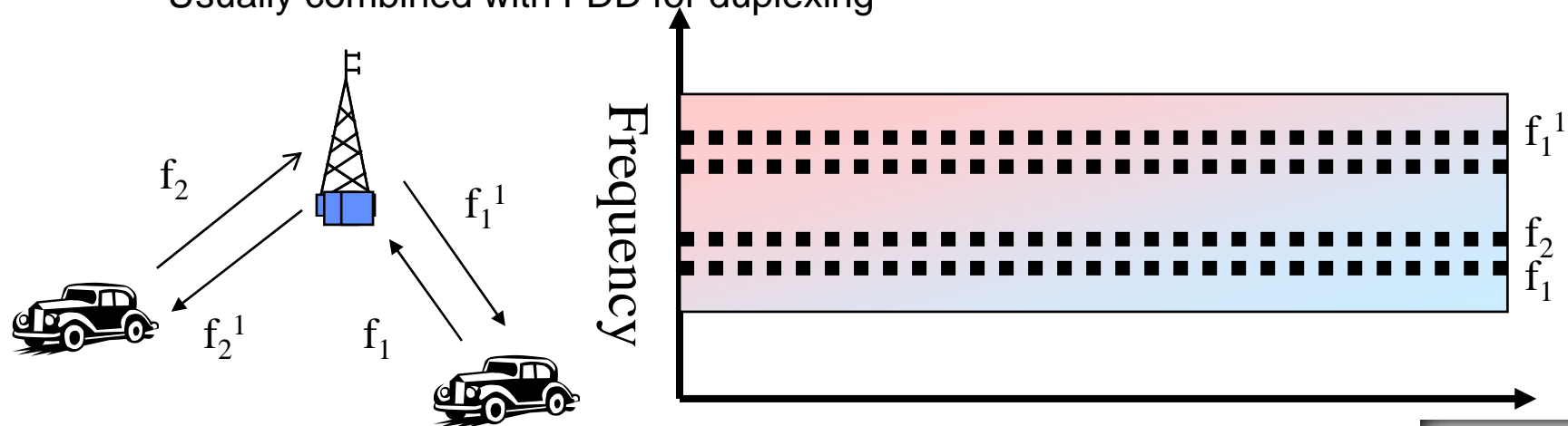
◆ Time Division Duplexing (TDD)

- Two distinct sets of time slots on the same frequency for the two directions
- Time latency because only quasi-duplex
- No need for RF duplexer



Frequency Division Multiple Access (FDMA)

- ◆ **Assign different frequency bands to individual users or circuits**
 - Frequency band (“channel”) assigned on demand to users who request service
 - No sharing of the frequency bands: idle if not used
 - Usually available spectrum divided into number of “narrowband” channels
 - ◆ Symbol time \gg average delay spread, little or no equalization required
 - Continuous transmission implies no framing or synchronization bits needed
 - Tight RF filtering to minimize adjacent band interference
 - Costly bandpass filters at basestation to eliminate spurious radiation
 - Usually combined with FDD for duplexing



Example-AMPS Cellular System

◆ User FDMA/FDD

- A channel is a pair of frequency duplexed simplex channels
- Each simple channel is 30 KHz
- Simple channels are separated by 45 MHz (allow cheap RF duplexers)
- Forward link 869-894 MHz, reverse link 824-849 MHz
- Two carriers per market share the channels

◆ Number of supported channels in AMPS

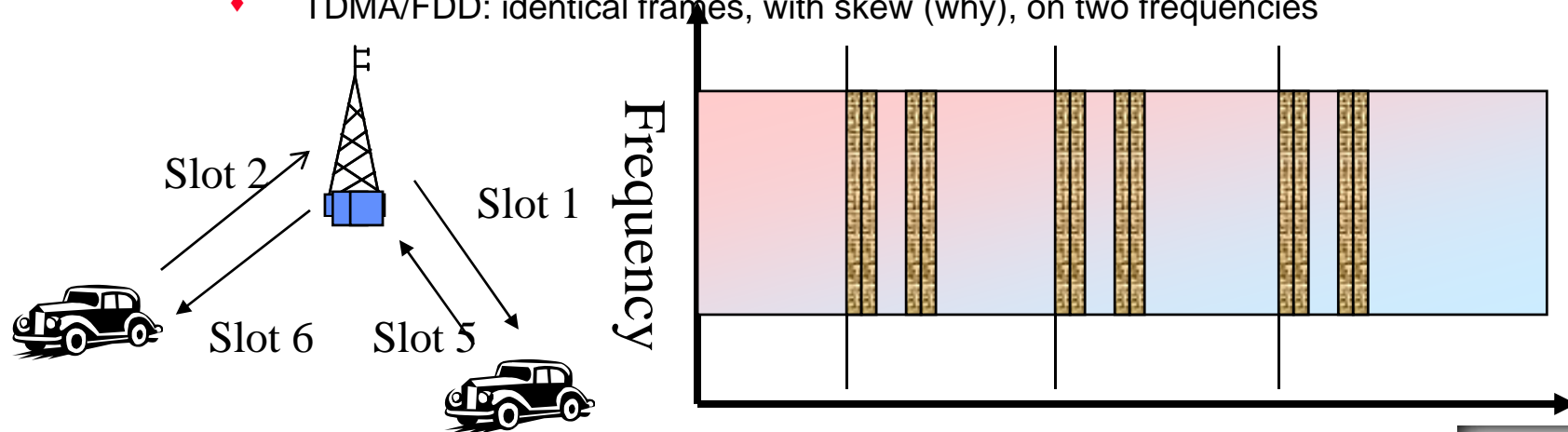
$$N = \frac{B_{total} - 2B_{guard}}{B_{channel}} = \frac{12.5MHz - 2(10kHz)}{30KHz} = 416$$

◆ Problem: set of active users is not fixed

- How is the FDMA/FDD allocated to a user who becomes active?
 - ◆ Static multiple access is not a complete solution .. Need a separate signalling channel with “demand-access”.
 - ◆ Pure FDMA is basically “dead” in the digital world

Time Division Multiple Access (TDMA)

- ◆ **Multiple user share frequency band via cyclically repeating “time slots”**
 - “channel”==particular time slot reoccurring every frame of N slots
 - Transmission for any user is non-continuous: buffer-and-burst digital data & modulation needed, lower battery consumption
 - Adaptive equalization is usually needed due to high symbol rate
 - Larger overhead-synchronization bits for each data burst, guard bits for variations in propagation delay and delay spread
 - Usually combined with either TDD or FDD for duplexing
 - ◆ TDMA/TDD: half the slots in a frame used for uplink, half downlink
 - ◆ TDMA/FDD: identical frames, with skew (why), on two frequencies



TDMA

- ◆ **More features**
 - Simply mobility & link control.. Snoop for other BSs during idle slots
 - Pulsating power envelop:interference with devices such as hearing aids

- ◆ **Possible enhancements to basic TDMA to integrate non-voice services**
 - Different # of slots per frame to different users (variable bit rate)
 - Dynamically reassign time slots for “bandwidth on demand”

Packet Radio

- ◆ **Packet Nature**
 - If we could deliver information by packet
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CSMA with Collision Detection/Avoidance

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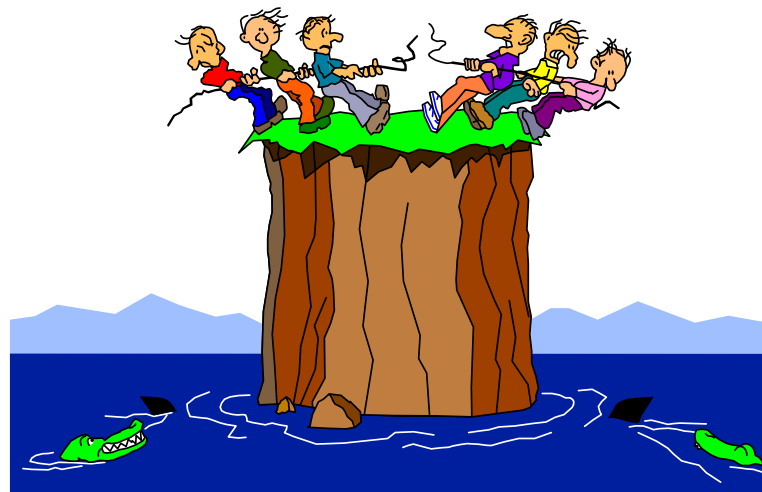
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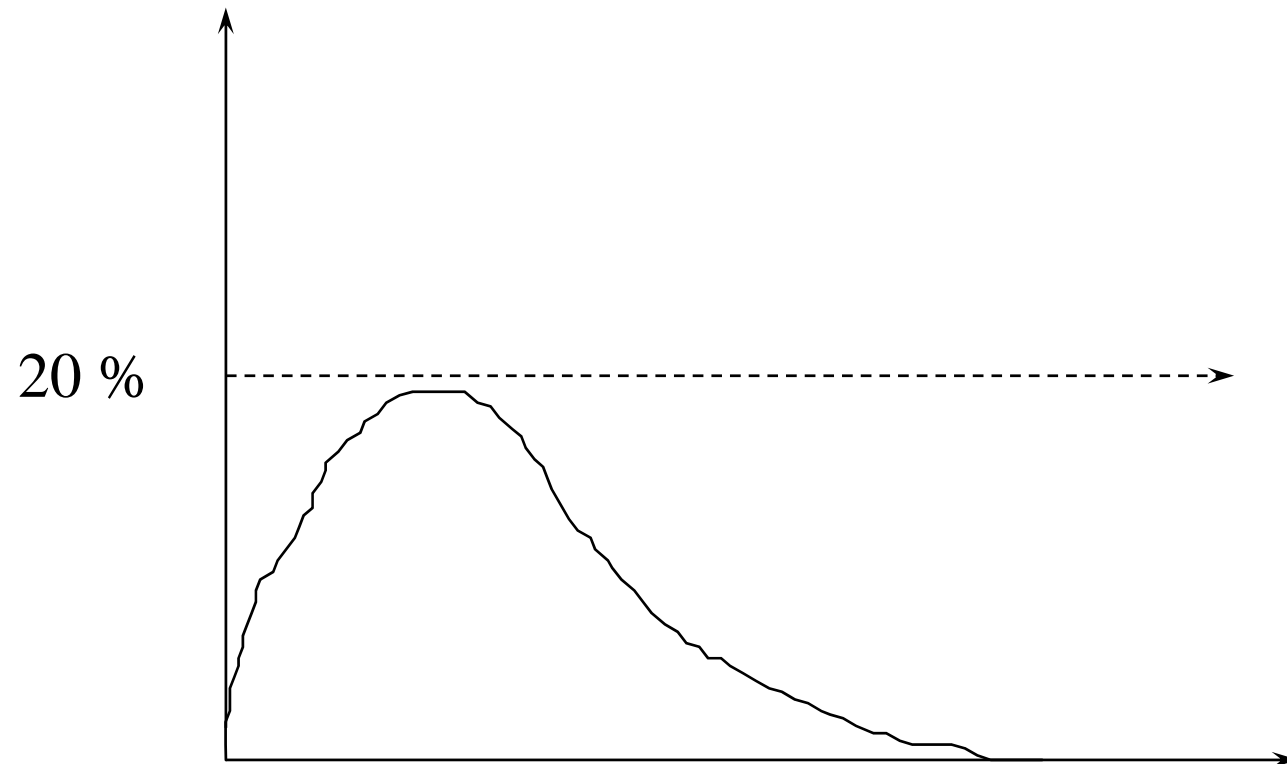


ALOHA System

- ◆ If you want, transmit
- ◆ If no acks
 - wait a random time
 - transmit the same packet again
- ◆ Problem ?
 - Collision ?
 - A lot of Users ?

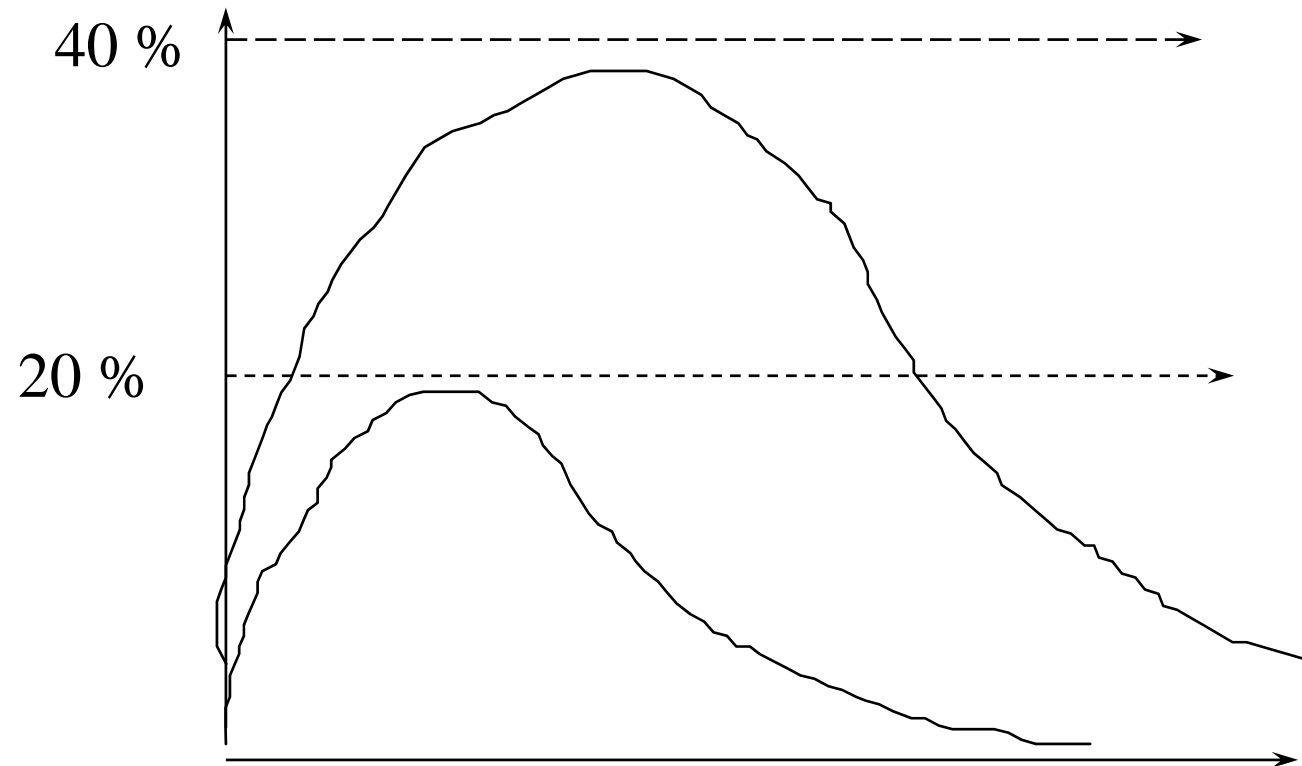


Pure ALOHA Throughput



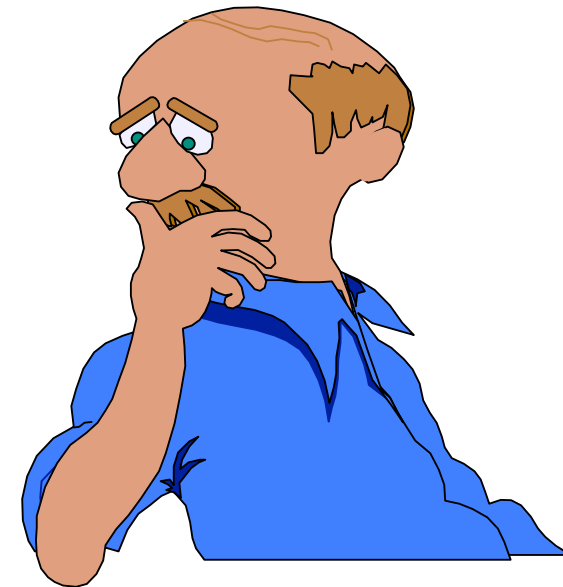
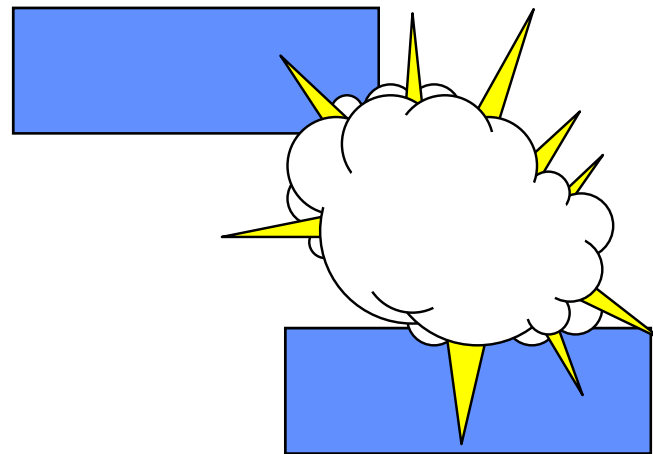
Traffic Load

Slotted ALOHA Throughput



Traffic Load

Slotted ALOHA



Maybe We could do
some arrangement ?

QoS & Delay



Traffic Load

DELAY

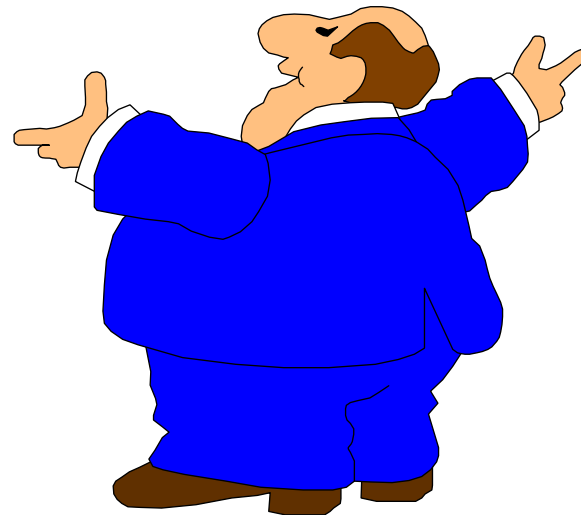
Whenever Users are many

- ◆ No one will succeed
- ◆ Collides all the time



Reason

- ◆ No one really listen to other people
- ◆ No one really cares



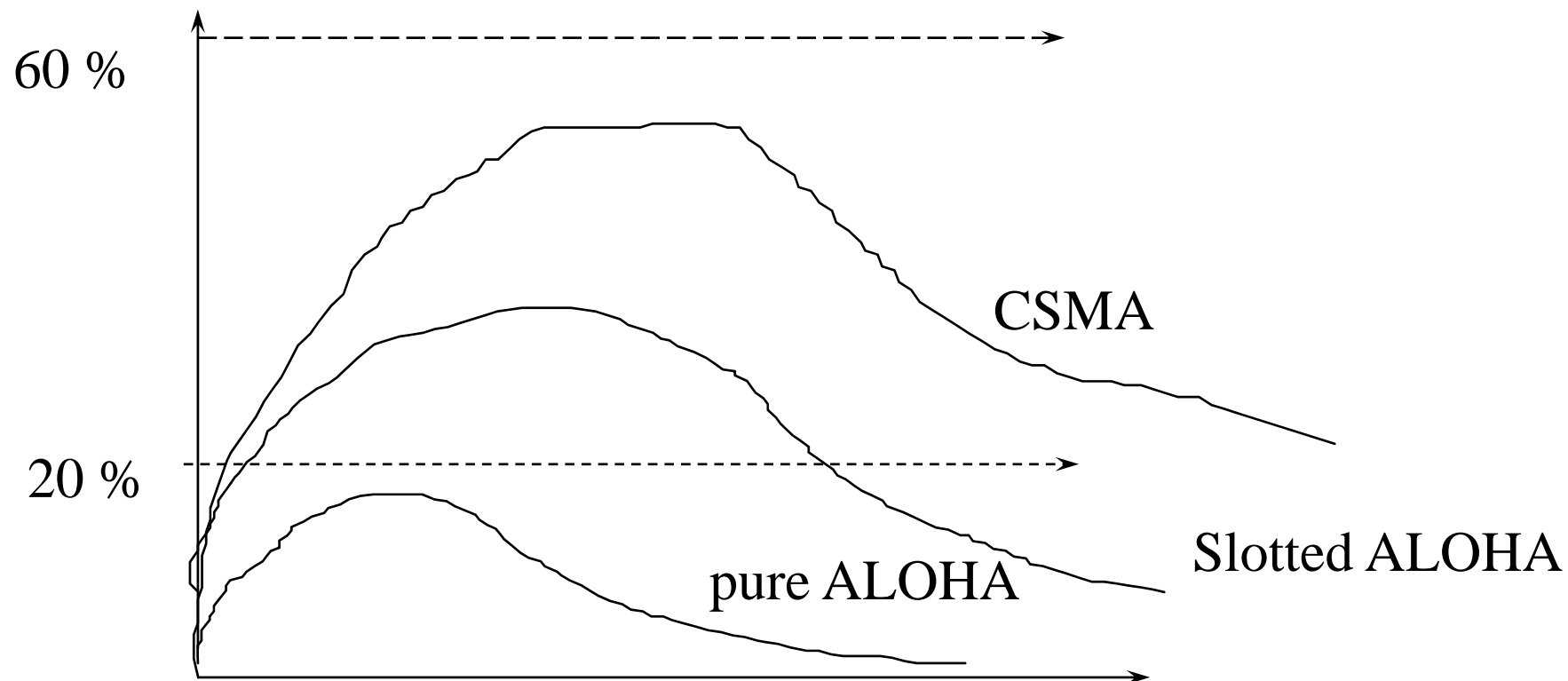
CSMA

- ◆ **Most LANs use CSMA**
- ◆ **Carrier Sense**
 - CSMA/CA: Collision Avoidance
 - CSMA/CD: Collision Detection

CSMA

- ◆ **Check if carrier is ok**
- ◆ **if the channel is free**
 - transmit
- ◆ **Otherwise, if the channel is busy**
 - wait a random time and try again
 - Back of a random time

CSMA



Traffic Load

Integrated CSMA/TDMA MAC Protocol

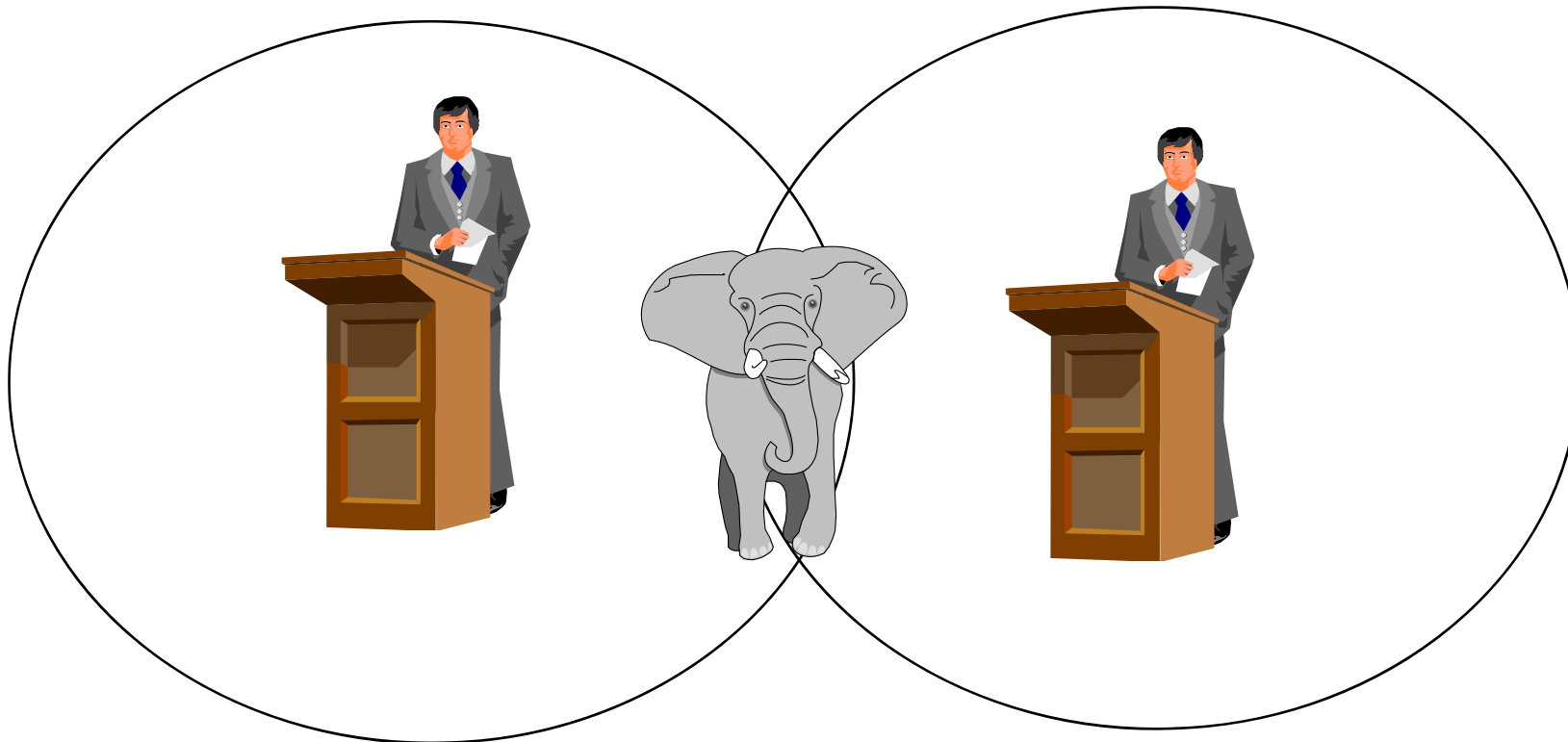
- ◆ Hybrid of reservation and Random Access
- ◆ A frame is segmented into:
 - Two reservation intervals for isochronous traffic
 - One interval for random access traffic

Can Support AP or Ad Hoc

- ◆ AP (Access Point)
- ◆ Ad HOC
 - Coordination Function will be distributed among all of the nodes of the ad hoc network

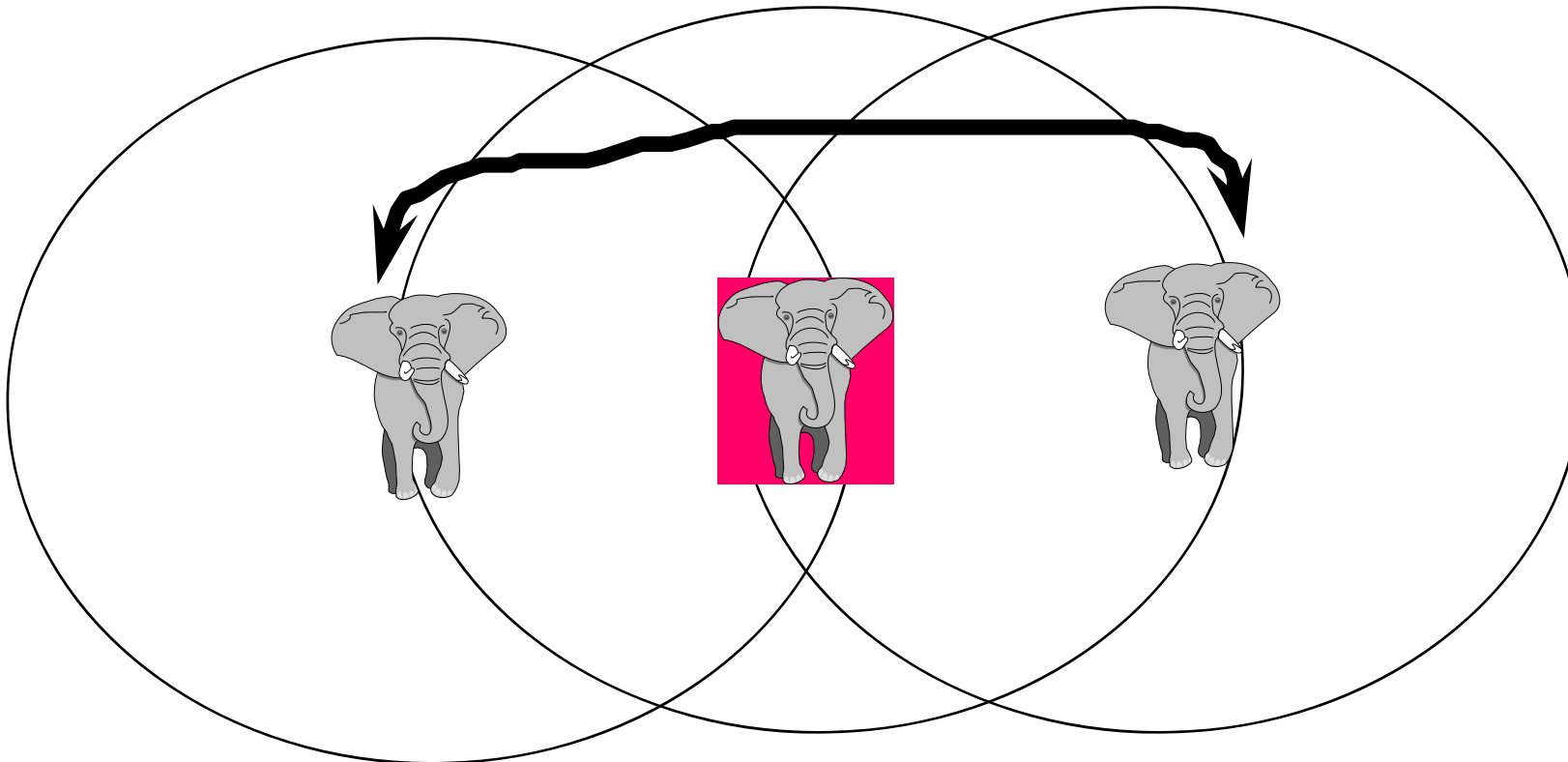
Challenge of Wireless Network

- ◆ Does “listen before you talk “ work ?



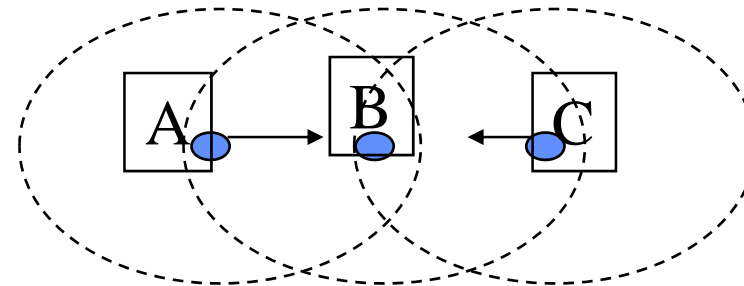
Hidden Terminal

- ◆ Due to transmission range

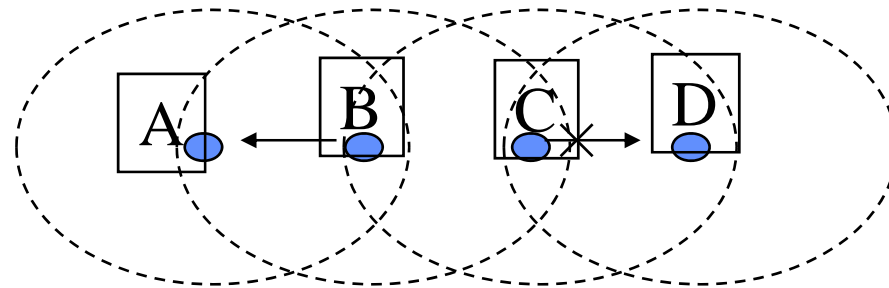


Carrier Sense Multiple Access (CSMA)

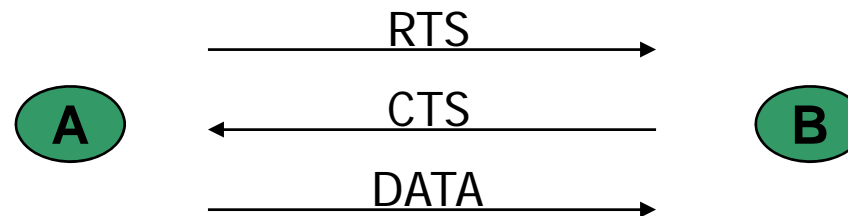
- ◆ To avoid collision, sender senses the carrier before transmission. But collision occurs at the receiver not transmitter.
- ◆ Hidden Terminal -



- ◆ Exposed Terminal-



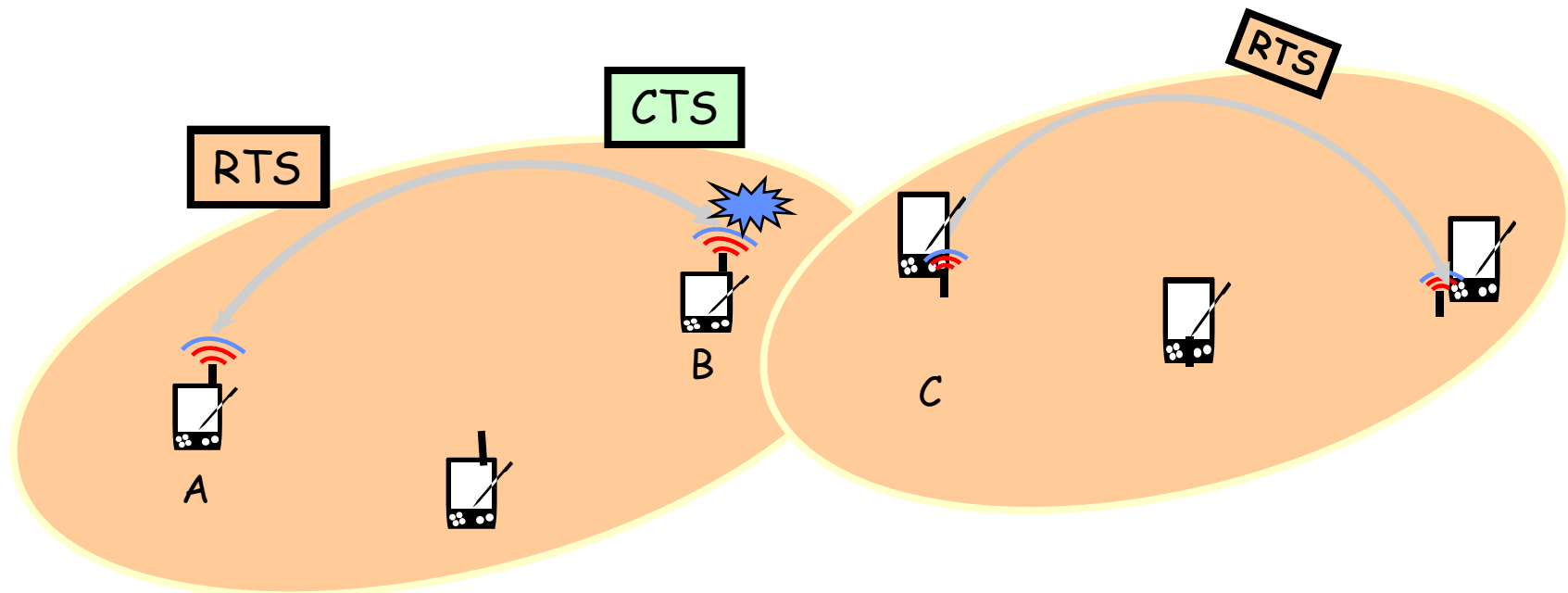
Multiple Access Collision Avoidance (MACA)



- ◆ Request-To-Send (RTS) packet: A to B.
- ◆ Clear-To-Send (CTS) packet: B to A.
- ◆ Node overhearing RTS will defer until A receive CTS.
- ◆ Node overhearing CTS will defer until B receive data.
- ◆ What do the above two features achieve (Hidden Terminal and Exposed Terminal)?

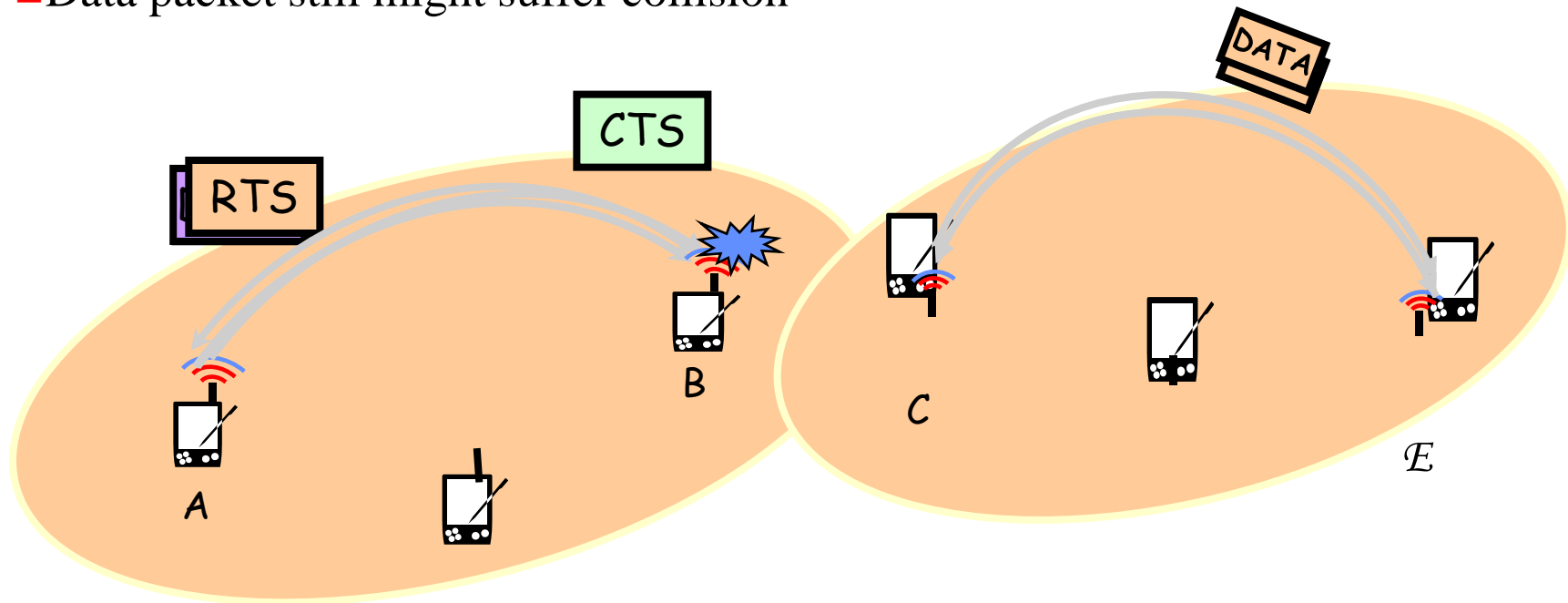
Hidden Terminal Problem Still Exists (1)

- Data packet still might suffer collision



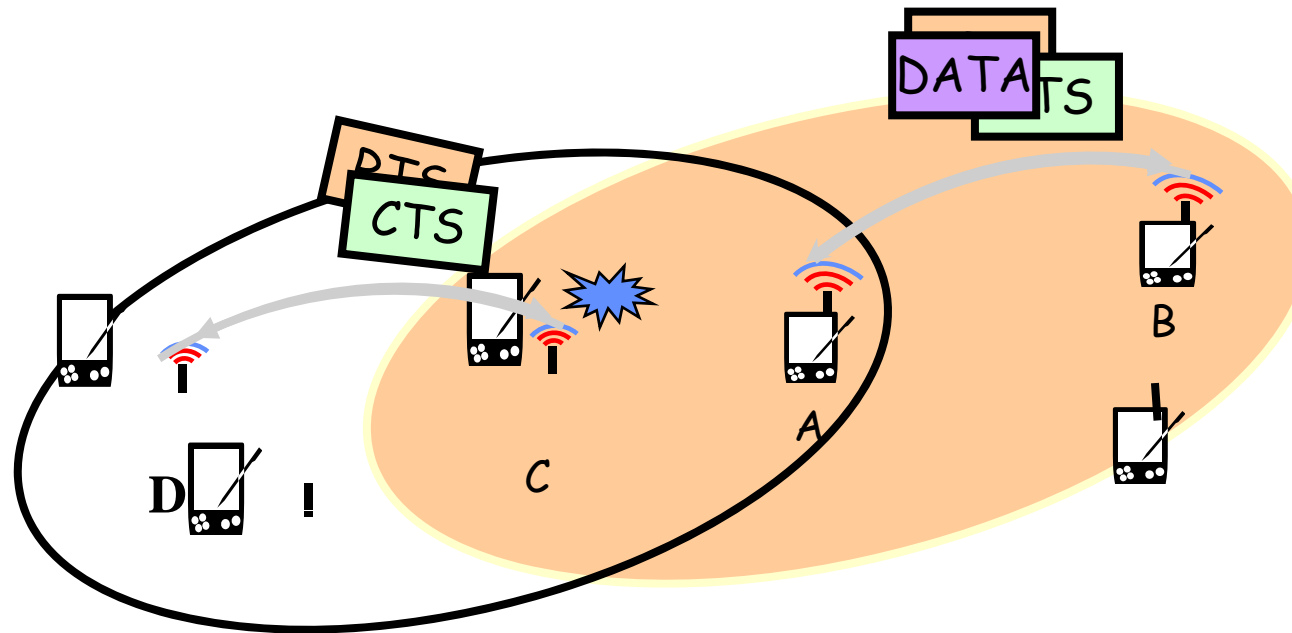
Hidden Terminal Problem Still Exists (2)

- Data packet still might suffer collision



Exposed Terminal Problem Still Exists

- Node C can not receive CTS



MACAW

Features

- ◆ **Backoff algorithm.**
- ◆ **Multiple Stream model.**
- ◆ **Basic Message Exchange**
 - ACK
 - DS
 - RRTS

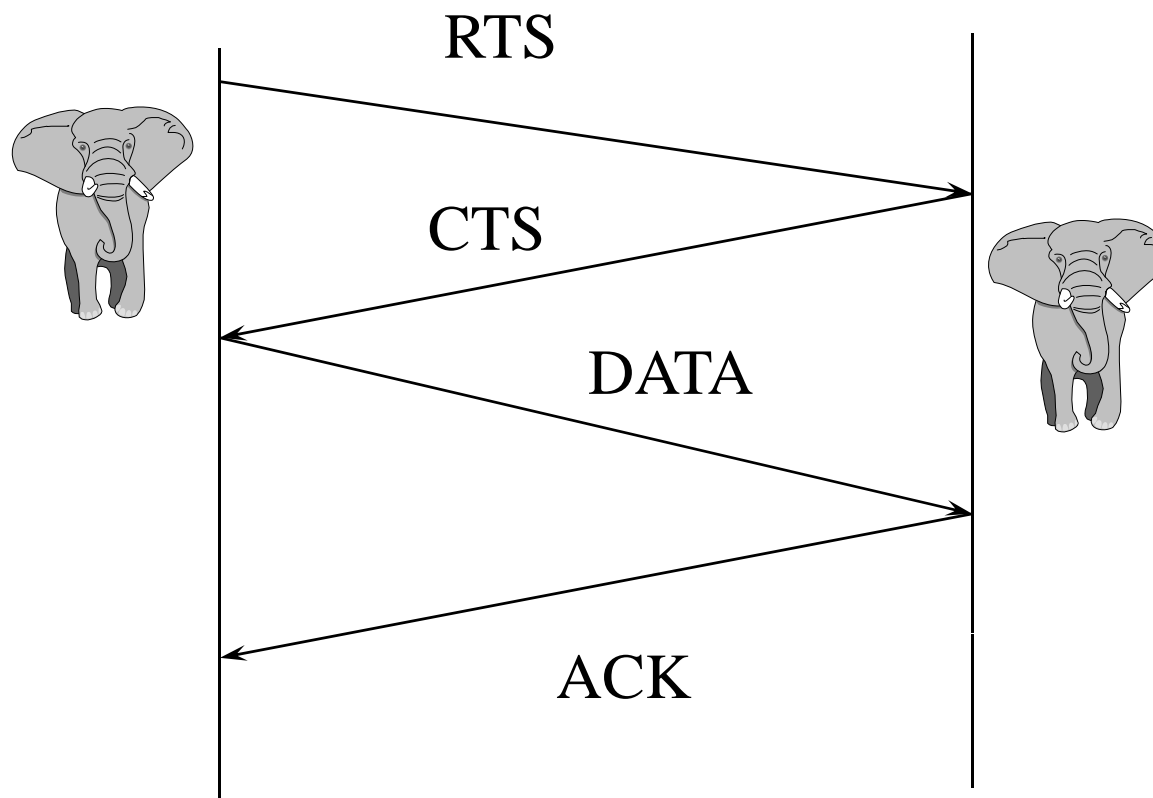
Backoff Algorithm

- ◆ **The algorithm used in MACA: Binary Exponential Backoff (BEB).**
 - Maintains a Backoff counter (BO)
 - BO is doubled after every collision
 - Reduced to minimal BO after every successful RTS-CTS exchange.
 - Sender waits for an interval chosen randomly between 1 and BO.
 - $F_{\text{inc}}(x) = \text{MIN} [2x, \text{BO}_{\text{max}}]$
 - $F_{\text{dec}}(x) = \text{BO}_{\text{min}}$
- ◆ **Results in unfair sharing of bandwidth.**

Modifications used in MACAW

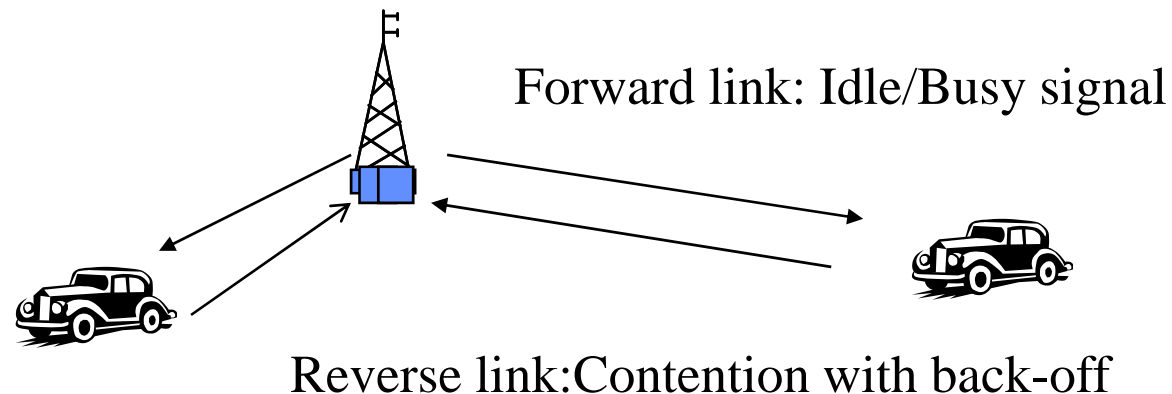
1. After every successful transmission all pads are made to have the same BO. (What is the problem with this?).
2. Gentler adjustment (MILD):
 - Upon collision $F_{inc}(x) = \text{MIN} [1.5x, BO_{max}]$.
 - Upon success $F_{dec}(x) = \text{MAX} [x-1, BO_{min}]$.

RTS/CTS/DATA/ACK



Data Sense Multiple Access (DSMA)

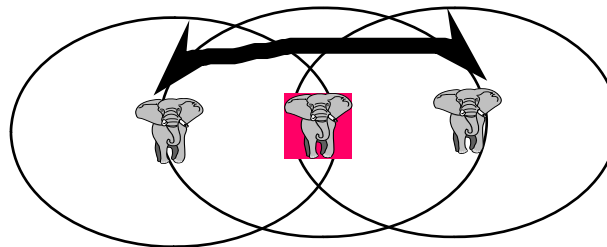
- ◆ Variation of CSMA-also called inhibit Sense Multiple Access
- ◆ Basestation transmits a busy/idle message on a forward control channel
- ◆ Mobile listens on the forward control channel for the busy/idle message
- ◆ Mobile transmits on the reverse channel only if busy/idle message indicates that the reverse channel is free
- ◆ Back-off and retransmit if collision occurs nevertheless
- ◆ Used in CDPD (Cellular digital packet data)



Problems in Contention-based Wireless Multiple Access



- ◆ **Near-Far effect-characterized by capture ratio of the receiver**
 - Strongest (near by) transmitter can capture the intended receiver
 - Weaker (far away) transmitters get ignored by the receiver
 - Depends on receiver and modulation used
 - Fairness terminal problem
- ◆ **Hidden terminal problem**
 - Terminal “hidden” from the transmitter may disrupt the receiver
 - Makes carrier sensing ineffective
 - A cannot detect collisions at B due to transmission from C
 - Solve by using RTS/CTS control frame to reserve medium



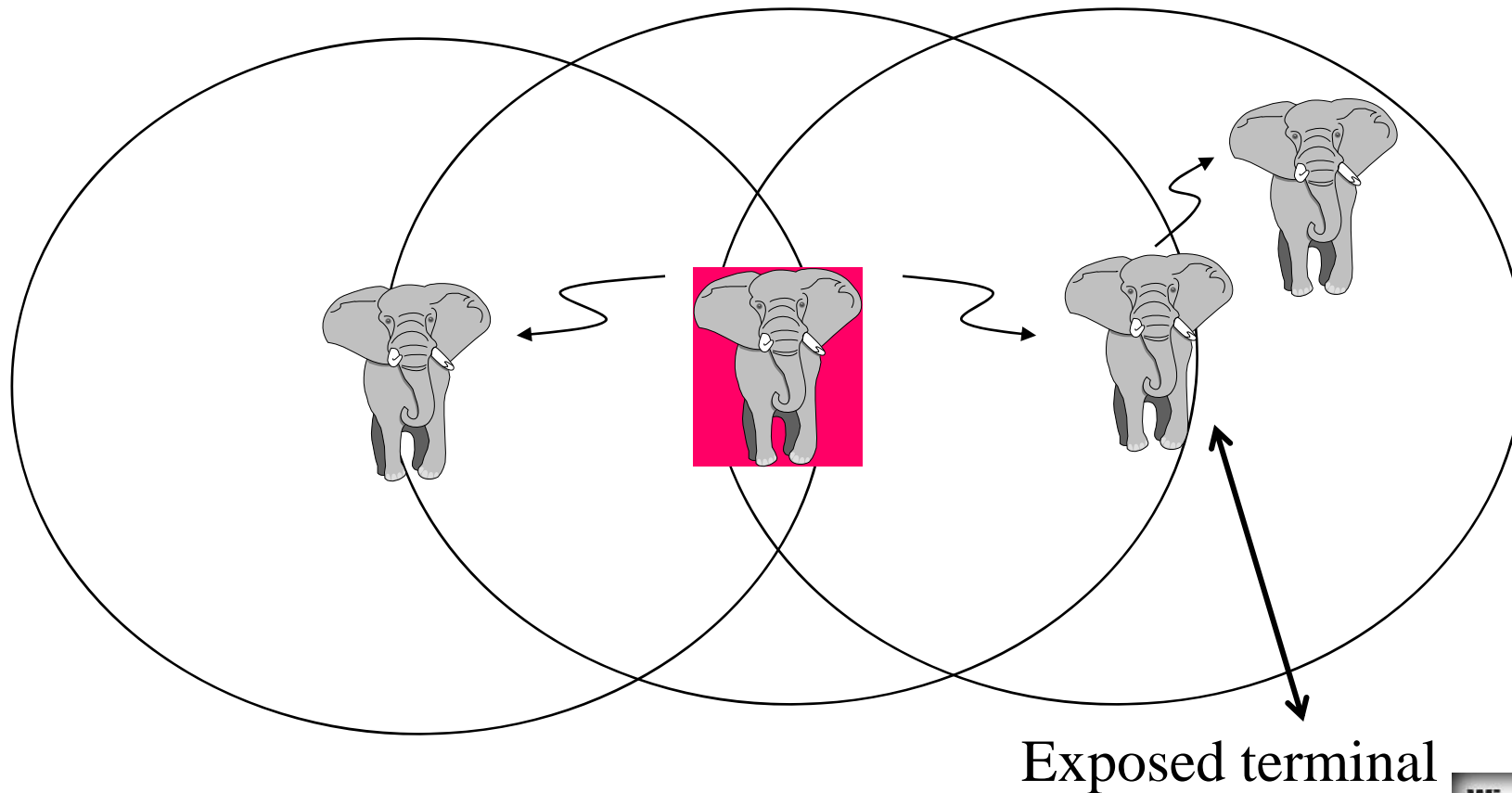
More on RTS/CTS

- ◆ **RTS/CTS serve to “reserve” the medium**
 - RTS contains length of proposed transmission
 - CTS also contains length of proposed transmission
 - MHs overhearing RTS defer all transmissions until after CTS would have finished (including receiver turnaround time)
 - MHs overhearing CTS defer for length of data packet transmission
 - Retransmission happen only if no CTS is received in reponse to RTS

- ◆ **Binary exponential backoff (BEB) has problems**
 - Does not provide fairness if every MH generate enough traffic to consume the channel
 - After collisions, the less-backed-off mobile wins eventually all but one MD are backed-off to BOMax

Exposed Terminal Problem

- ◆ **C will sense channel busy, and defer, but doesn't need to**
 - The C to D transmission can take place but is delayed



CSMA/CD?

- ◆ Collision Detection ?
- ◆ If a collision is detected, stop transmitting the present packet ?
- ◆ Is CSMA/CD possible ?
 - transmit and receive at the same time ?
 - CSMA wireless network, transmit and receive at the same frequency band
 - unlike Cellular System, uplink and downlink

IEEE 802.11 MAC

- ◆ **Support for multiple access PHYs; ISM band DSSS and FHSS, IR @ 1 and 2 Mbps**
- ◆ **Efficient medium sharing without overlap restrictions**
 - Multiple networks in the same area and channel space
 - Distributed Coordination Function: using CSMA/CA
 - Based on carrier sense mechanism
- ◆ **Robust against interference (e.g. co-channel interference)**
 - CSMA/CA+ACK for unicast frame with MAC level retransmission
- ◆ **Protection against Hidden terminal problem: Virtual Carrier Sense**
 - Via parameterized use of RTS/CTS with duration information
- ◆ **Provision for Time Bounded Services via Point Coordination Points**
- ◆ **Configurations: ad hoc & distributed system connecting access points**
- ◆ **Mobile-controlled hand-offs with registration at new basestation**

Schedule Access-Reservation-based Protocols



- ◆ Also called “Demand Assigned Multiple Access”
- ◆ Center agent that acts a slot scheduler
- ◆ Sender request “reservations” for future time slots
- ◆ Central agent assigns a slot
- ◆ Data transmission in the assigned slot is done without contention
- ◆ Assumption is that data packets \gg reservation request packets
- ◆ Overhead of reservation and acknowledgement messages
- ◆ Trades higher throughput (up to 80% utilization) for higher latency

Order MAC Techniques

- ◆ **Token Bus and Token Ring**
 - Token are passed among nodes
 - How about wireless network ?
 - ◆ Nodes might leave ?
 - ◆ Break the Order
 - ◆ Take away the token

Basic Scenario

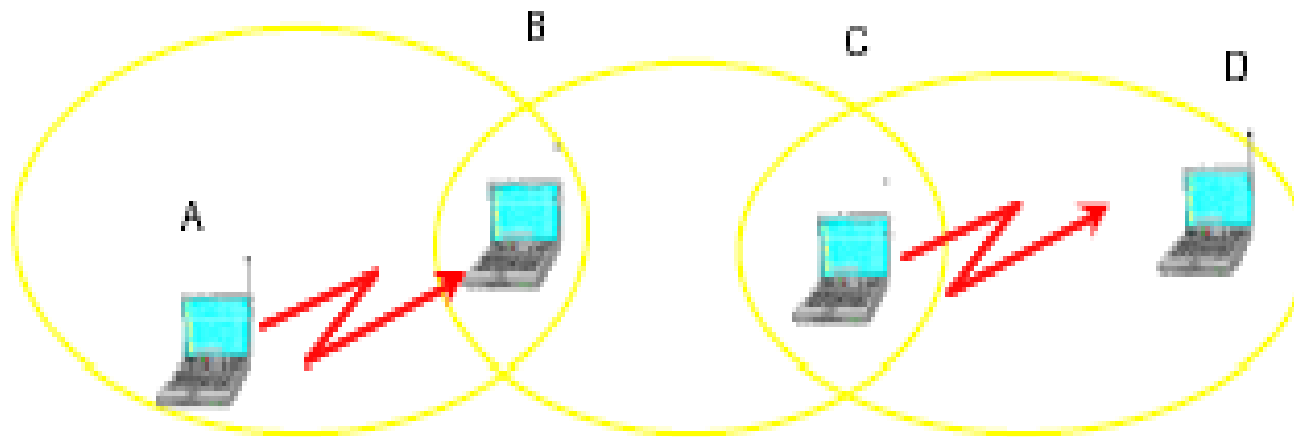


Fig. 1. A is sending a packet to B when C should decide whether to transmit to D.

Hidden and Exposed Stations

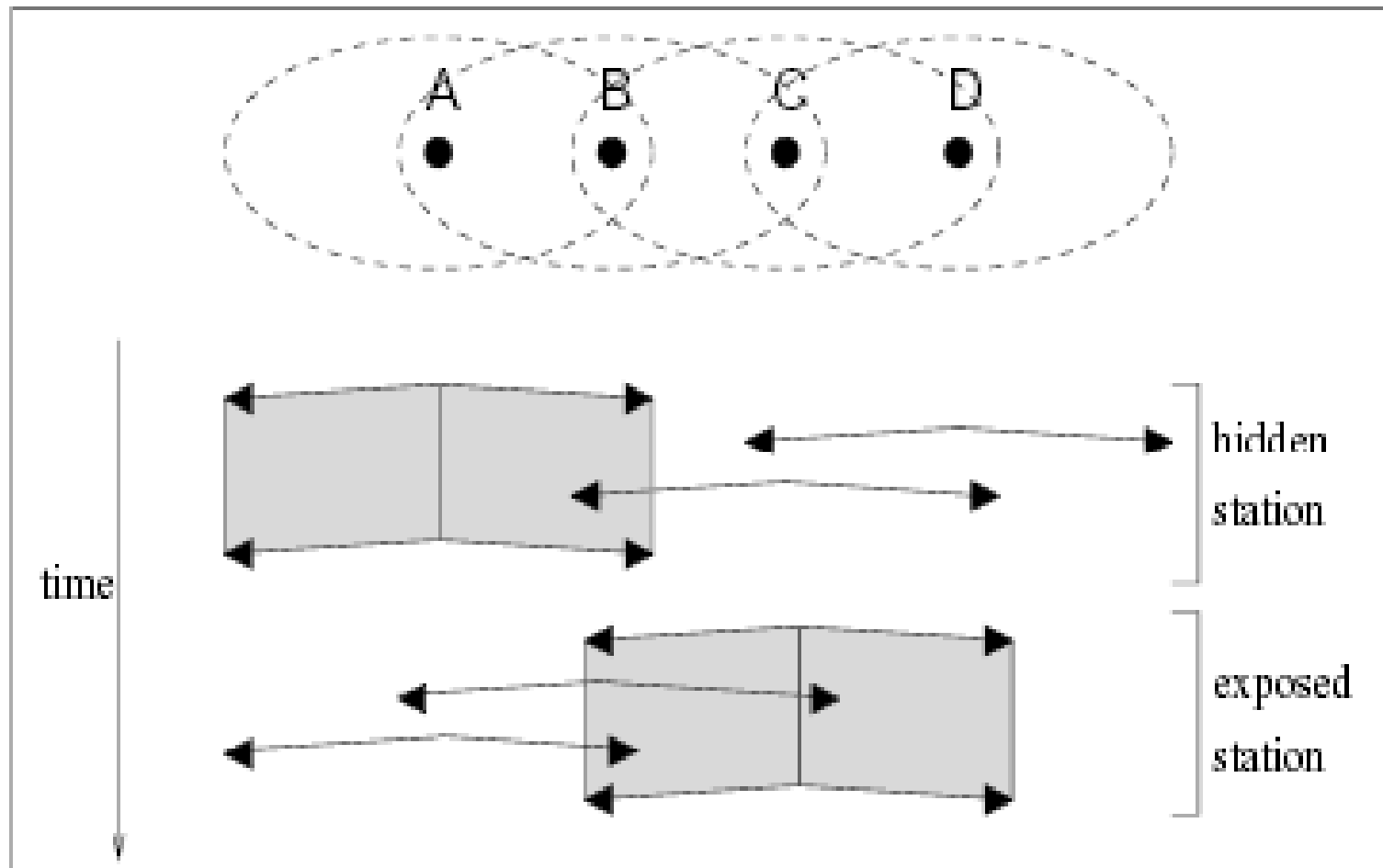


Figure 1: Hidden and Exposed Stations

Capture Effect/Near Far Problem

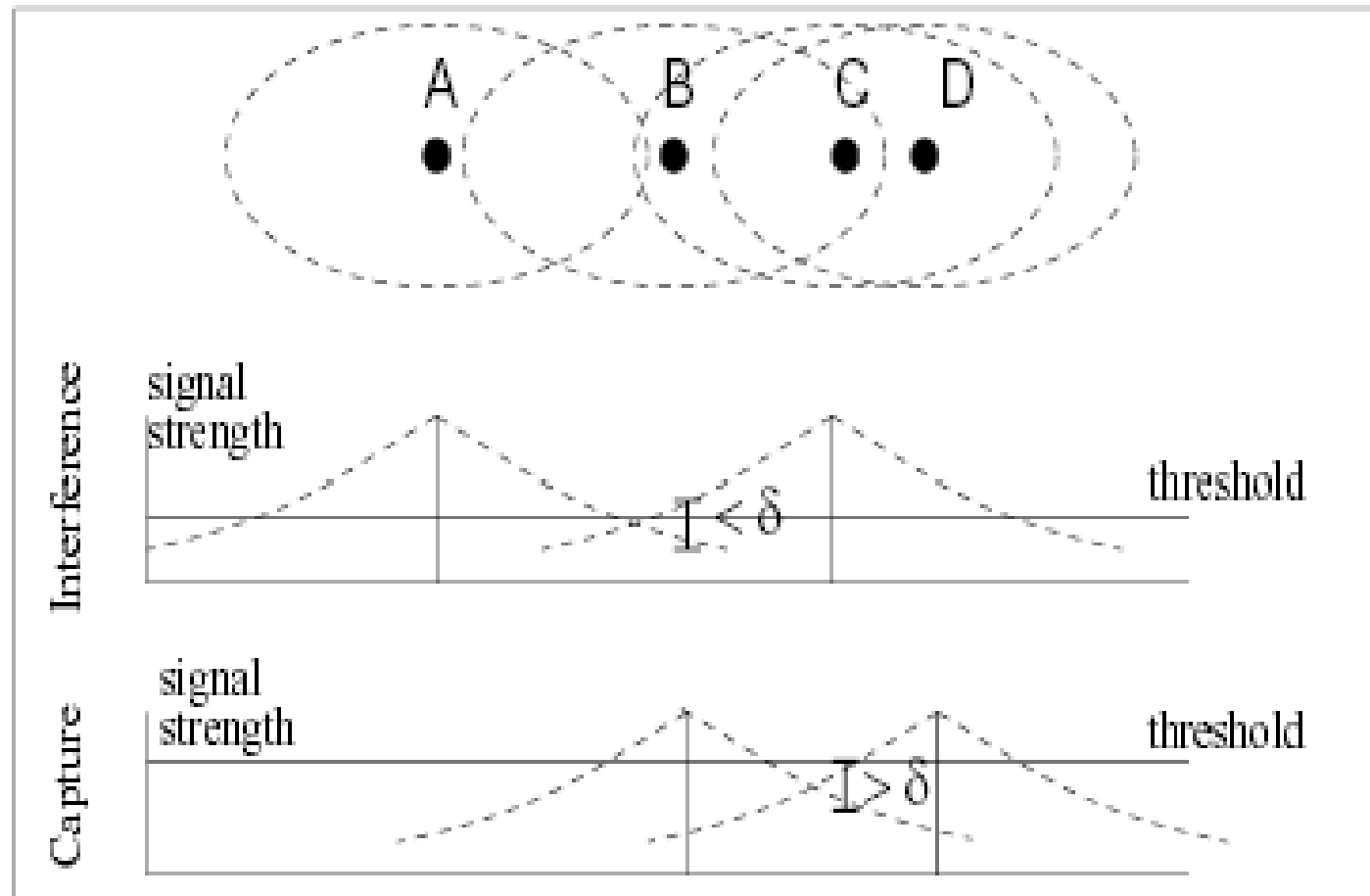
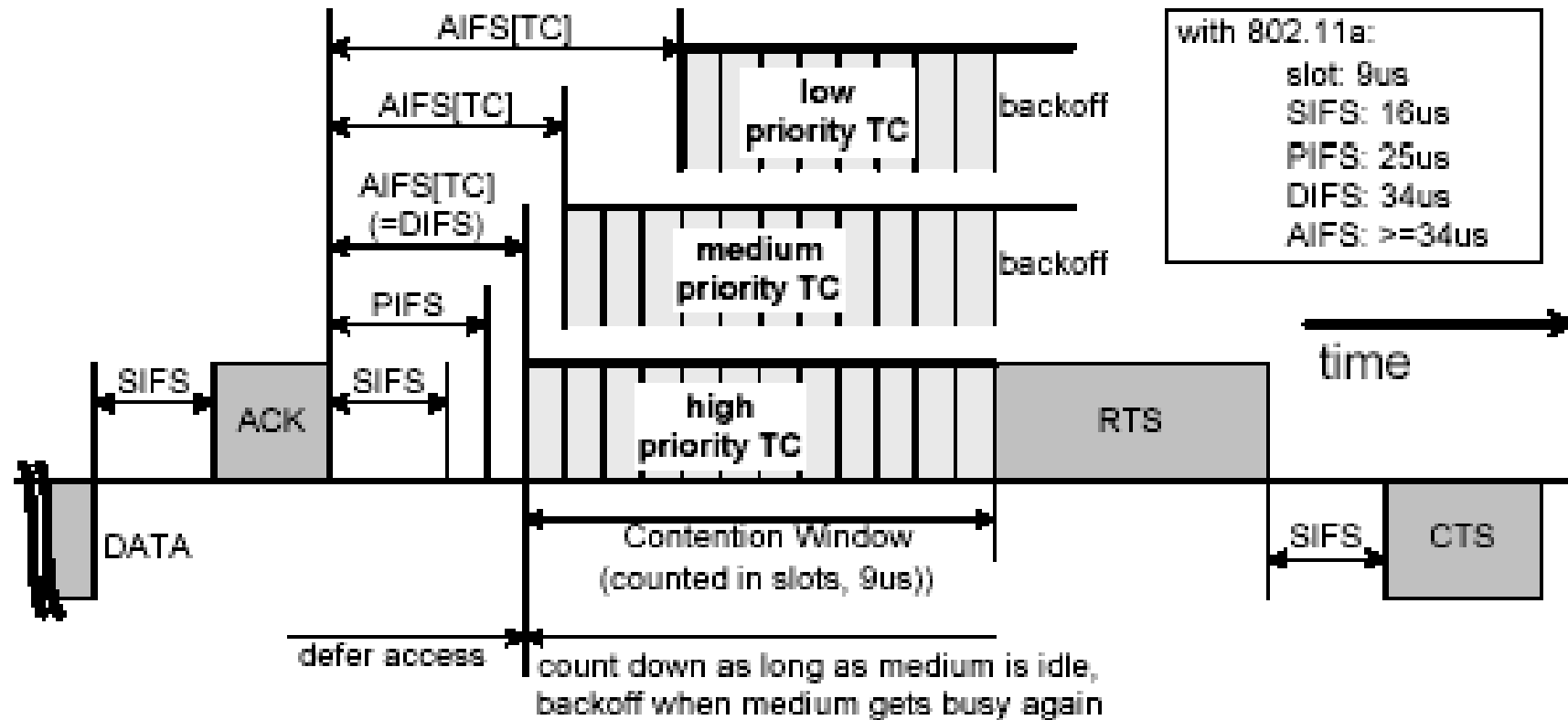
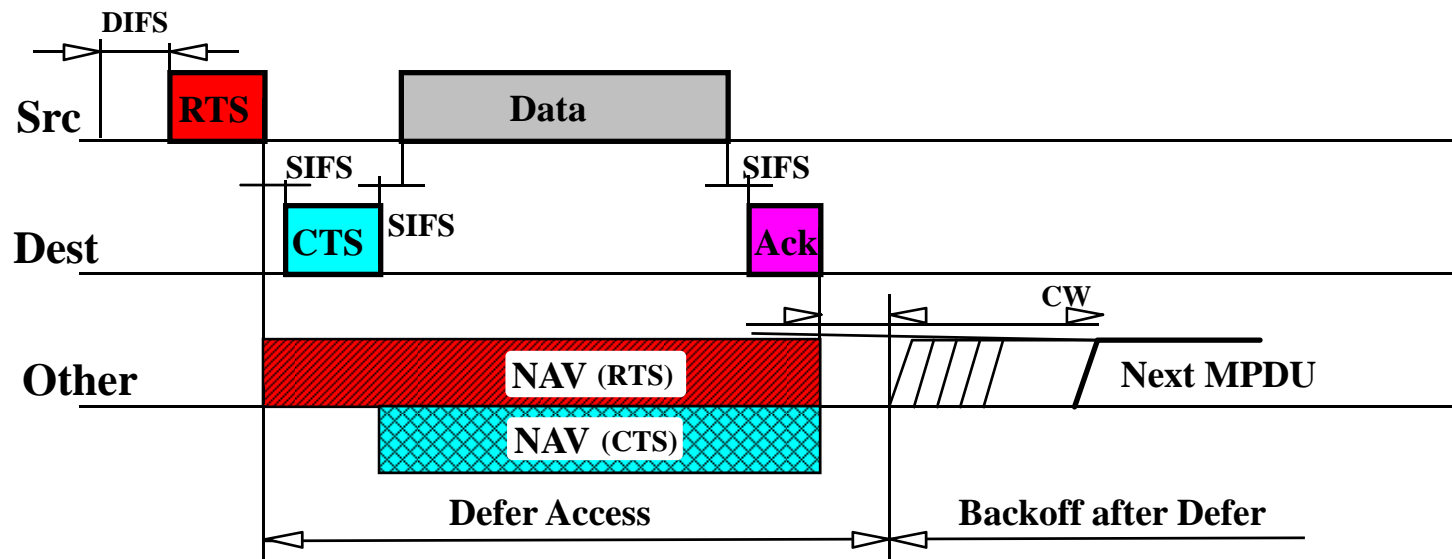


Figure 2: Interference and Capture

802.11 E



802.11



Interference Issue for CSMA/CA

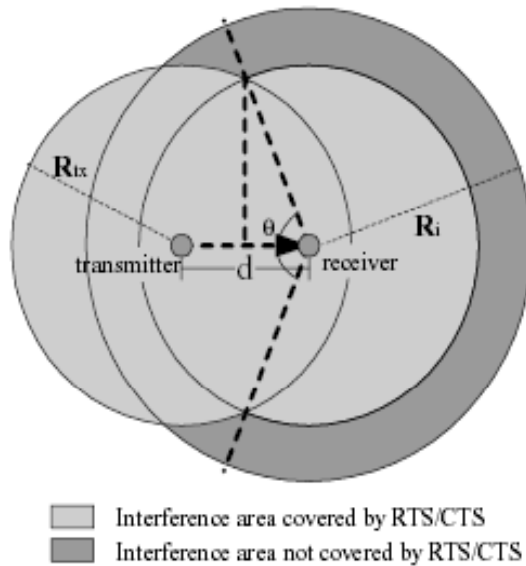


Fig. 1. Effectiveness of RTS/CTS handshake when d is larger than $T_{SNR}^{-\frac{1}{k}} * R_{tx}$ and smaller than R_{tx} .

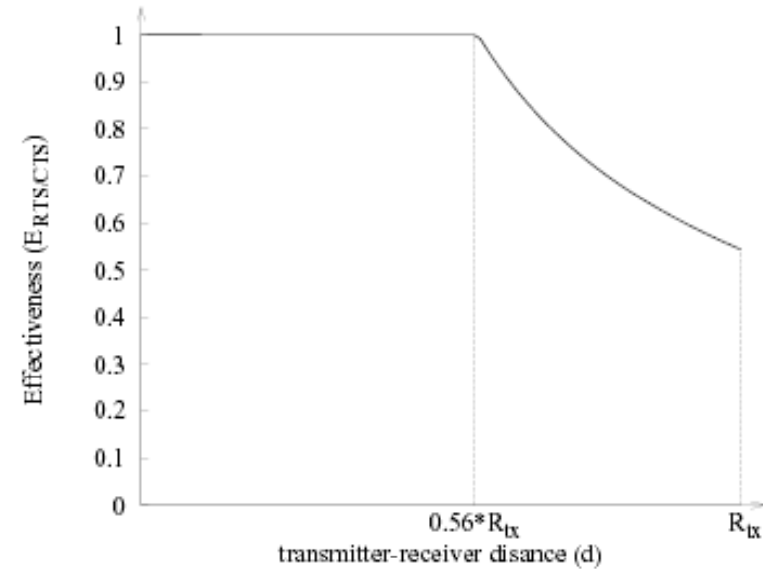
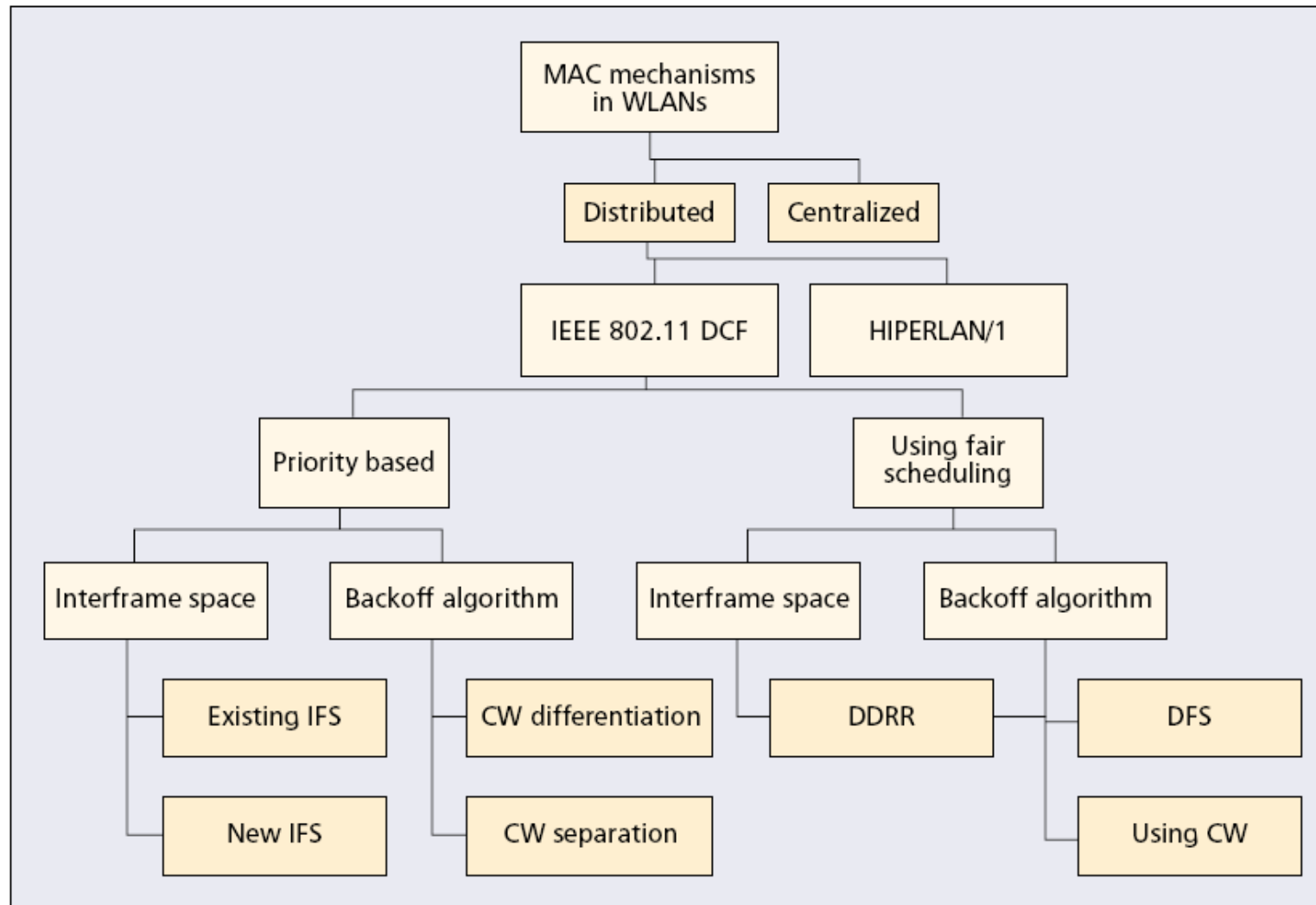


Fig. 2. Effectiveness of RTS/CTS handshake for TWO-RAY GROUND model and SNR threshold as 10.

QoS issue for 802.11

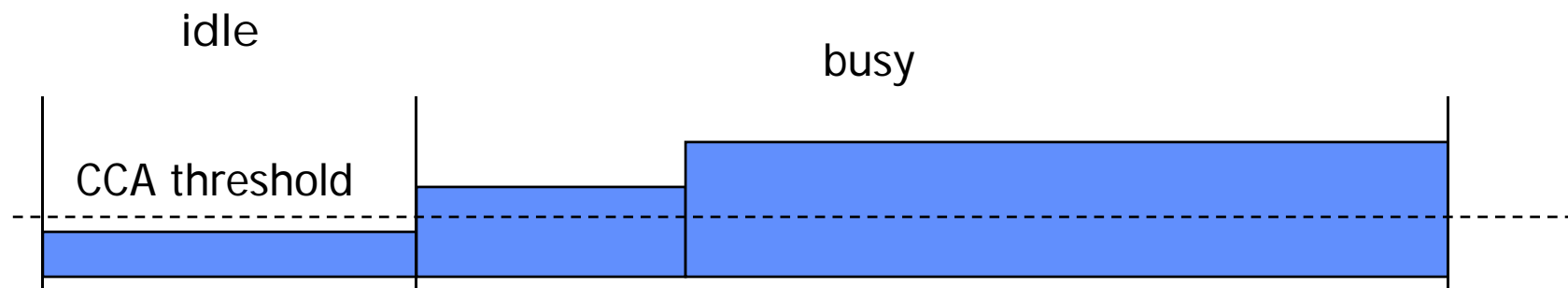


High-Density (HD) WLAN

- ◆ In HD-WLAN, its overall capacity can be expressed as $L \times S \times C$
 - L – per link capacity
 - C – number of simultaneous trans. Per channel.
 - S – the number of non-interfering channels
- ◆ Hence, the issues of HD-WLAN is
 - How to increase the performance of S.
- ◆ Co-Channel Inference (CCI)

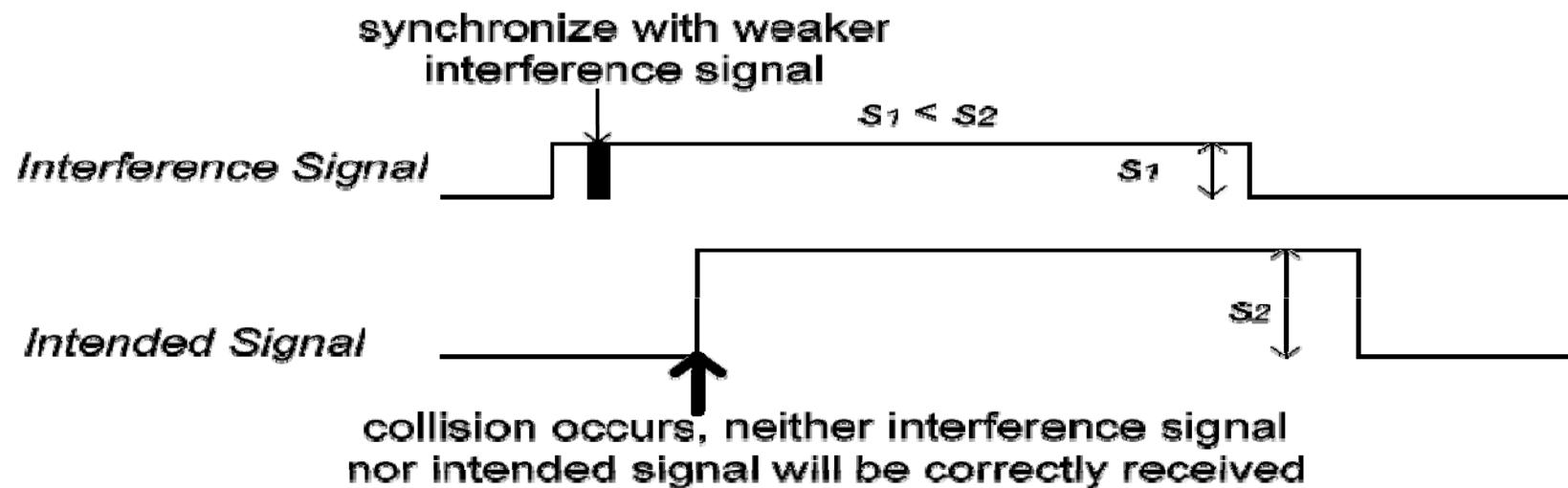
Clear Channel Assessment (CCA)

- ◆ A station performs CCA before a data trans. to sample the energy in the channel.
- ◆ The station will proceed only if the sampled energy is below a threshold known as the **CCA threshold**.



Receiving Sensitivity (RS)

- ◆ Today's consumer 802.11 radios are often not able to preempt a receiving process to capture a newly-arrived strong signal.
- ◆ This issue called “stronger-last” collision”.



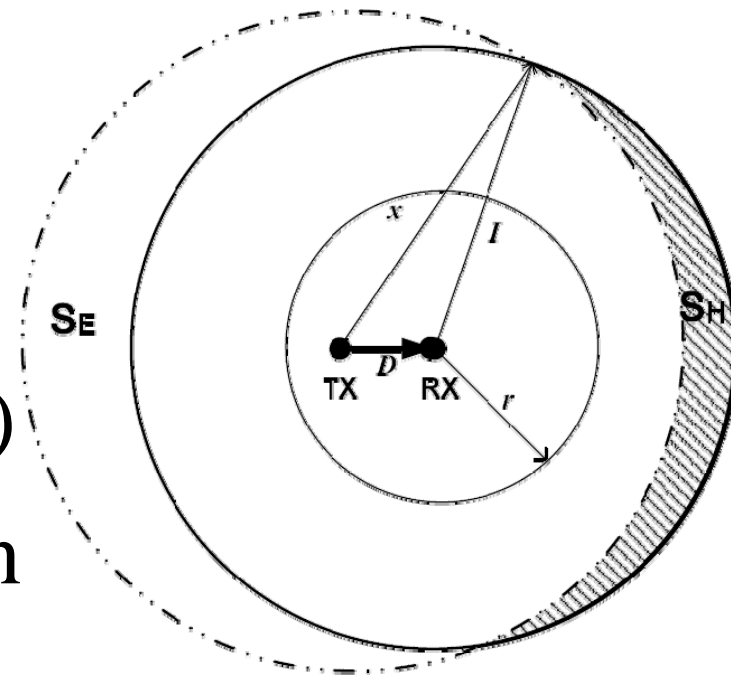
Analytical Model for RS/CCA Adapt.

- ◆ In 802.11 WLAN research, the logarithm path loss model is widely used to show **average** SS at receiver.

$$P_{RX}(d) = P_{RX}(\bar{d}) \left(\frac{d}{\bar{d}} \right)^{\gamma}$$

$\gamma = 2$ free-space (LOS)

$\gamma = 4$ ground reflection



Only Strong signals triggers Recv.

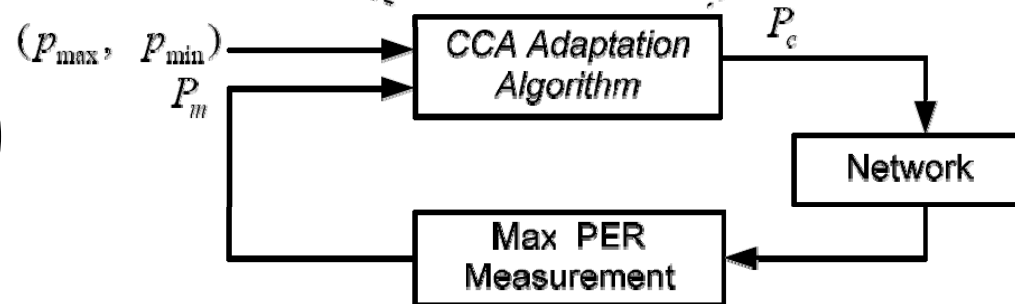
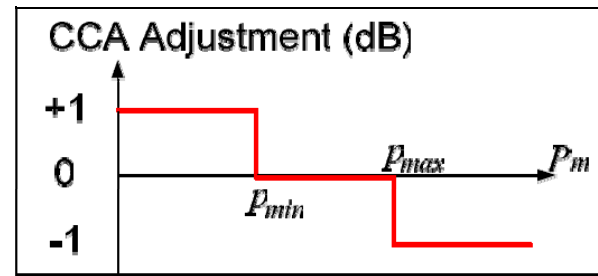
- ◆ most of the weak signal that causes strong-last collision will be from device in **co-channel cells**.
- ◆ Hence, let $P_r = RSSI$ be the **RS threshold**, and RSSI stands for **receive signal strength indicator**.
- ◆ However, signal strength is not constant.

$$P_r = \bar{s} - \sigma$$

CCA adaptation algorithm

- The maximum of measured PER values is used with a simple linear adaptation algorithm.

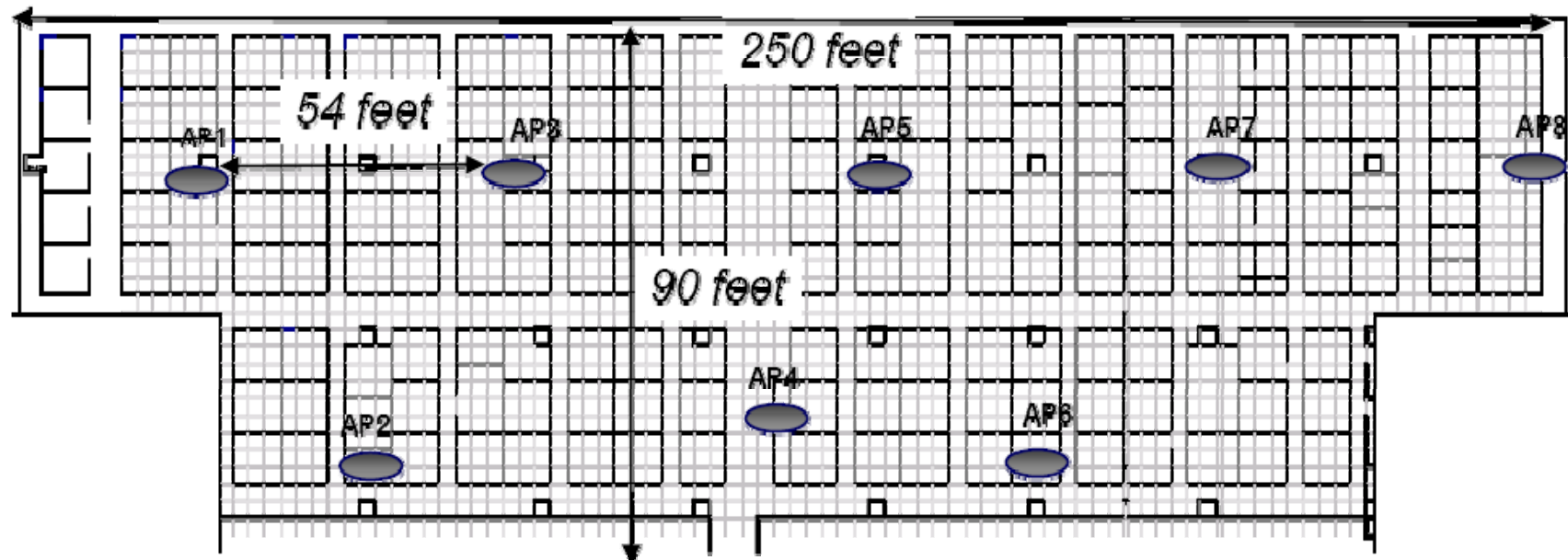
$$P_c = \begin{cases} \max \left(P_{c_{\min}}, P_c (dBm) \right) \\ P_c \\ \min \left(P_{c_{\max}}, P_c + \delta \right) \end{cases}$$



Experimental Topology

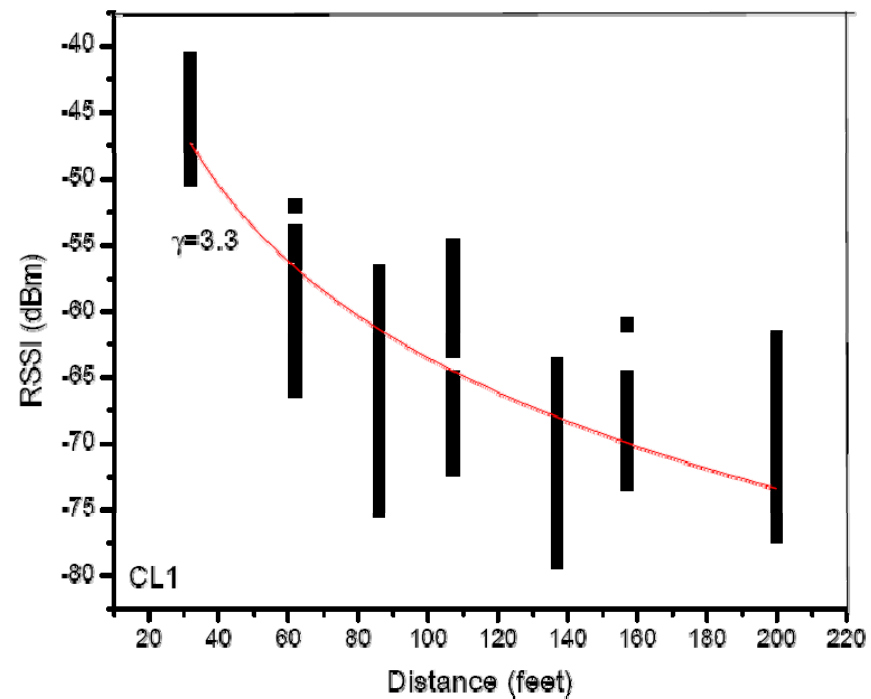
◆ Testbed Setup

- 8APs, (cisco Aironet 1130 802.11ABG)
- N clients with Centrino 2200 and WAG511(11a)



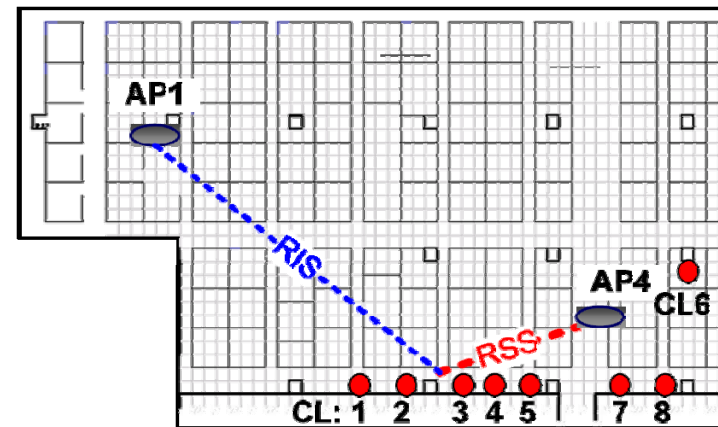
Experimental – Channel Characterization

- ◆ **6 clients** are deployed, one in **each corner** of the network.
- ◆ HD-WLAN is config. in 802.11g **channel 1** using **11dbm** as trans. power.
- ◆ CL: 3.3, 3.9, 3.3, 3.6, 3.9, 3.5.

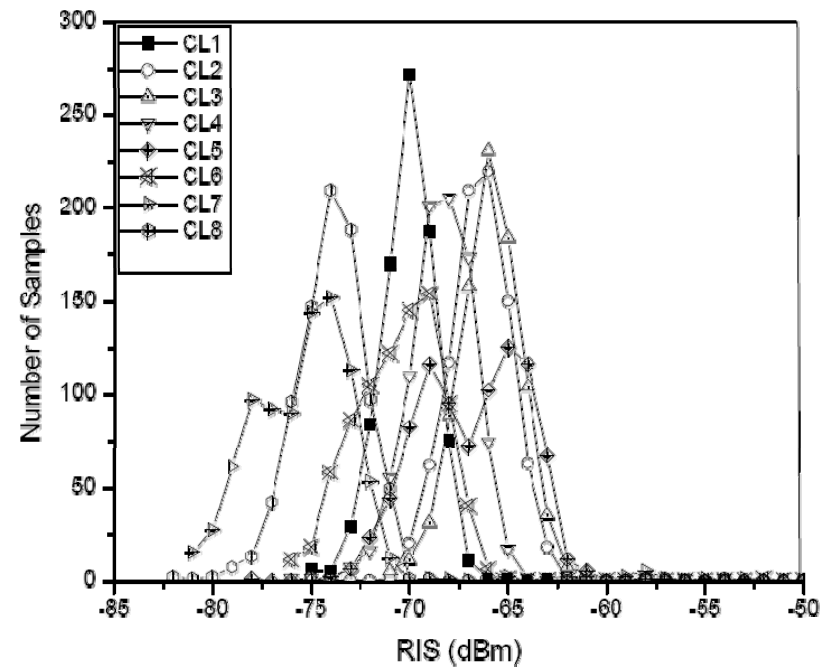
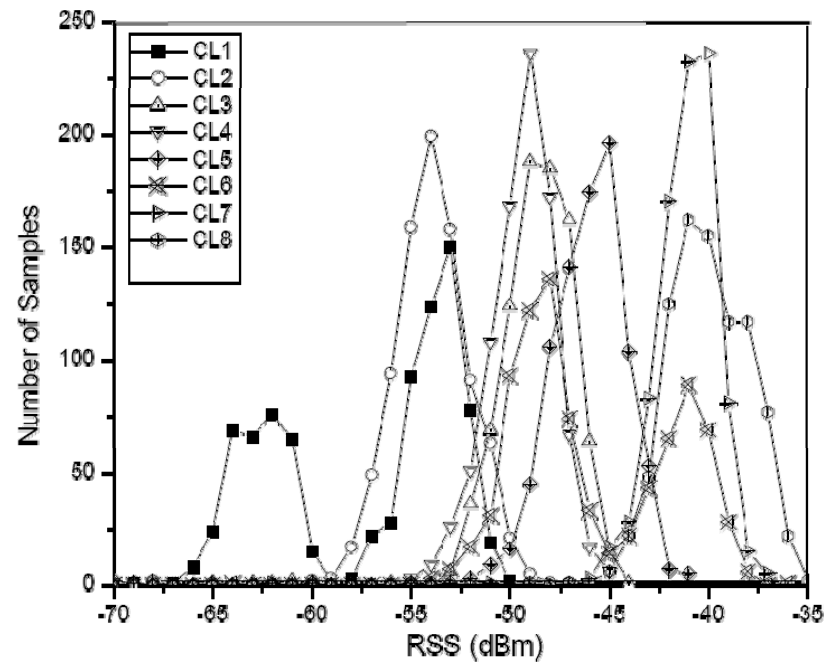


Channel Characterization

- ◆ Next, **CL1-8** are deployed to measure the RSSI between **AP1** and **AP4**.
- ◆ In each run, CL samples RSSI received from AP1 and AP4 with a **10-second** interval from **4000seconds**.



Results of Channel Characterization



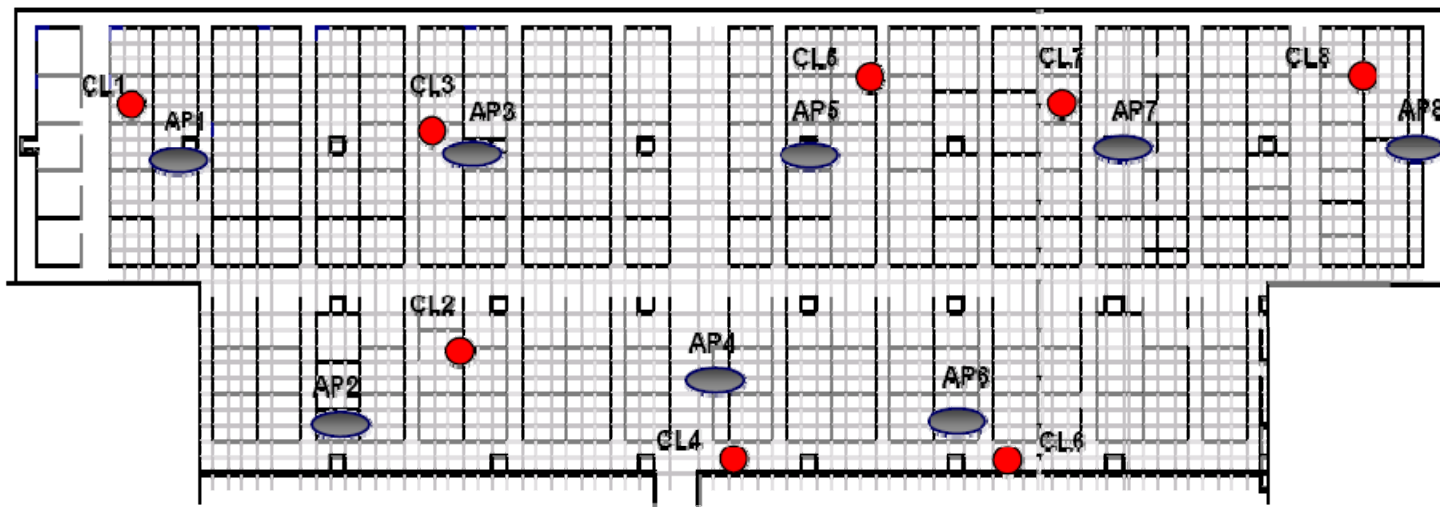
RS Adaptation

- ◆ Downlink, UDP traffic to all active CLs with packet size 1400bytes.

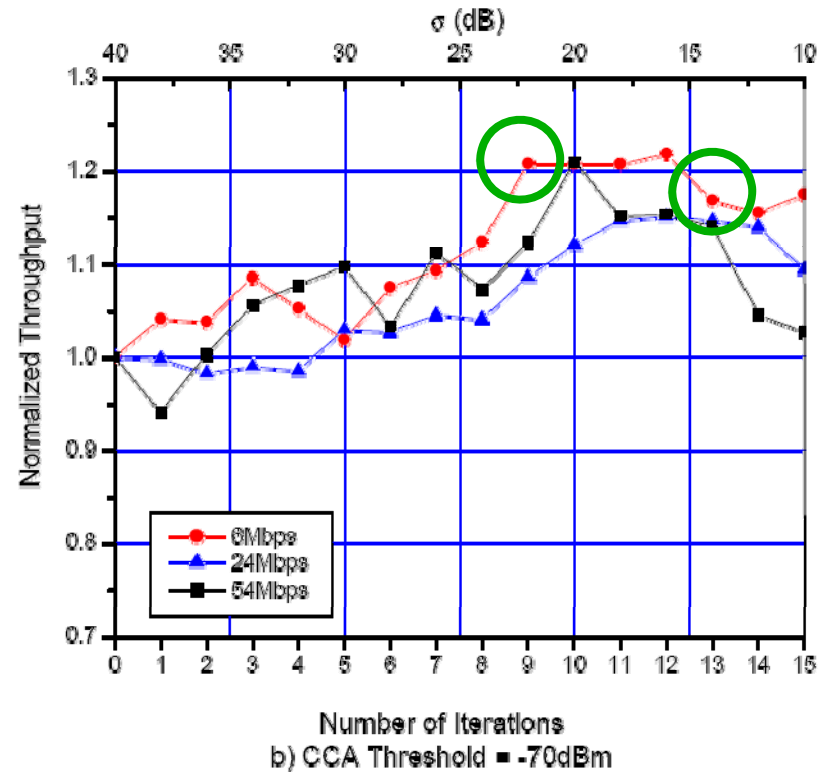
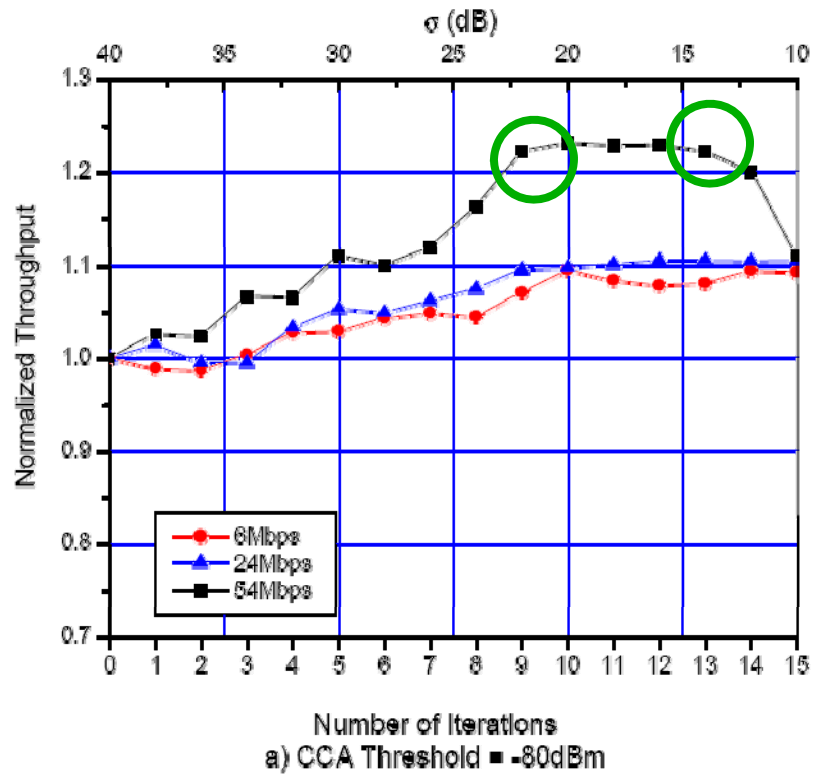
$$P_m = \max(p_i) |_{i=\{1,2,\dots,n\}};$$

$$Th = \frac{1}{n} \sum_i x_i;$$

$$Fi = \frac{(\sum_i x_i)^2}{n \sum_i x_i^2},$$



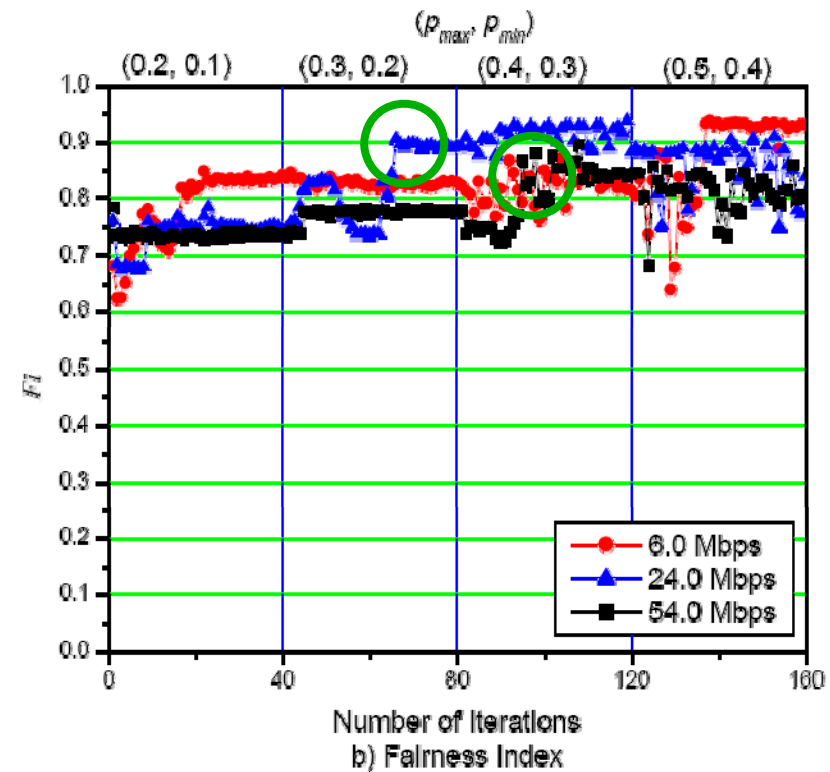
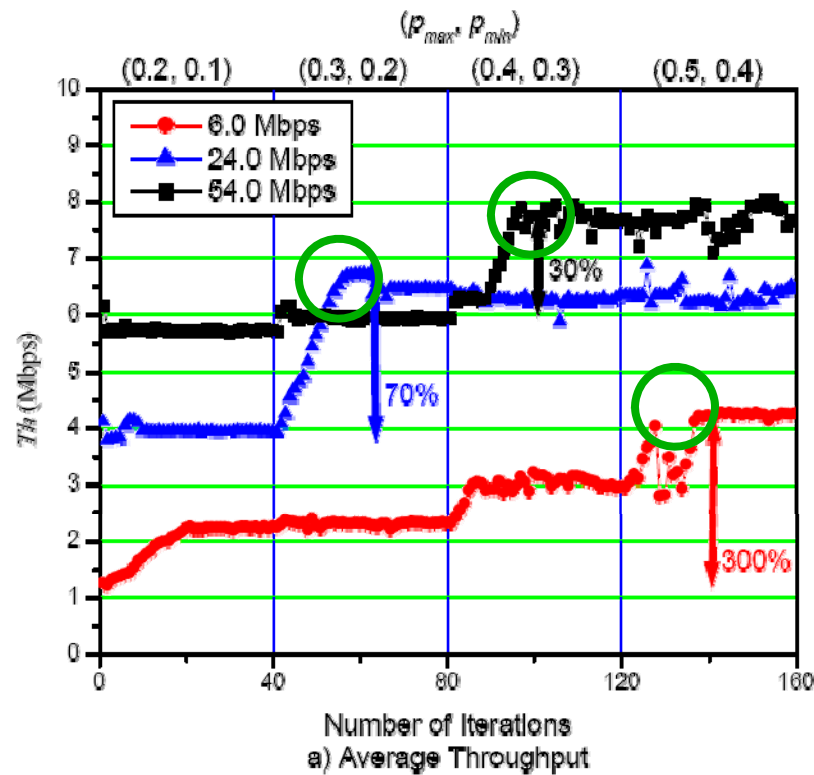
RS Adaptation Results



CCA Adaptation

- ◆ Next, we investigate the effect of the Pm target with CCA adaptation.
- ◆ Four targets
 - $(p_{\max}, p_{\min}) = \{(0.2, 0.1), (0.3, 0.2), (0.4, 0.3), (0.5, 0.4)\}$ are tested in sequence
 - with total 160 iterations and
 - each one staying 40 iterations.

CCA Adaptation results



Dynamic CSMA Scheme

