無線網路多媒體系統
Wireless Multimedia System

Lecture 8: Wireless TCP

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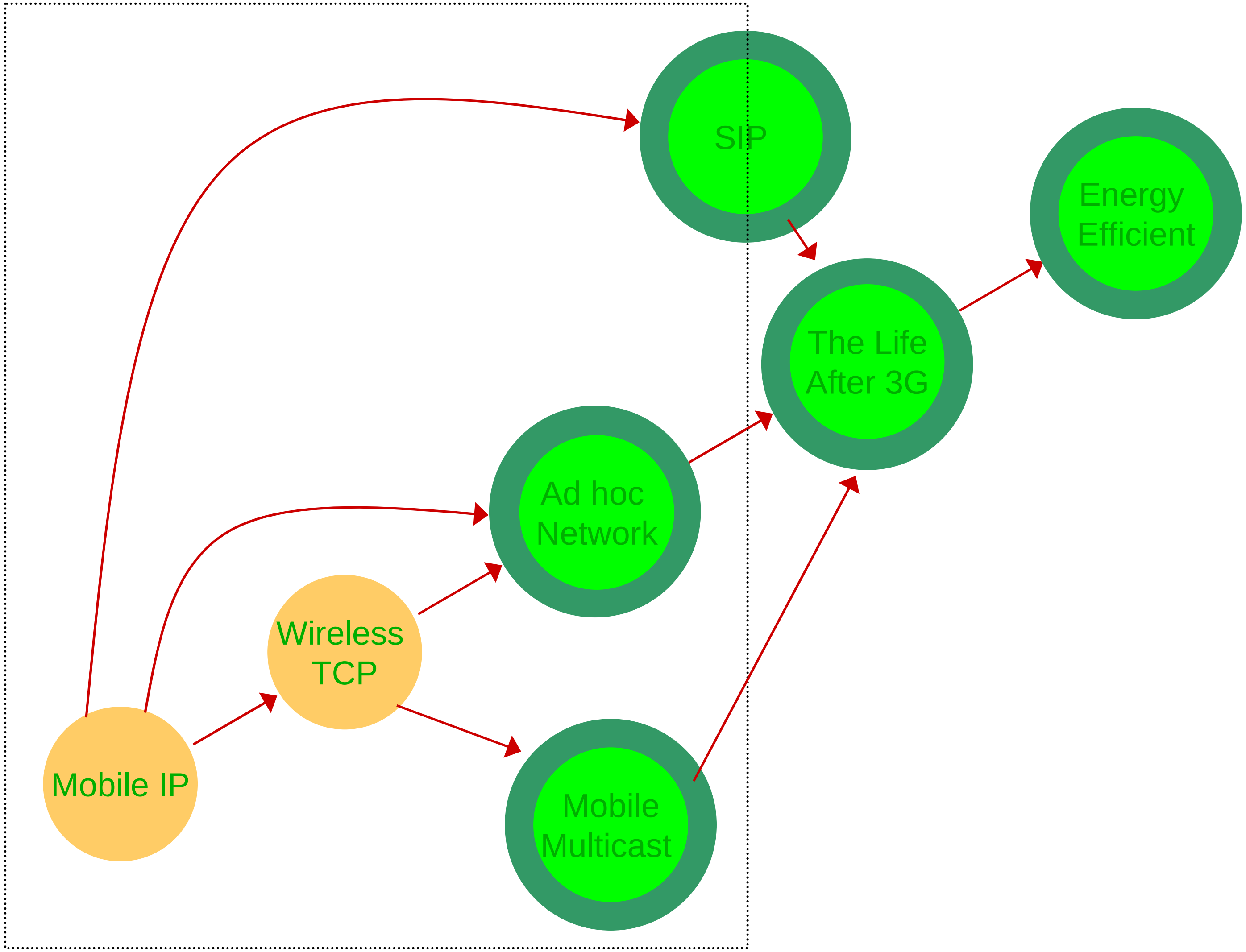
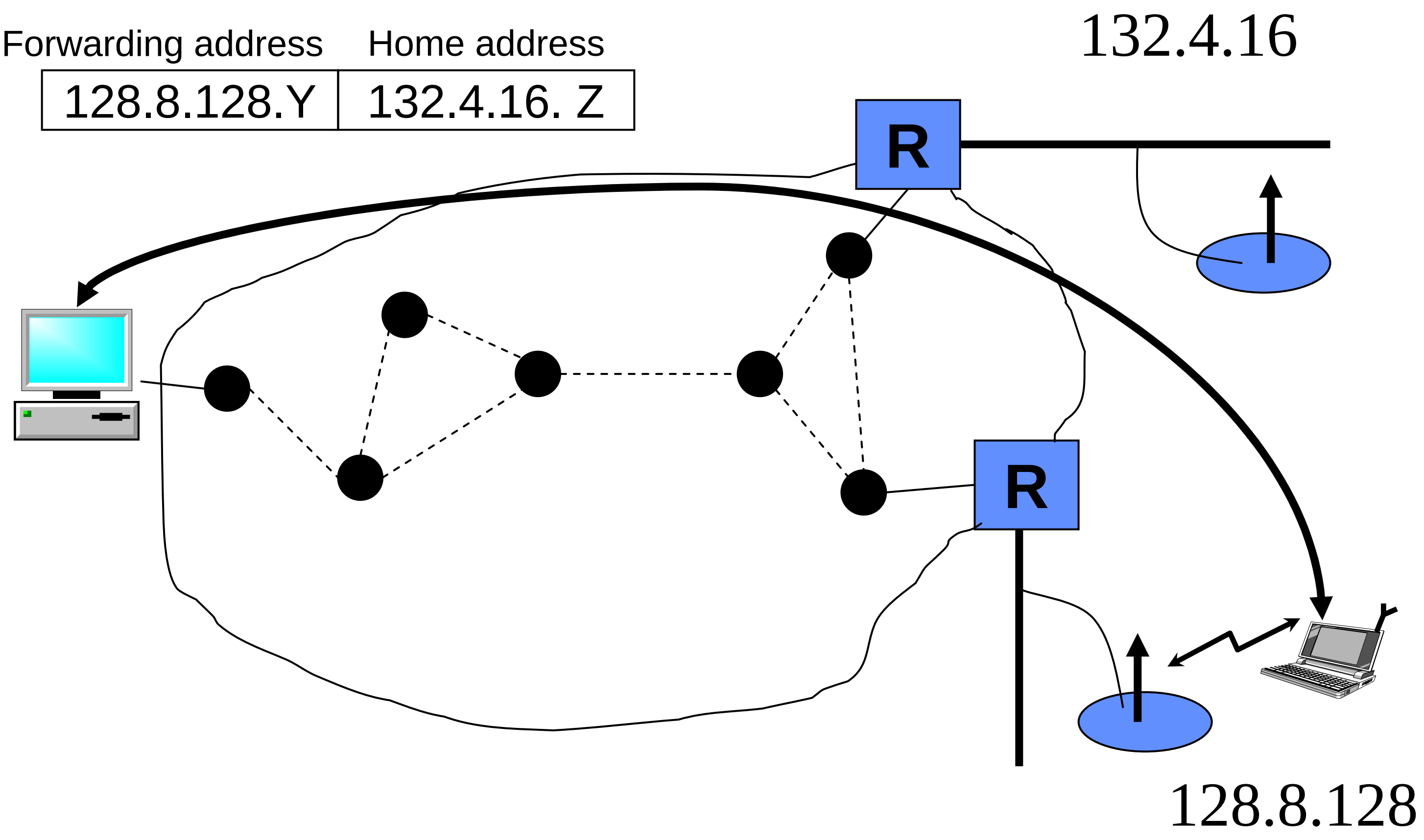
Coming Issues

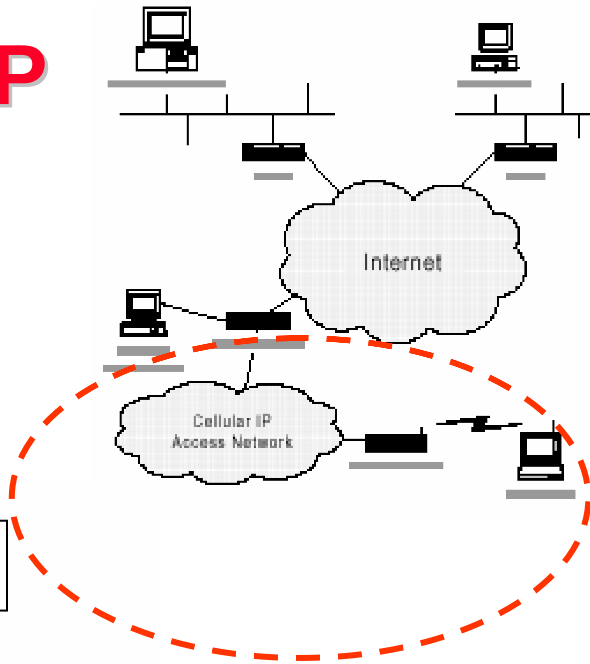


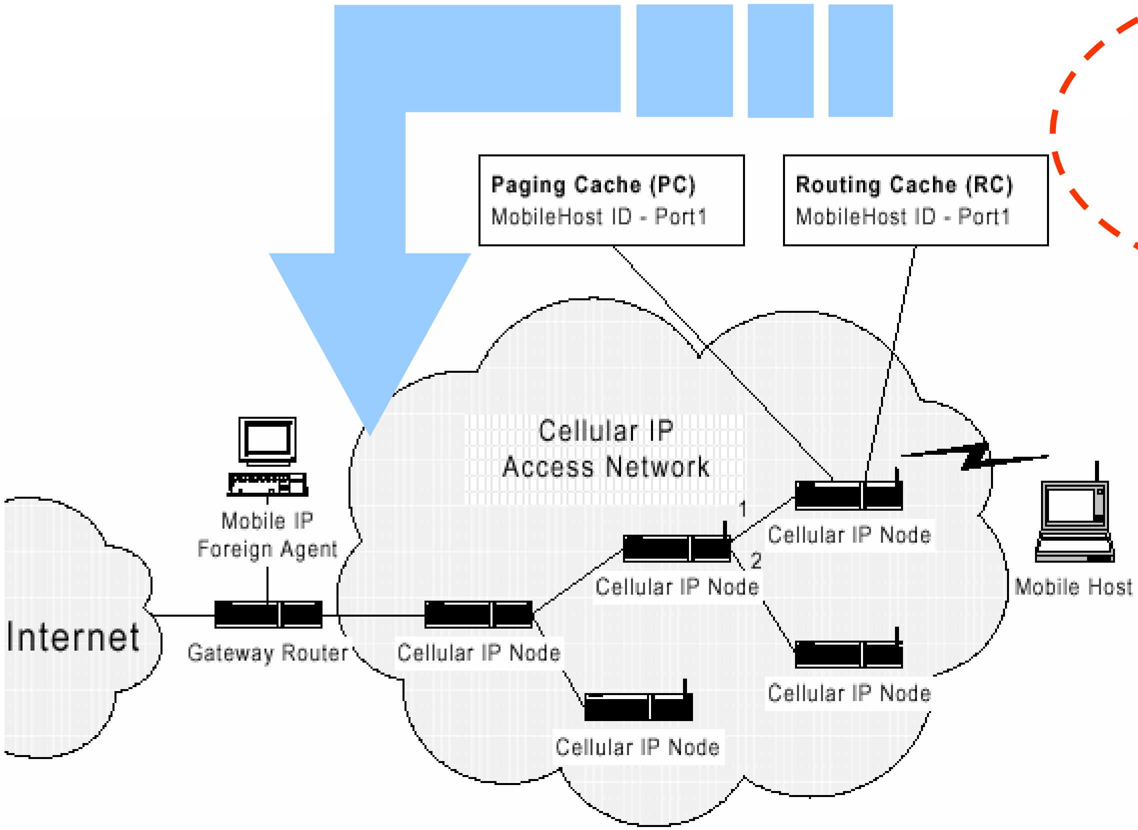
Mobile IP



Forwarding address
Home address
128.8.128.Y
132.4.16. Z
132.4.16
R
R
128.8.128

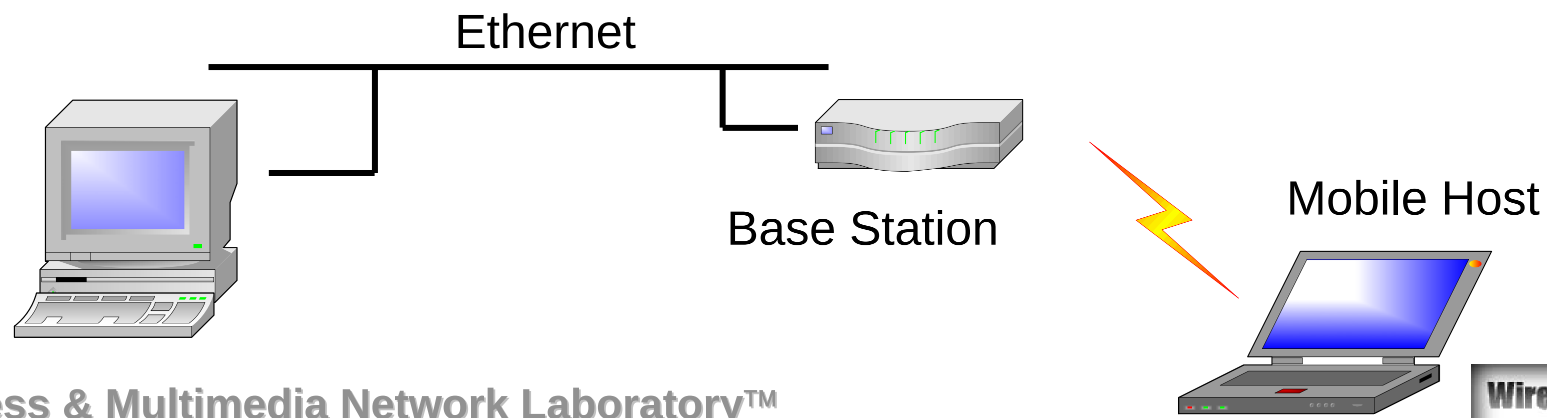
Cellular IP


Internet
Cellular IP
Access Network


Paging Cache (PC)
MobileHost ID - Port1
Routing Cache (RC)
MobileHost ID - Port1
Cellular IP
Access Network
Mobile IP
Foreign Agent
Internet
Gateway Router
Cellular IP Node
Cellular IP Node
Cellular IP Node
Cellular IP Node
Cellular IP Node
1
2
Mobile Host

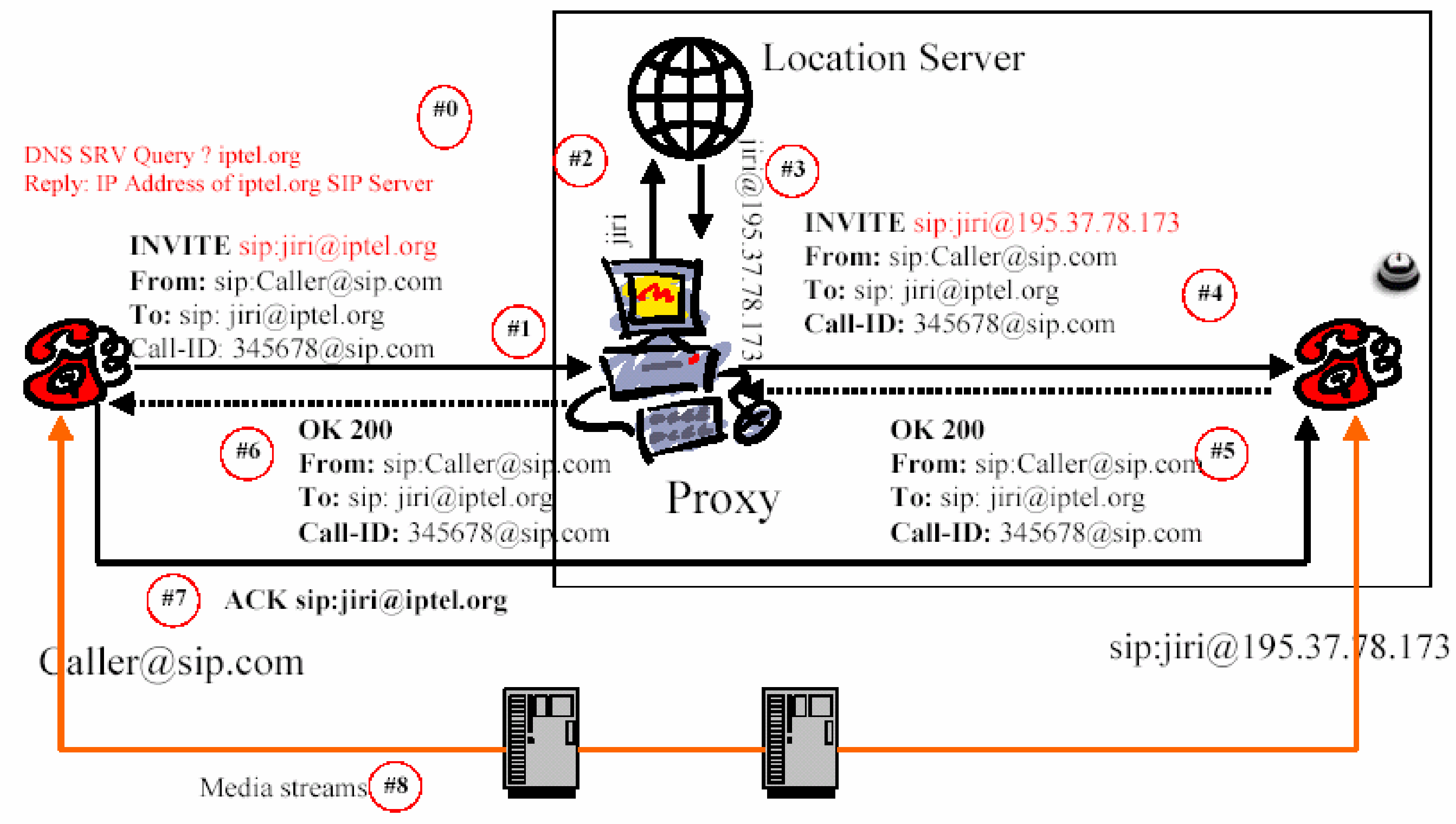
Wireless TCP

- TCP turned to perform well in traditional network where the packet losses occur mostly because of congestion.
- In the wireless environment
  - **Non-congestion** losses caused by wireless link
  - The degraded performance of TCP is mostly due to **mistaking wireless losses for congestion**.



SIP Protocol

Mobility support



Mobile Multicast

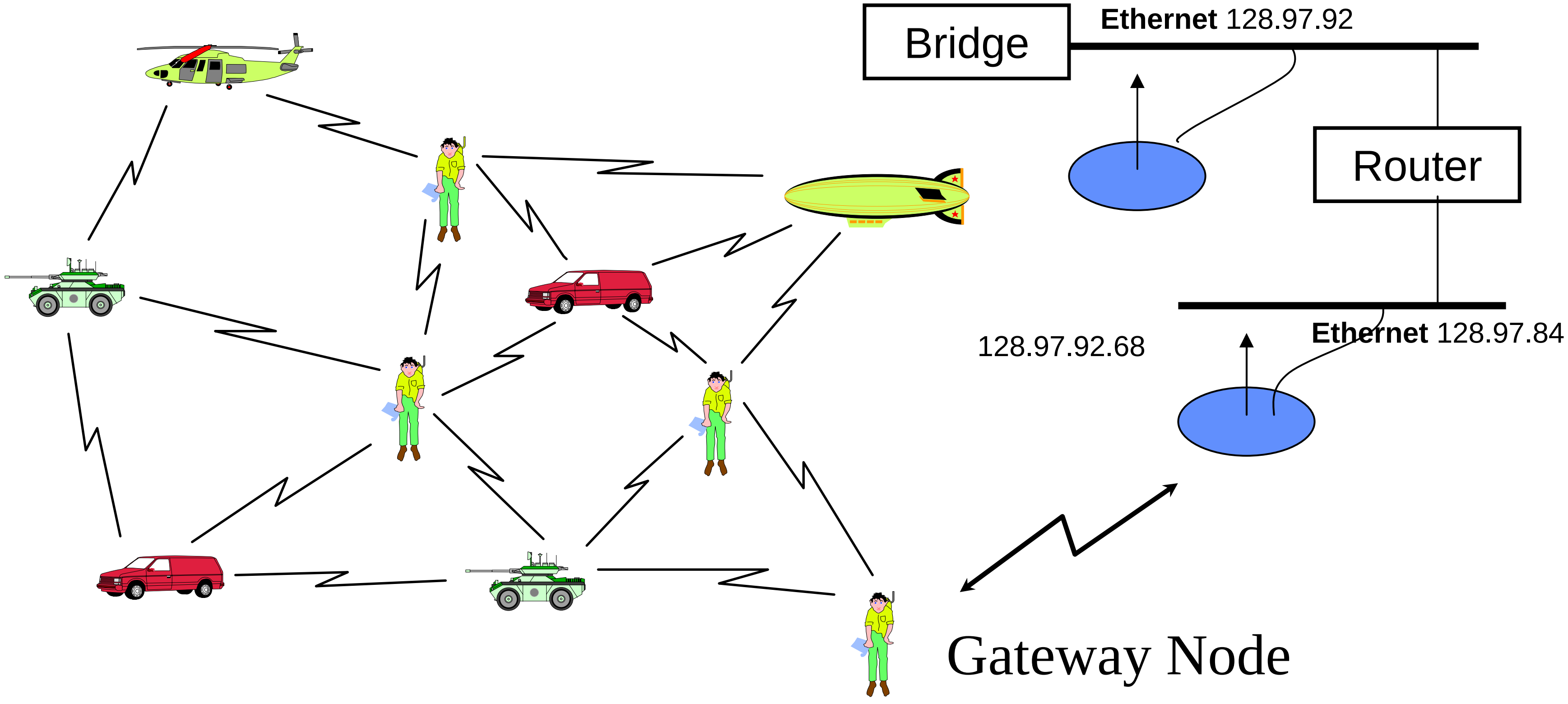
- Mobile Network~ Mobile IP
- Application Requirements: updates to replicated databases, Inter-process communication among cooperating processes
- Resource Conservations~ Single Copy in…Multicast IP

Mobile Multicast

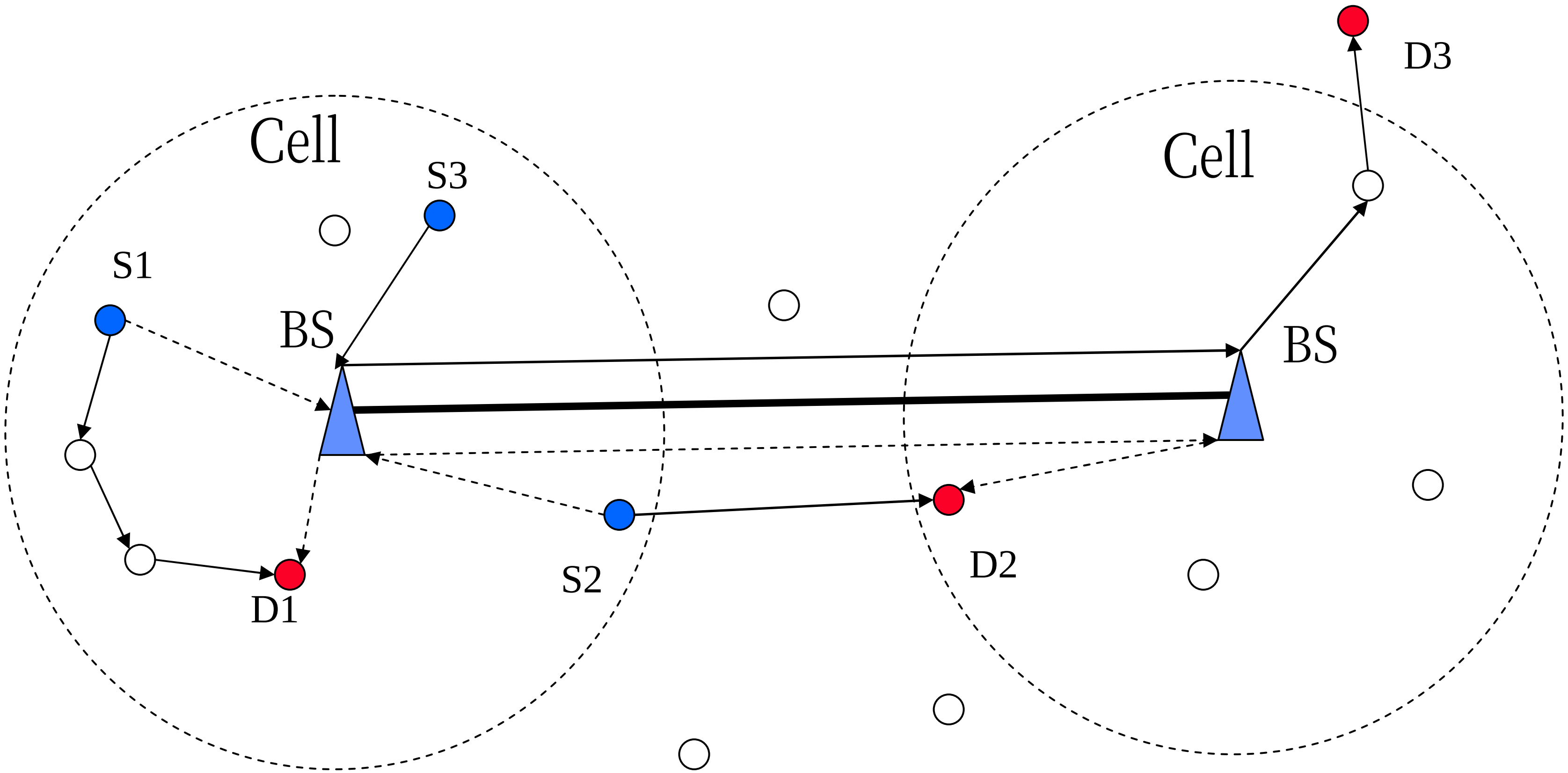
One to Many Mobile Multicasting Services

Internet Interconnection and Mobile IP

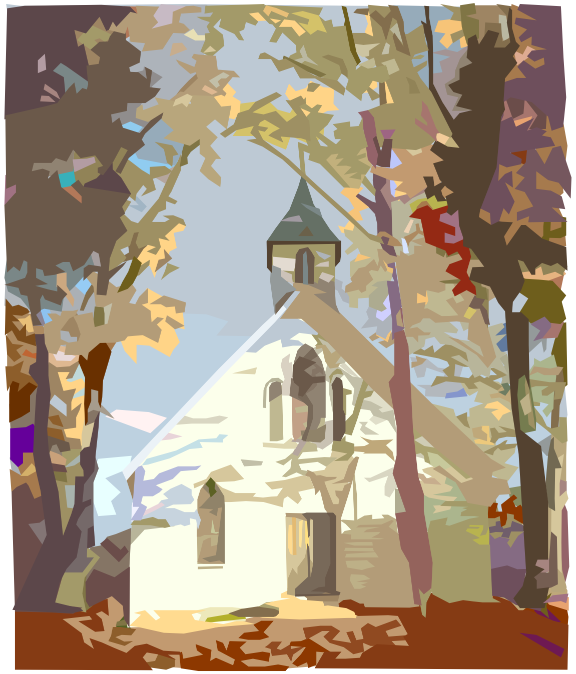
- DSR support the seamless interoperation between an ad hoc network and the Internet



Ad hoc & Cellular System



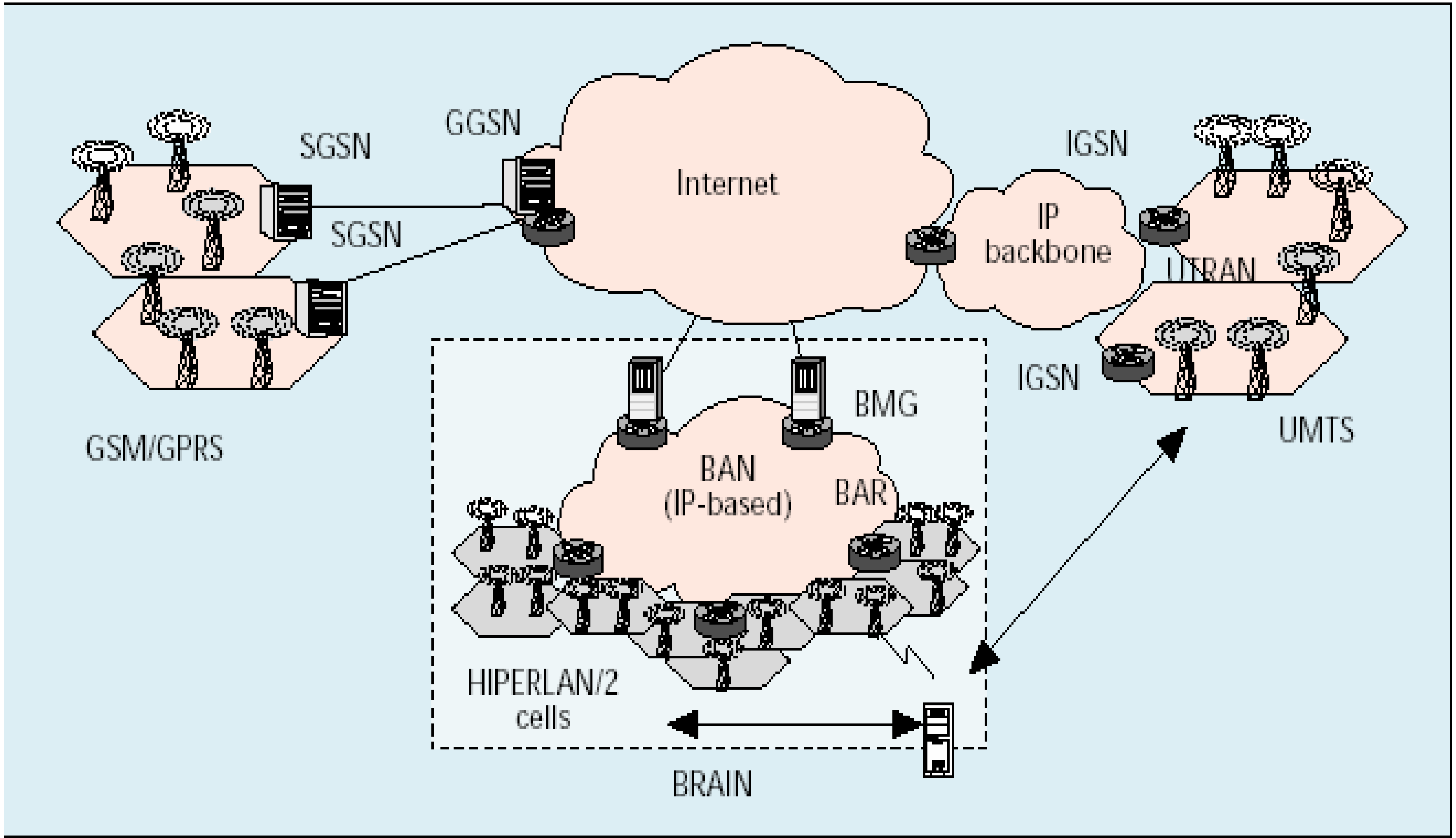
QoS Support for an All-IP System Beyond 3G



BRAIN

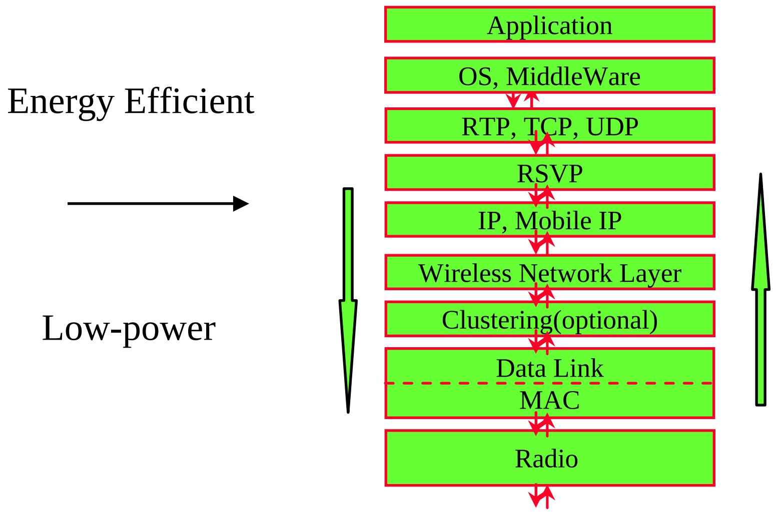
- Broadband Radio Access for IP-Based Networks
  - Cellular systems, fixed networks, and wireless LANs
  - Personal mobility, adapted for the terminal and link bandwidth
  - End-to-end QoS
  - A new QoS model for applications (BRENTA)
  - The radio link improvements
- IP-aware RAN (Radio Access Network)
  - Better support to IP applications
  - IP infrastructure will be widely available
- Protocol must be redesigned
  - Resource Management
  - Terminal mobility
  - RAN and terminal must have IP Stack

BRAIN (Broadband Radio Access for IP-based Network)



Energy and Power Efficient

- As wireless networks become an integral component of the modern communication infrastructure, energy efficiency will be an important design consideration due to the limited battery life of mobile terminals.
- This paper presents a comprehensive summary of recent work addressing energy efficient and low-power design within all layers of the wireless network protocol stack.



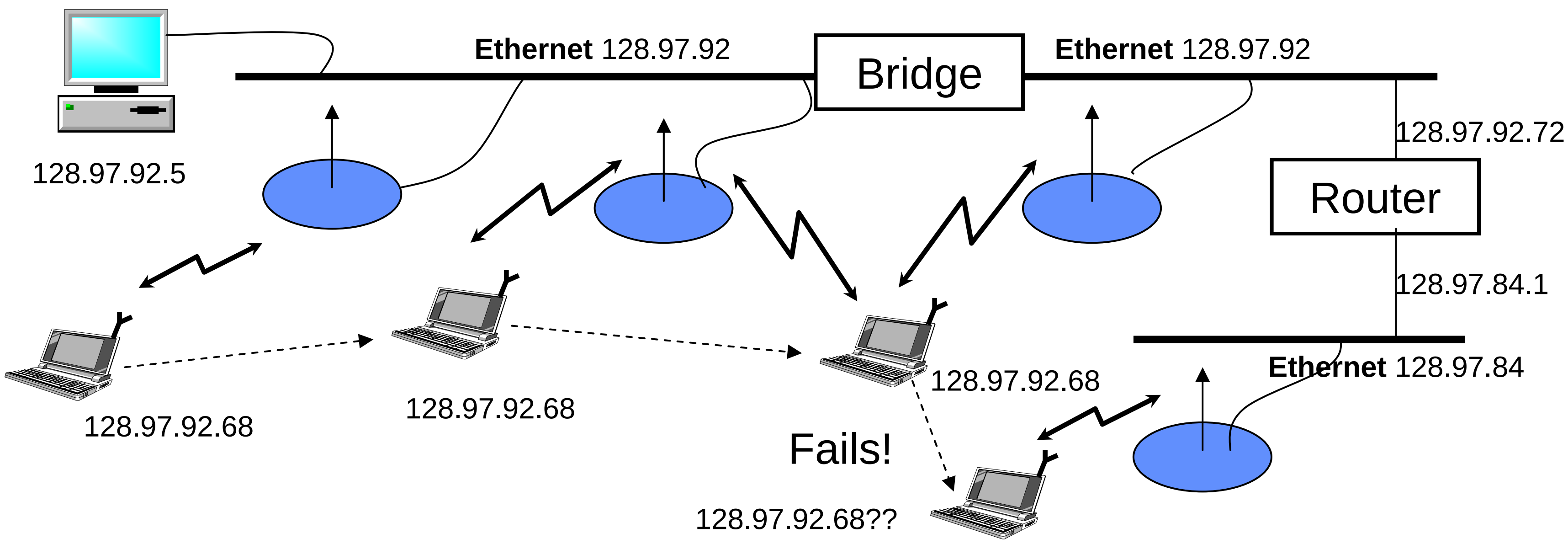
Agenda

- Basic TCP
- Impact of Mobility & Wireless on TCP performances
- Solutions for Wireless TCP
- Midterm (next week)

Reading

- [Balakrishnan95], Harri Balakrishnan, Srinivasan Seshan, Elan Amir and Randy H. Katz, “Improving TCP/IP Performance over Wireless Networks”, ACM Mobicom95
- [Balarkrishnan97], Harri Balarkrishna, Venkat N, Padmanabhan, Srinivasan Seshan and Randy Katz, “A Comparison of Mechanisms for Improving TCP Performance over Wireless Links”, IEEE JSAC 97.
- Reference: [Mario2001], Saverio Mascolo, Claudio Casetti,Mario Gerla, Renwang“TCP Westwood: Bandwidth Estimation for Enhanced Transport over Wireless Links”, Mobicom2001

Mobility in Wireless LANs: Basestation as Bridges

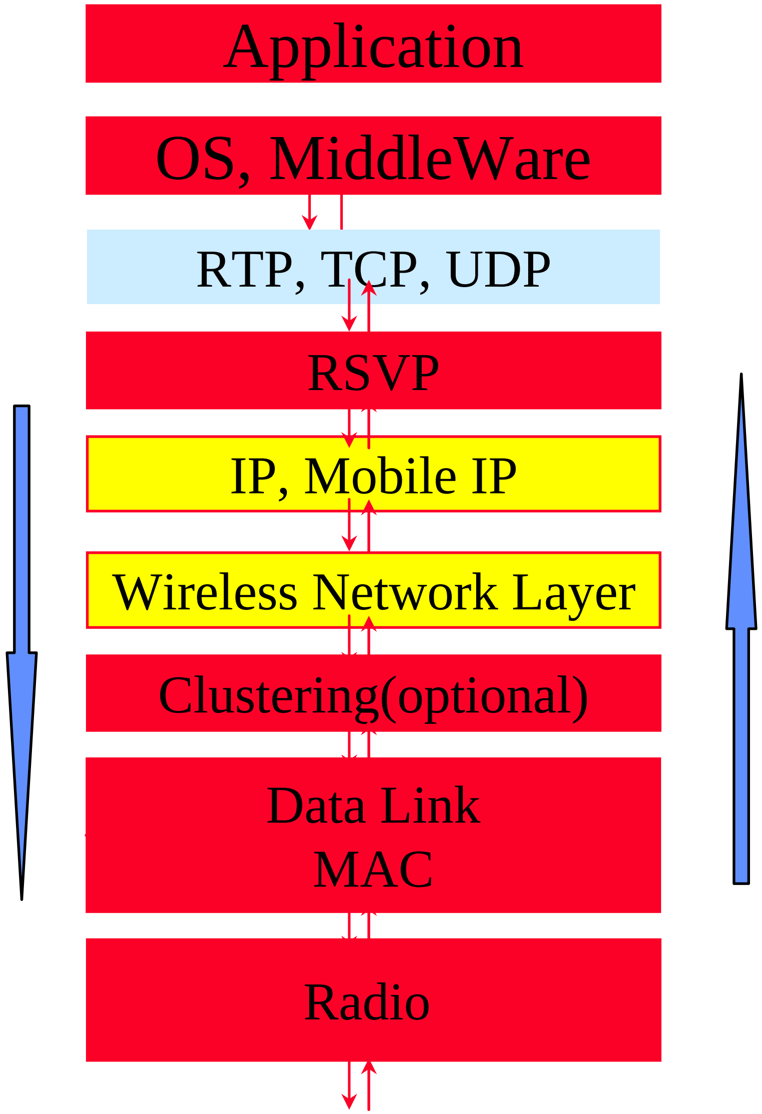


- Basestations are bridges(layer 2) – i.e. they relay MAC frames
  - Smart bridges avoid wasted bandwidth
- Works the within an ethernet(or other broadcast LAN)
  - Fails across network boundaries, and in switched LANs(e.g. ATM)

QoS and Multimedia Traffic Support

Adaptive
Algorithm

by QoS
Requirement



Mobility
Unpredictable
channel

by QoS
Information

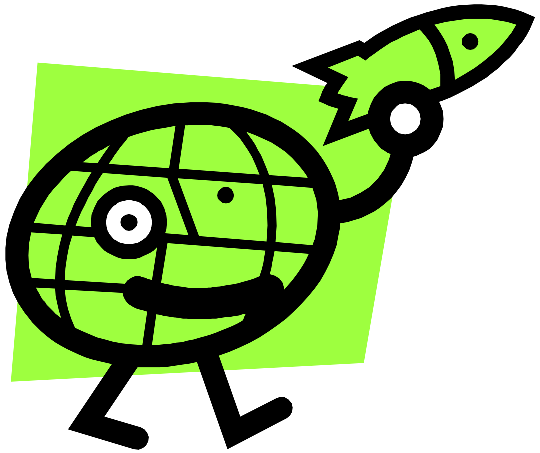
Background

- With the growth of wireless device, wireless network access will become popular, but…
- Import the protocol from the wire network to wireless network…
- Packet losses occur in wireless due to the lossy links, not network congestion
- In traditional TCP, it can not distinguish the difference between that lossy link and network congestion

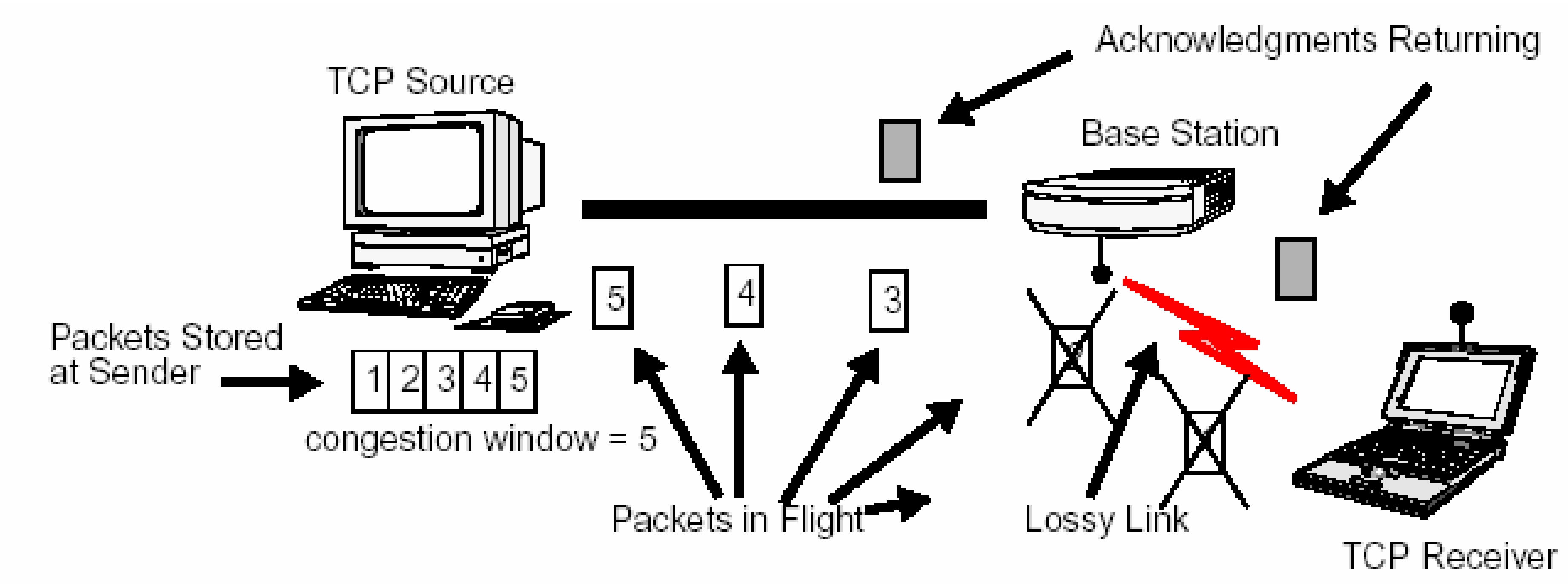
Characteristics of Wireless & Mobility

- Limited Bandwidth
  - Small frame sizes (MTU) to keep latency small
- High bit error rates
  - Small frame sized to keep packet loss probability small
- Time varying bit error rate
  - Fading, frequency collisions etc.
- QoS (loss rate, delay) degradation during hand-off
  - Due to network layer rerouting
  - Due to link layer procedures
- QoS degradation after hand-offs
  - Lack of resource at new basestation
  - Less optimal route

Basic End-to-End Control (Transport)



Typical loss situation



UDP (Connectionless, Unreliable)

Possible Multicast, Real Time Traffic, TCP-Friendly

Impact on Connectionless, Unreliable Transport Protocol

- Example: effect on UDP applications
- Increase in end-to-end packet losses
  - Error on wireless link
  - Packet loss during hand-offs
- Drop in application throughput
  - Errors on wireless link
  - Packet loss during hand-off
- Pauses in interactive applications
  - Burst errors on wireless link
  - Packet loss during hand-off
  - Delay increase due to buffering & re-sequencing during hand-offs
- Application level impact is much more complex!

TCP (Connection Oriented, Reliable)

Data Transmission, WWW, flow control, error control

TCP Basics

- Sliding window protocol: Go-Back N ARQ
  - Transfers a byte stream in “segments”, not fixed user blocks, logical timer associated with each segment that is sent
  - 32-bit sequence number indicated byte number in stream
    - Window is max number of outstanding unACK’ed bytes in network
- Cumulative acknowledgement scheme (original TCP)
  - Ack’s all bytes up through n
  - Piggybacked on data packets in reverse direction
- Control of sender’s window size
  - Min (receiver’s advertized window, congestion window)
  - Three goals
    - Flow control to avoid receiver buffer overflow
    - Congestion control to react to congestion in network layer & below
    - Congestion avoidance
- Segment loss is assumed to be a result of congestion in routers
  - Reasonable for wired network since BER on fiber is better than 10^{-12}

TCP's End to End Congestion Control

- Window-based congestion control
  - Cwnd: congestion window size
  - Ssthresh: slow start threshold (for slow down of increase)
- Timeout is an indicator of segment loss
- Timeout value
  - Using estimated average of ACK delay and expected deviation
- On timeout
  - Segment is assumed lost and is attributed to congestion
  - One-half of current window Is recorded in ssthresh
  - Cwd is reduced to 1
  - Timeout value is increased in case packet was delayed

TCP's End-to-end Congestion Control

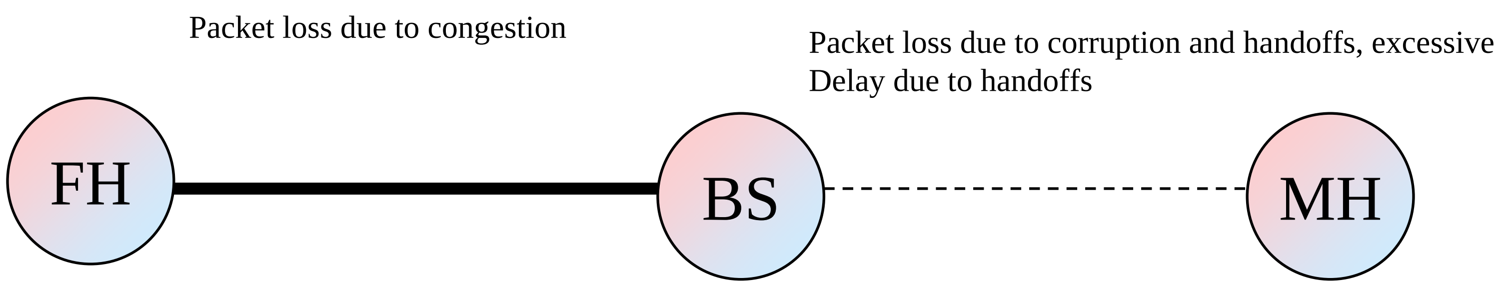
- On new ACK
  - Everything okay, so allow larger congestion window
  - Two ways of increasing cwnd
    - Phase1: slow start until cwnd <= ssthresh
      - Fast (exponential) increase of cwd
    - Phase2: congestion avoidance
      - Slow (additive) increase of cwnd
- Duplicate ACKs
  - Two causes: lost segment, misordered segment
  - >=3 duplicate ACKs in a row are a good indication of a lost segment but data is still flowing
  - Fast Retransmit and Fast Recovery
    - Missing segment is retransmitted without waiting for timeout
    - One half of current window is recorded in ssthresh
    - Congestion avoidance is done but not slow start

Challenges of Mobility and Wireless on Network Performance

TCP Performance

The Problem

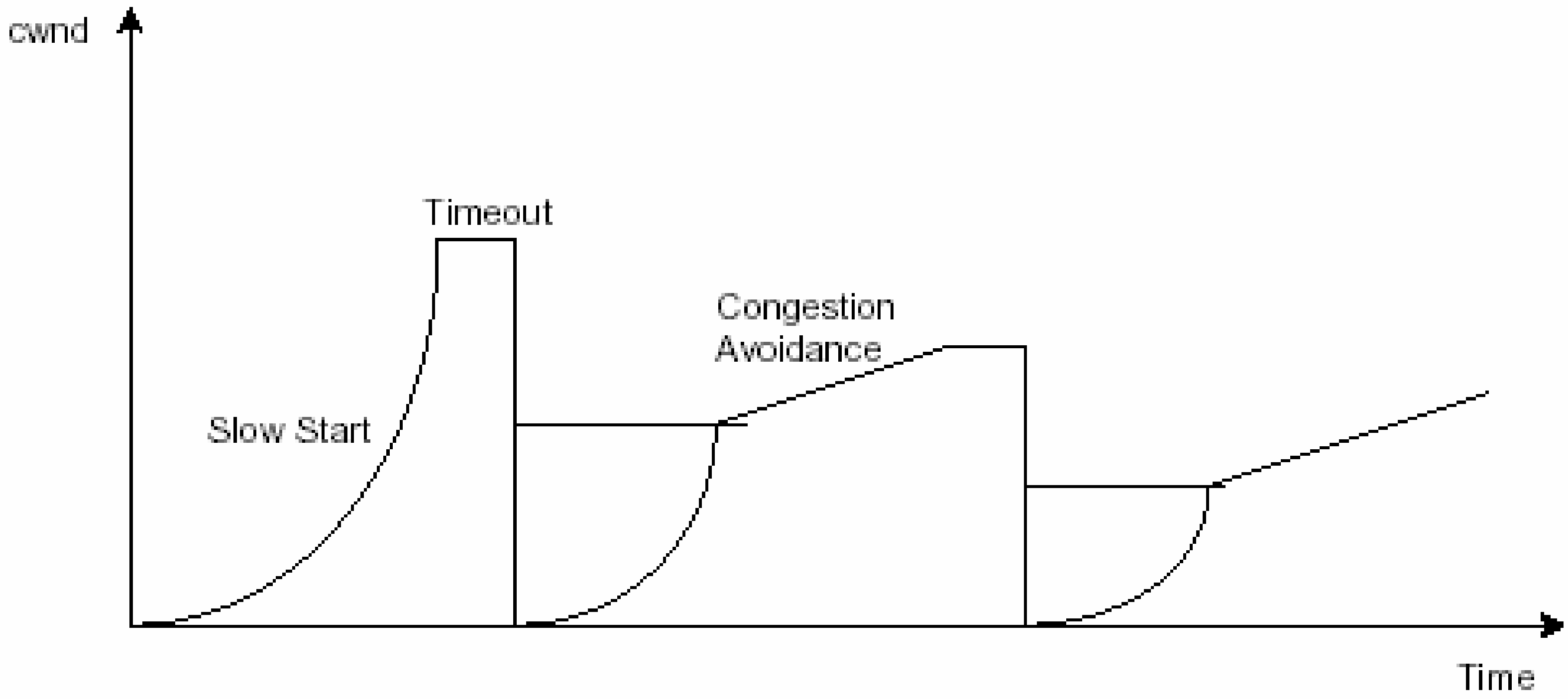
- In Wireless and mobile networks, segment loss is likely not due to congestion
  - Packet corruption due to high BER on wireless link (noise, fading)
  - Packet delay and losses during handoffs
- But, TCP invokes congestion control nevertheless
- Mistaking wireless errors and handoffs for congestion causes
  - Significant reductions in throughput (window size decreases, slow start)
  - Unacceptable delays (low resolution TCP times ~500ms, back-off)



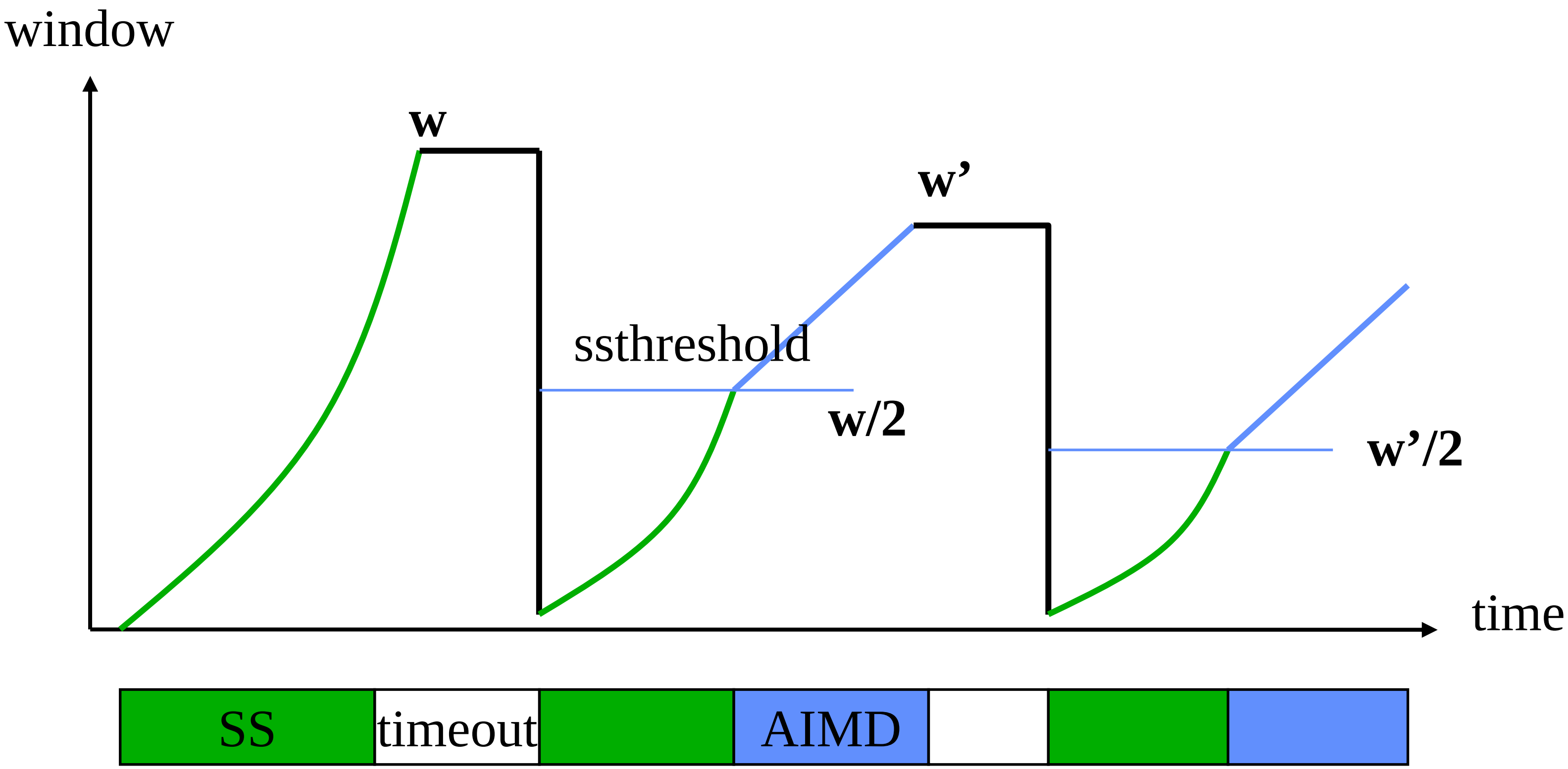
Example graph

cwin<=ssthresh → slow start

cwin>ssthresh → congestion avoidance

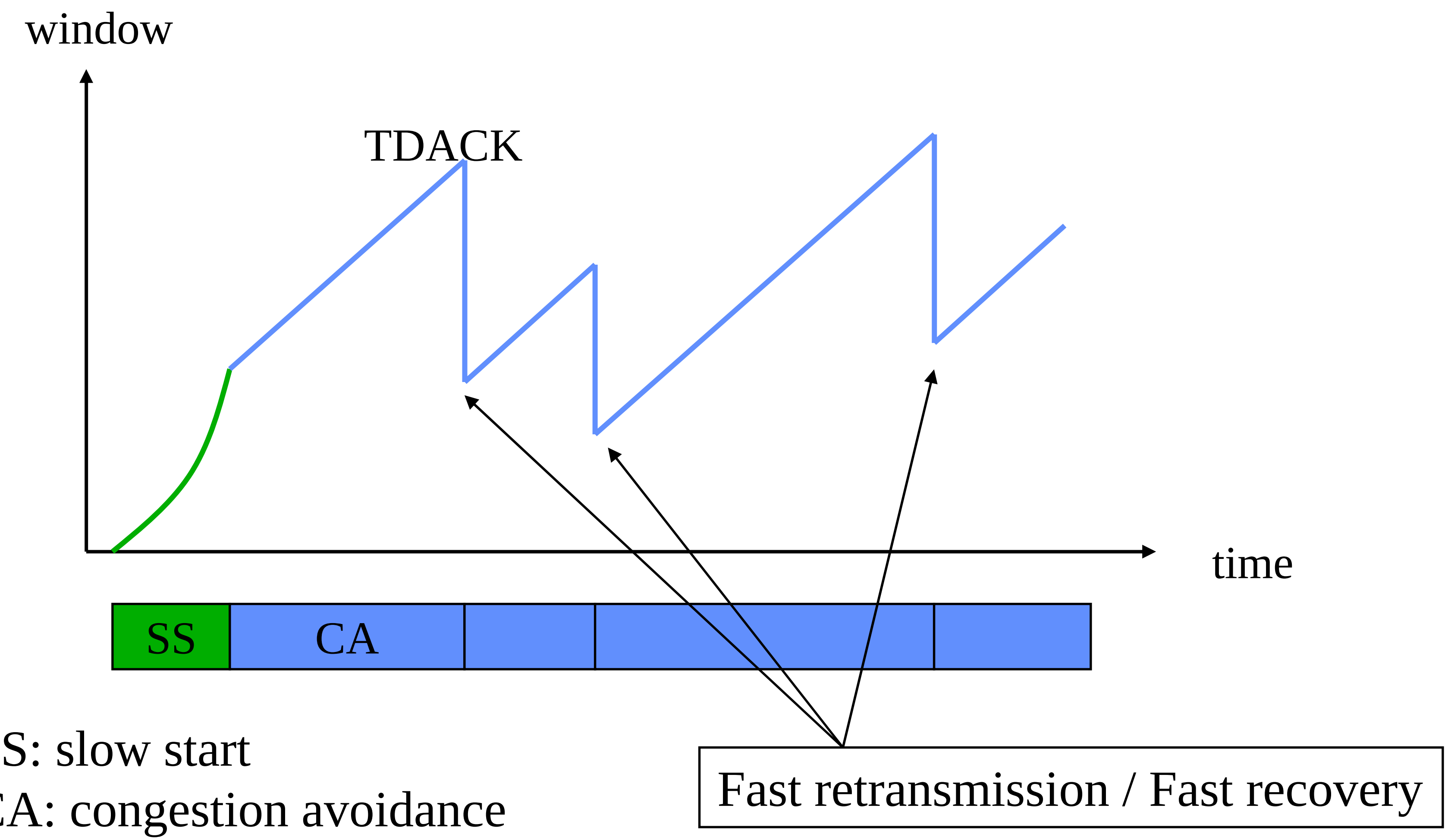


Slow Start of TCP Reno



ssthreshold : slow-start threshold

Congestion Avoidance of TCP Reno



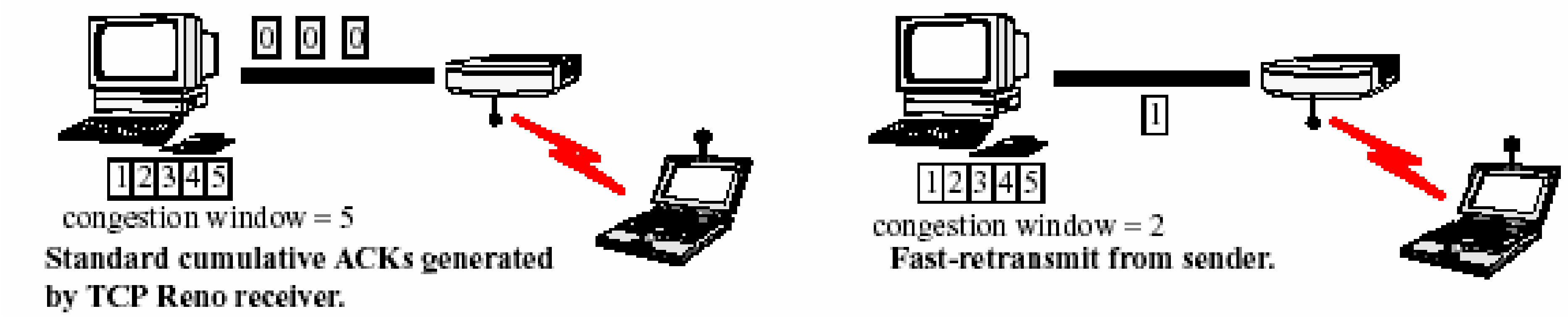
SS: slow start

CA: congestion avoidance

Fixes?

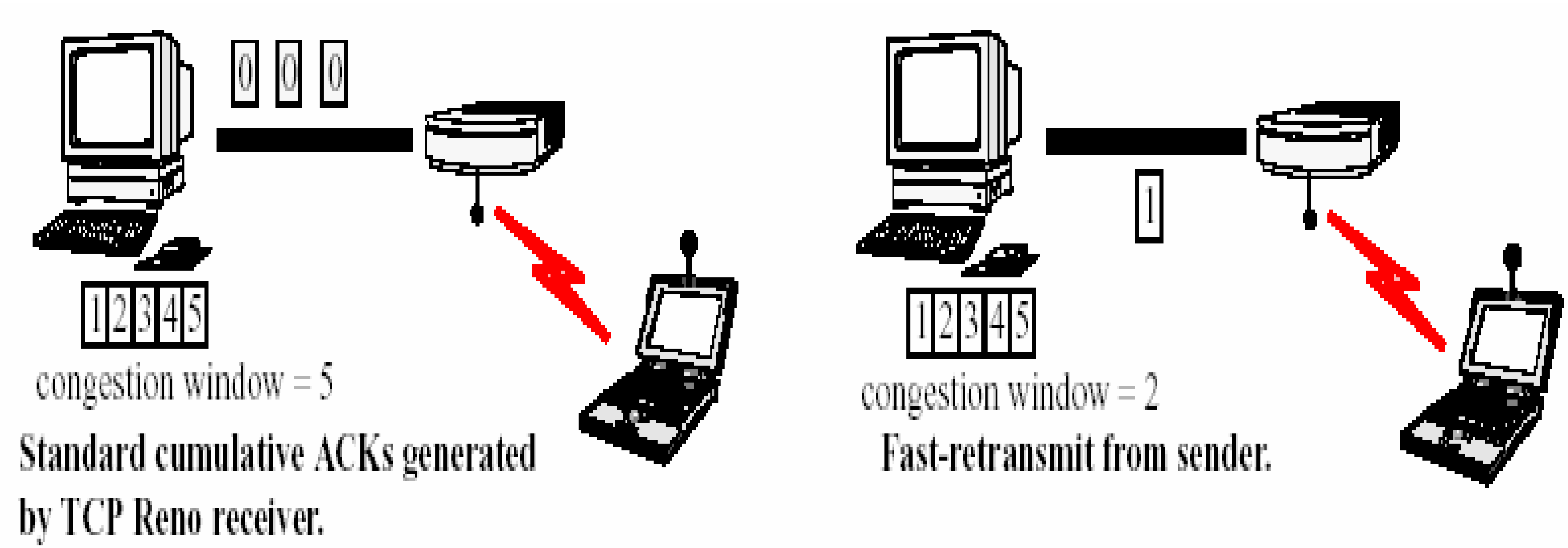
- Fix TCP
  - TCP really a hack in many ways..
  - Separate congestion control from error control
  - Move away from cumulative ACK
- Fix lower layer to make TCP work better
  - Improve the wireless link
- Use something different
  - Something totally new
  - Something different for the wireless part

Normal TCP



0
0
0
1 2 3 4 5
congestion window = 5
Standard cumulative ACKs generated by TCP Reno receiver.
1
1 2 3 4 5
congestion window = 2
Fast-retransmit from sender.

Fast-Retransmit Scheme



Solutions for WTCP (I)

Split the connection into two parts

Split Connection Approaches

- Main Idea: split MH ↔ FH connection into two MH ↔ BS & BS ↔ FH
  - Separate flow control and reliable delivery mechanisms
  - Intermediate higher layer agent at the base-station
  - Session layer hides the split connection
- Two approaches:
  - Both FH ↔ BS & BS ↔ MH segments use TCP: Rutger’s Indirect-TCP
    - e.g. uses MTCP (Multiple TCP) over BS ↔ MH
  - BS ↔MH uses specialized protocol
    - e.g. uses SRP (Selective Repeat) over BS ↔MH
    - Error and flow control optimized for lossy wireless link

Pros & Cons of Split-Connection Approaches

- Pros
  - FH is shielded from wireless link behavior
  - Handoff is transparent to FH
  - Relative easy to implement
  - Requires no modification to FH
  - Can use specialized protocol over wireless link
- Cons
  - Loss of end-to-end semantics
  - Application relink with new library
  - Software overhead: efficiency and latency
  - Large handoff latency

Solutions for WTCP (II)

Lower layer to make TCP work better

Link-level Error Control

- FEC and ARQ on wireless link to increase its reliability
  - Improves performance independent of transport protocol
- Disadvantage
  - Coupling between link level and end-to-end retransmission may lead to degraded performance at high error rates
  - Does not address the delay and losses due to handoffs

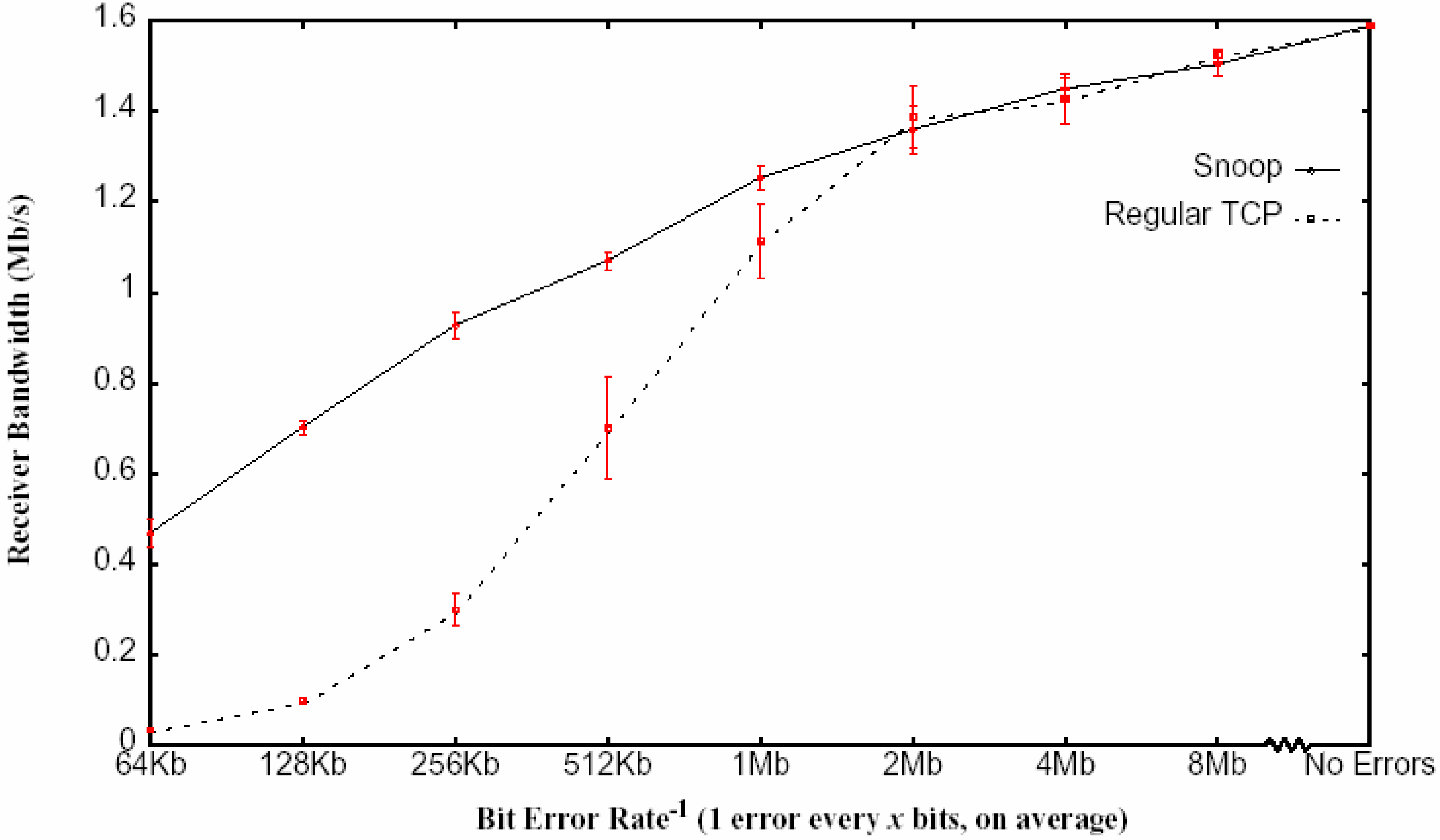
Solutions for WTCP (III)

Snoop, Make it look like!

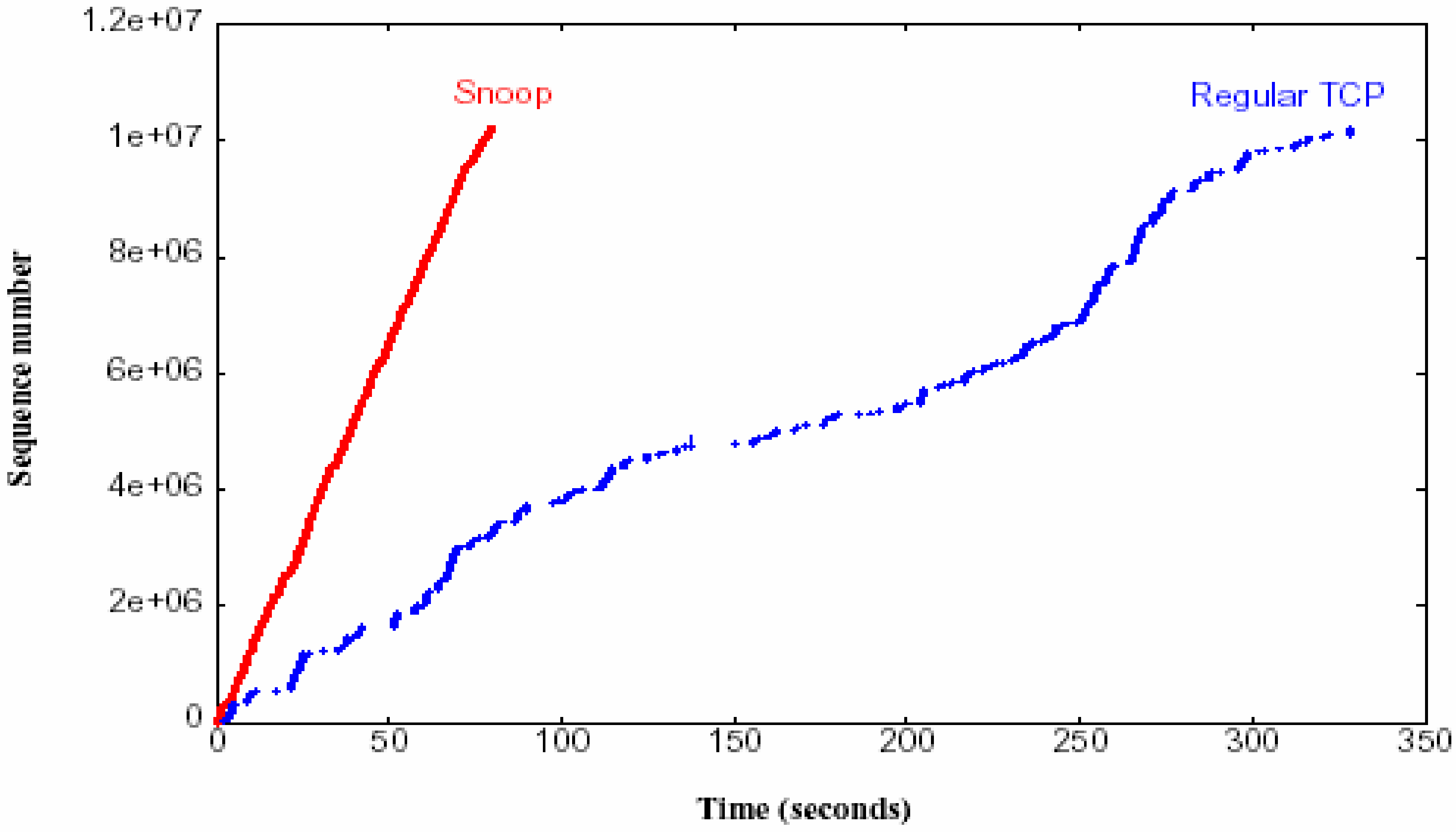
Snoop TCP

- Basic Idea for transfer of data to MH
  - Snoop Module: Modify network layer routing code at BS
  - Cache un-acknowledged TCP data going to MH at BS
  - Perform local retransmissions over wireless link
    - Policies to deal with ACKs from MH and timeout
    - Used duplicate ACKs to identify packet losses
  - Shields sender from wireless link
    - Transient conditions of high BER, temporary disconnection
- Basic idea for transfer of data from MH
  - BS detects missing packets and generated NACKs for MH, expoits SACK option for TCP
  - MH re-sends the packets, requires modifying TCP code at MH
- Features
  - Speedups of up to x20 over regular TCP depending on bit error rate
  - Maintain end-to-end semantics
  - Does not address the handoff problem

Performance of the Snoop Mechanism



Performance of the Snoop Mechanism



Comparison of Wireless TCP Techniques

- End-to-End proposals
  - Selective ACKs
    - Allows sender to recover from multiple packet losses without resorting to course timeout
  - Explicit Loss Notification (ELN)
    - Allow sender to distinguish between congestion vs. other losses
- Split-connection proposal
  - Separate reliable connection between BS & MH
    - May use standard TCP or, special techniques such as SACK, or NACK
- Link-layer proposal
  - Hide link-layer losses via general local retransmission and FEC
  - Make link-layer TCP aware
    - Snoop agent to suppress duplicate ACKs

Main Conclusions of [Balakrishnan97]

- Simple link layers do not quite work
  - Adverse interaction of times is actually a minor problem
  - Fast retransmission and associated congestion control gets triggered and cause performance loss
- Reliable link layer with TCP knowledge works well
  - Shielding sender from duplicate ACKs due to wireless losses improves throughput by 10-30%
- No need to split end-to-end connections
  - I-TCP does as bad because sender stalls due to buffer space limit at BS
  - Using SAK or BS-MH link works well
- SACK and ELN helps significantly
  - Help avoid timeous
  - e.g. ELN helped throughput by x2 over vanilla TCP-Reno
  - But still do 15% to 35% worse than TCP-aware link layer schemes

Introduction

- TCP Westwood (TCPW) is a sender-side modification of TCP Reno in wire as well as wireless network
- TCPW can estimate the E2E b/w and the improvement is most significant in wireless network with lossy links
- TCPW sender monitors the ACK reception and from it estimates the data rate
- The sender uses the b/w estimate to properly set the cwin and ssthresh

Filtering the ACK reception rate

- Sample of bandwidth

- We $$b_k = \frac{d_k}{t_k - t_{k-1}}$$ ter to average sampled
  mea